Induction, Episode VI: Return of the I.H.

Part a: Recursive Definition

Ian Ludden

By the end of this lesson, you will be able to:

• Understand how to read a recursive definition, e.g., compute selected values or objects produced by that definition.

- Understand how to read a recursive definition, e.g., compute selected values or objects produced by that definition.
- Know that a recursive definition, and an inductive proof, require both a base case and an inductive step/formula.

- Understand how to read a recursive definition, e.g., compute selected values or objects produced by that definition.
- Know that a recursive definition, and an inductive proof, require both a base case and an inductive step/formula.
- Define the Fibonacci numbers.

Recursion: divide problem into smaller problem(s)

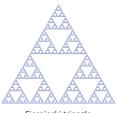
- Recursion: divide problem into smaller problem(s)
- Coding (mergesort and other "divide-and-conquer" algorithms)

- Recursion: divide problem into smaller problem(s)
- Coding (mergesort and other "divide-and-conquer" algorithms)
- Art (esp. fractal art)



Droste effect

- Recursion: divide problem into smaller problem(s)
- Coding (mergesort and other "divide-and-conquer" algorithms)
- Art (esp. fractal art)
- Closely related to induction



Sierpinski triangle



Droste effect

Some Classic Examples

Some Classic Examples

Example 1: Factorial base \mathbb{C} pecursive formula Define $f: \mathbb{N} \to \mathbb{N}$ by f(0) = 1 and $f(n) = n \cdot f(n-1)$ for n > 0. $f(S) = 5 \cdot f(\mathbb{H}) = 5 \cdot \mathbb{H} \cdot f(3) = \dots = 5!$ $f(n) = \begin{cases} f(n-1) & \text{otherwise} \\ f(n-1) & \text{otherwise} \end{cases}$ $f(n) = \begin{cases} f(n-1) & \text{otherwise} \\ f(n-1) & \text{otherwise} \end{cases}$

Some Classic Examples

Example 1: Factorial

Define $f : \mathbb{N} \to \mathbb{N}$ by f(0) = 1 and $f(n) = n \cdot f(n-1)$ for n > 0.

Example 2: Fibonacci two bake cases

Define
$$F : \mathbb{N} \to \mathbb{N}$$
 by $F(0) = 0$, $F(1) = 1$, and $F(n) = F(n-1) + F(n-2)$ for $n > 1$.



Preview: Connection to Induction

Preview: Connection to Induction

Example 3: Implicit Summation

Define $g: \mathbb{Z}^+ \to \mathbb{Z}^+$ by

$$g(n) = \begin{cases} 2 & \text{if } n = 1 \\ n(n+1) + g(n-1) & \text{otherwise.} \end{cases}$$

$$g(n) = h(n+1) + (h-1)(n-1+1) + g(n-2)$$

$$= h(h+1) + (h-1) + (h-2)(n-1) + g(n-3)$$

$$\vdots$$

$$= \sum_{i=1}^{n} i(i+1) = \frac{h(n+1)(h+2)}{3} \quad \text{"closed form"}$$

Recap: Learning Objectives

- Understand how to read a recursive definition, e.g., compute selected values or objects produced by that definition.
- Know that a recursive definition, and an inductive proof, require both a base case and an inductive step/formula.
- Define the Fibonacci numbers.