

$$\Rightarrow \text{total time } O\left(\left(\frac{n}{w'}\right)^2\right) \leq \left(O\left(\frac{n^2}{(\log_2 n)^2}\right)\right)$$

Ex 3: 3SUM for ints

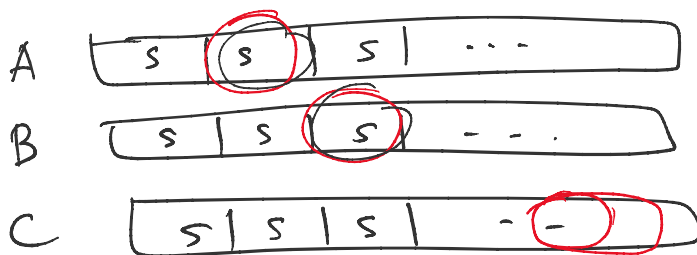
Baran, Demaine, Patrascu '05 $O\left(\frac{n^2}{\log^2 n} (\log \log n)^2\right)$.

idea - hashing again!

$h(x) = x \bmod p$. with rand $p \in (R/2, R)$
for a small R .

Warm-up with Convul-3SUM:

divide into $\frac{n}{s}$ blocks of s numbers



$\Rightarrow \left(\frac{n}{s}\right)^2$ subprobs of $O(s)$ numbers
after hashing, $O(s \log R)$ bits

Solve each subproblem by table lookup
of size $2^{O(s \log R)} \ll n$

choose $s = \frac{\delta \log n}{\log \log n}$ $R = \log^c n$

$$\Rightarrow O\left(\left(\frac{n}{s}\right)^2\right) = O\left(\frac{n^2}{\log^2 n} (\log \log n)^2\right)$$

for $a_i + b_j \neq c_{ij}$,
 $\Pr(a_i + b_j = c_{ij}) = h(c_{ij}) \lesssim \tilde{O}\left(\frac{1}{R}\right)$

for $a_i + b_j \neq c_{ij}$,

$$\Pr[h(a_i + b_j) = h(c_{ij})] \lesssim \tilde{O}\left(\frac{1}{R}\right)$$

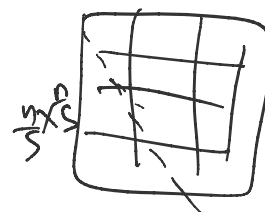
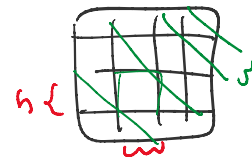
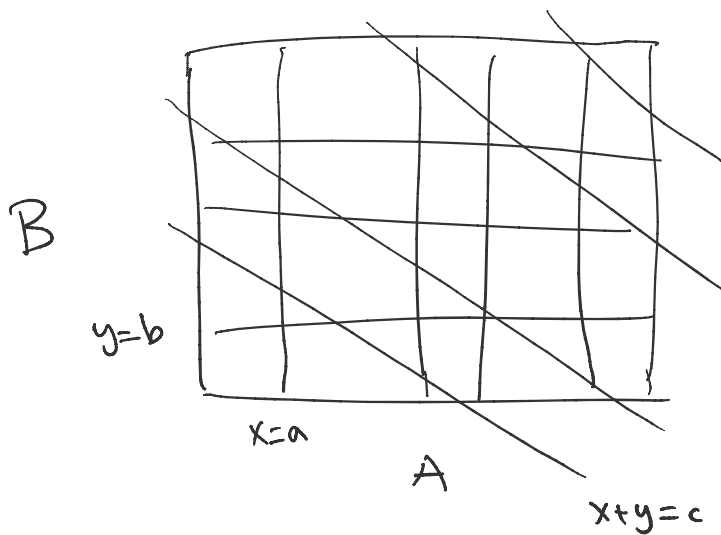
$$\Rightarrow \text{expected \# false positives} \lesssim \tilde{O}\left(\frac{n^2}{R}\right)$$

$$\Rightarrow \text{time for checking these is } \tilde{O}\left(\frac{n^2}{R} \cdot S\right) \ll O\left(\frac{n^2}{\log^{c-2} n}\right)$$

final bd $O\left(\frac{n^2}{\log^2 n} (\log \log n)^2\right)$.

for 3SUM; sort A, B, C

then divide into $\frac{n}{s}$ blocks of size s



$$O\left(\left(\frac{n}{s}\right)^2\right) \text{ cells} \Rightarrow O\left(\left(\frac{n}{s}\right)^2\right) \text{ subprobs of size } O(s)$$

rest is same.

