

Lecture 12: Circular Convolution

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ECE 401: Signal and Image Analysis

- ① Review: DFT
- ② Computational Complexity: Fast Fourier Transform
- ③ Circular Convolution
- ④ Zero-Padding
- ⑤ Summary

Outline

- 1 Review: DFT
- 2 Computational Complexity: Fast Fourier Transform
- 3 Circular Convolution
- 4 Zero-Padding
- 5 Summary

Review: DFT

The DFT (discrete Fourier transform) of any signal is $X[k]$, given by

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi kn}{N}}$$
$$x[n] = \frac{1}{N} \sum_0^{N-1} X[k] e^{j \frac{2\pi kn}{N}}$$

Particular useful examples include:

$$f[n] = \delta[n] \leftrightarrow F[k] = 1$$
$$g[n] = \delta[((n - n_0))_N] \leftrightarrow G[k] = e^{-j \frac{2\pi kn_0}{N}}$$

Properties of the DTFT

Properties worth knowing include:

- ③ Periodicity: $X[k + N] = X[k]$
- ① Conjugate symmetric: $x[n] \text{ Real} \leftrightarrow X[k] = X^*[-k]$
- ② Linearity:

$$z[n] = ax[n] + by[n] \leftrightarrow Z[k] = aX[k] + bY[k]$$

- ③ Time Shift: $x[n - n_0] \leftrightarrow e^{-j\frac{2\pi kn_0}{N}} X[k]$, paying attention to the fact that $x[n]$ is periodic in time
- ④ Frequency Shift: $e^{j\frac{2\pi k_0 n}{N}} x[n] \leftrightarrow X[k - k_0]$, paying attention to the fact that $X[k]$ is periodic in frequency

Convolution

Convolution (finite impulse response filtering) is a generalization of weighted local averaging:

$$y[n] = h[n] * x[n] \equiv \sum_m x[m]h[n-m] = \sum_m x[n-m]h[m]$$

- If all samples of $h[n]$ are positive, then it's a weighted local averaging filter
- If the samples of $h[n]$ are positive for $n > 0$ and negative for $n < 0$ (or vice versa), then it's a weighted local differencing filter

Frequency Response

- **Tones in → Tones out**

$$x[n] = e^{j\omega n} \rightarrow y[n] = H(\omega)e^{j\omega n}$$

$$x[n] = \cos(\omega n) \rightarrow y[n] = |H(\omega)| \cos(\omega n + \angle H(\omega))$$

$$x[n] = A \cos(\omega n + \theta) \rightarrow y[n] = A|H(\omega)| \cos(\omega n + \theta + \angle H(\omega))$$

- where the **Frequency Response** is given by

$$H(\omega) = \sum_m h[m]e^{-j\omega m}$$

Putting it all together

If the input signal is

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi kn}{N}}$$

... then the output of $y[n] = h[n] * x[n]$ is ...

$$y[n] = \frac{1}{N} \sum_{k=0}^{N-1} H[k] X[k] e^{j \frac{2\pi kn}{N}}$$

$$H[k] \equiv H\left(\frac{2\pi k}{N}\right)$$

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Convolution is an $\mathcal{O}\{N^2\}$ Operation

Suppose we convolve an N -sample $h[n]$ with an N -sample $x[n]$. Consider carefully where the nonzero terms are:

$$y[n] = \sum_{m=n}^{N-1} h[m]x[n-m], \quad 0 \leq n \leq 2N-1$$

Let's do a computational complexity analysis:

- There are $2N = \mathcal{O}\{N\}$ output samples.
- The m^{th} output sample results from a sum of length $N - m = \mathcal{O}\{N\}$
- The total complexity is therefore $\mathcal{O}\{N^2\}$

Fast Fourier Transform (FFT)

- The regular DFT formula, $X[k] = \sum x[n]e^{-j\frac{2\pi kn}{N}}$, is also $\mathcal{O}\{N^2\}$ (N terms in the sum, N output coefficients).
- In 1965, Cooley and Tukey found a way to re-arrange the terms of the DFT so it can be computed in $\mathcal{O}\{N \log_2 N\}$ multiply-add operations. Their algorithm is called the “fast Fourier transform” (FFT).
- For example, suppose $N = 1024$, then
 - $N^2 = 1,048,576$ operations
 - $N \log_2 N = 10,240$ operations (a factor of 100 savings)

Can we use the FFT to speed up convolution?

