

# CS 173: Discrete Structures, Fall 2009

## Solutions to review problems for second quiz

### 1. Big-O notation

For each of the following pairs of functions state whether  $f(n) = O(g(n))$  or  $f(n) = \Omega(g(n))$  or  $f(n) = \Theta(g(n))$

(a)  $f(n) = \lceil n \rceil^2$  and  $g(n) = \lfloor n \rfloor^2$ .

**Solution:**  $f(n) = \Theta(g(n))$ . The two functions are equal for integer inputs. For non-integer inputs, floor and ceiling can differ by at most 1. So the difference between the squares will be no more than  $(n^2 + 2n + 1) - n^2 = 2n + 1$ . As  $n$  gets large, this grows more slowly than  $n^2$ .

(b)  $f(n) = (\log_{10}(n))^2$  and  $g(n) = n$ . The log function grows more slowly than  $n$ , and the square of the log function grows even more slowly.

**Solution:**  $f(n) = O(g(n))$ .

(c)  $f(n) = n^{2^n}$  and  $g(n) = n^{n^2}$ , because  $2^n$  grows much faster than  $n^2$ .

**Solution:**  $f(n) = \Omega(g(n))$ .

(d)  $f(n) = n!$  and  $g(n) = n^n$ .

**Solution:**  $f(n) = O(g(n))$ . If you expand both functions into products, their first terms are the same, but then the later terms are much smaller for  $n!$ . This difference gets bigger as  $n$  gets larger.

(e)  $f(n) = 2^n + n$  and  $g(n) = 3^n$

**Solution:**  $f(n) = O(g(n))$ . You can ignore the  $n$  term in  $f$ . As  $n$  gets large, the difference between  $2^n$  and  $3^n$  gets larger, because you keep multiplying in more factors of  $\frac{3}{2}$ .

Determine whether each statement below is true or false.

(f) If  $f(n) = \Theta(g(n))$  and  $h(n) = \Theta(g(n))$  then  $f(n)h(n) = \Theta(g(n))$ .

**Solution:** False.  $f(n)h(n) = O(g(n)g(n)) \rightarrow \forall x > k, f(x)h(x) \leq cg(x)g(x)$ .  $cg(x)$  cannot be replaced by a constant.

(g) If  $f(n) = \Omega(g(n))$  and  $h(n) = \Omega(g(n))$  then  $f(n) + h(n) = \Omega(g(n))$ .

**Solution:** True. We have,  $f(x) \geq c_1g(x)$  for all  $x > k_1$  and  $h(x) \geq c_2g(x)$  for all  $x > k_2$ . Hence we have that  $f(x) + h(x) \geq c_1g(x) + c_2g(x) = (c_1 + c_2)g(x)$  for all  $x > \max(k_1, k_2)$ .

(h) If  $f(n) = O(g(n))$  then  $g(n) = \Omega(f(n))$ .

**Solution:** True. We saw this in lecture. (It's one way to define  $\Omega$ .)

(i) If  $f(n) = \Theta(g(n))$  then  $g(n) = \Theta(f(n))$

**Solution:** True.  $\Theta$  is a type of equality. If it's true in one direction, it's true in the other.

(j) If  $f(n) = \log_a(n)$  for  $a > 2$  then  $f(n) \neq \Theta(\log_2(n))$ .

**Solution:** False. Changing the base of  $\log_a n$  multiplies all the outputs by a number that doesn't depend on  $n$  (it only depends on the old and new bases).

When the right answer is given as  $f(n) = \Theta(g(n))$  above, you could also (optionally) say that  $f(n) = O(g(n))$  and  $f(n) = \Omega(g(n))$ .