Remember this formula:

$$y[n] = \frac{1}{N} \sum_{k=0}^{N-1} H[k]X[k]e^{j\frac{2\pi kn}{N}}$$

How about the following?

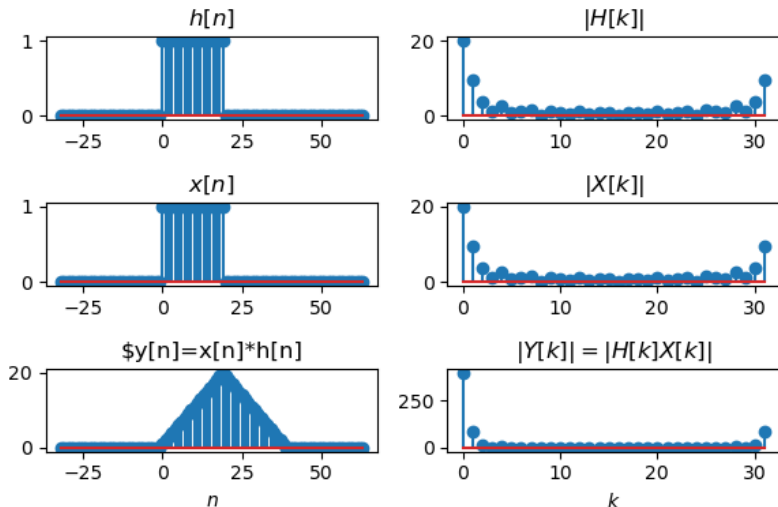
- 1 Take the FFT of $x[n]$ to find $X[k]$ ($\mathcal{O}\{N \log N\}$)
- 2 Take the FFT of $h[n]$ to find $H[k]$ ($\mathcal{O}\{N \log N\}$)
- 3 Multiply them to find $Y[k]$ ($\mathcal{O}\{N\}$)
- 4 Take the inverse FFT to find $y[n]$ ($\mathcal{O}\{N \log N\}$)

Is the total complexity $\mathcal{O}\{N \log N\}$, much faster than regular convolution? YES!! Does it compute the desired answer? NO!!!
Why not, and how can we fix it?

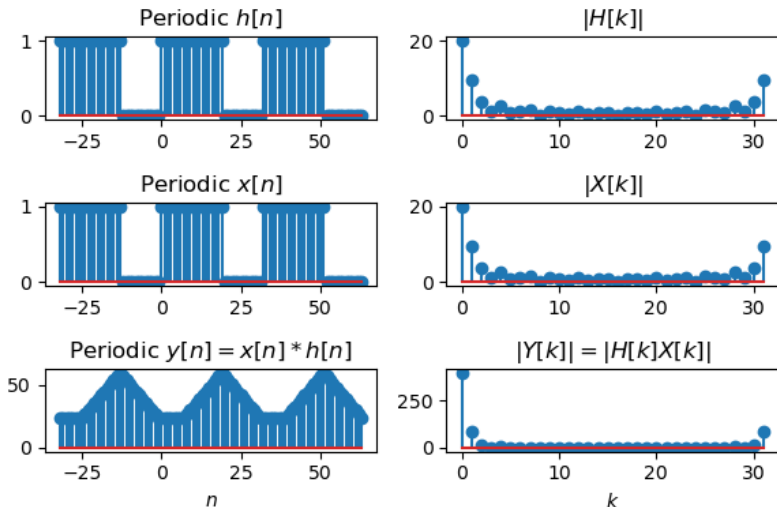
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Multiplying two DFTs: what we think we're doing



Multiplying two DFTs: what we're actually doing



Circular convolution

Suppose $Y[k] = H[k]X[k]$, then

$$\begin{aligned}y[n] &= \frac{1}{N} \sum_{k=0}^{N-1} H[k] X[k] e^{j \frac{2\pi kn}{N}} \\&= \frac{1}{N} \sum_{k=0}^{N-1} H[k] \left(\sum_{m=0}^{N-1} x[m] e^{-j \frac{2\pi km}{N}} \right) e^{j \frac{2\pi kn}{N}} \\&= \sum_{m=0}^{N-1} x[m] \left(\frac{1}{N} \sum_{k=0}^{N-1} H[k] e^{-j \frac{2\pi k(n-m)}{N}} \right) \\&= \sum_{m=0}^{N-1} x[m] h[((n-m))_N]\end{aligned}$$

... where the notation $((n-m))_N$ means “modulo N ,” i.e., we need to pay attention to the periodicity.

Circular convolution

Multiplying the DFT means **circular convolution** of the time-domain signals:

$$y[n] = h[n] \circledast x[n] \leftrightarrow Y[k] = H[k]X[k],$$

Circular convolution ($h[n] \circledast x[n]$) is defined like this:

$$h[n] \circledast x[n] = \sum_{m=0}^{N-1} x[m] h[(n-m)_N] = \sum_{m=0}^{N-1} h[m] x[(n-m)_N]$$

Linear convolution example

Circular convolution example

Quiz!

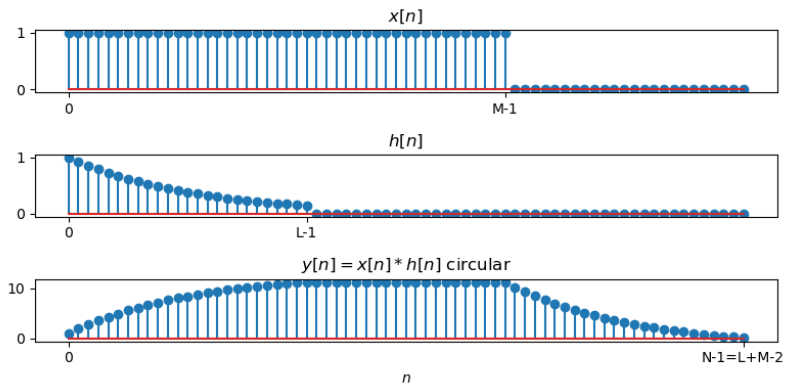
Go to the course webpage, try the quiz!

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How long is $h[n] * x[n]$?

If $x[n]$ is M samples long, and $h[n]$ is L samples long, then their linear convolution, $y[n] = x[n] * h[n]$, is $M + L - 1$ samples long.



Zero-padding turns circular convolution into linear convolution

How it works:

- $h[n]$ is length- L
- $x[n]$ is length- M
- As long as they are both zero-padded to length $N \geq L + M - 1$, then
- $y[n] = h[n] \circledast x[n]$ is the same as $h[n] * x[n]$.

Zero-padding turns circular convolution into linear convolution

Why it works: Either...

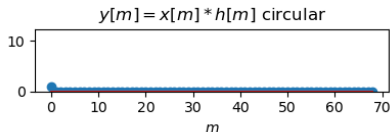
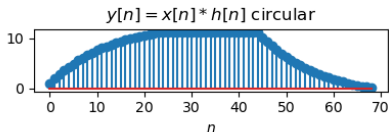
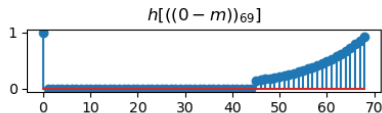
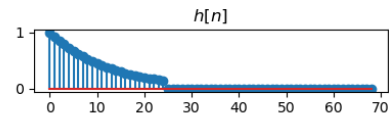
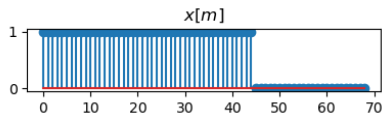
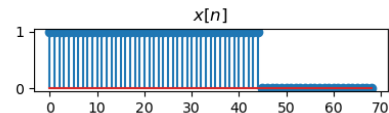
- $n - m$ is a positive number, between 0 and $N - 1$. Then $((n - m))_N = n - m$, and therefore

$$x[m]h[((n - m))_N] = x[m]h[n - m]$$

- $n - m$ is a negative number, between 0 and $-(L - 1)$. Then $((n - m))_N = N + n - m \geq N - (L - 1) > M - 1$, so

$$x[m]h[((n - m))_N] = 0$$

Case #2: $n - m$ is negative, so it wraps around, but N is long enough so that the wrapped part of $h[((n - m))_N]$ doesn't overlap with $x[m]$



Zero-padding turns circular convolution into linear convolution

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Circular convolution

Multiplying the DFTs, instead of implementing convolution directly, can save 100-fold computation or better. However, multiplying the DFT means **circular convolution** of the time-domain signals:

$$y[n] = h[n] \circledast x[n] \leftrightarrow Y[k] = H[k]X[k],$$

Circular convolution ($h[n] \circledast x[n]$) is defined like this:

$$h[n] \circledast x[n] = \sum_{m=0}^{N-1} x[m]h[((n-m))_N] = \sum_{m=0}^{N-1} h[m]x[((n-m))_N]$$

Circular convolution is the same as linear convolution if and only if $N \geq L + M - 1$.