What about reals?

APSP for Reals

first consider decision tree complexity

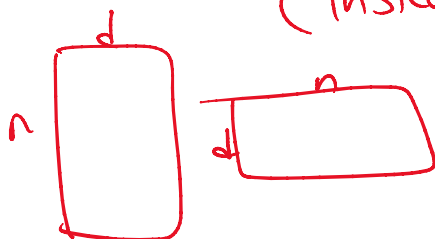
= # comparisons on real numbers

Fredman '75: APSP can be solved
in $O(n^{2.5})$ comps!

Lemma can compute $(\min, +)$ -MM of
 $n \times d$ & $d \times n$ matrix

using $\tilde{O}(d^2 n)$ comps

(instead of $O(dn^2)$)



Pf: key idea - $a_{ik} + b_{kj} \stackrel{?}{\leq} a_{ik'} + b_{kj}$
 $\iff \underline{a_{ik} - a_{ik'}} \stackrel{?}{\leq} b_{kj} - b_{kj}$
"Fredman's trick"

Sort $\{ a_{ik} - a_{ik'} : i \in [n], k, k' \in [d] \}$
 $\cup \{ b_{kj} - b_{kj} : j \in [n], k, k' \in [d] \}$

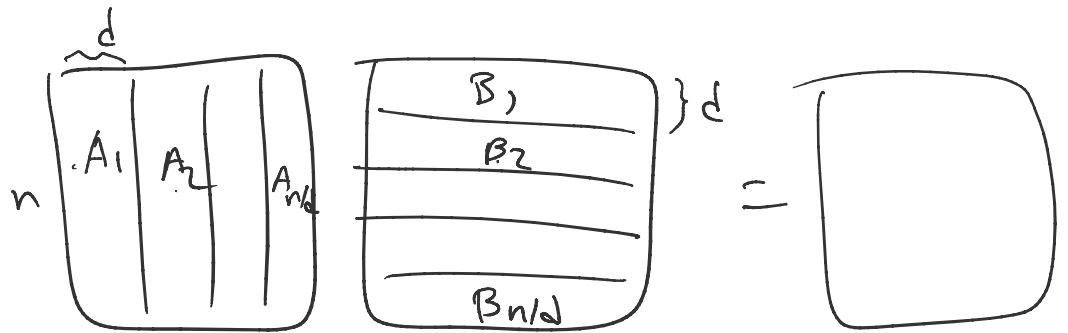
$\Rightarrow O(d^2 n)$ numbers
 $O(d^2 n \log(d^2 n))$ time

Afterwards, no more real comps needed! \square

Thm

can compute (min,t)-MM of 2 $n \times n$ ^{real} matrices
in $\tilde{O}(n^{2.5})$ comps

Pf: reduce to $\frac{n}{d}$ products of $n \times d$ & $d \times n$



$$\Rightarrow \tilde{O}\left(\frac{n}{d} \cdot d^2 n + \frac{n}{d} \cdot n^2\right)$$
$$= \tilde{O}\left(d n^2 + \frac{n^3}{d}\right)$$

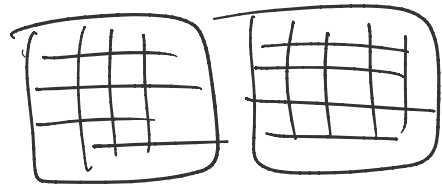
Choose $d = \sqrt{n}$.

\square

Improving Time:

reduce to $\left(\frac{n}{s}\right)^3$

(min,t)-MM of $s \times s$ matrices



build decision tree
of size/
time

$$2^{\tilde{O}(s^{5/2})} \ll n$$

choose $s \approx 8 \log^{2/5} n$

choose $s \approx 8 \log^{2/5} n$

$$\Rightarrow \sim O\left(\left(\frac{n}{s}\right)^3 \cdot s^{2.5}\right) = O\left(\frac{n^3}{\sqrt{s}}\right) \approx O\left(\frac{n^3}{\log^{1/5} n}\right)$$

3SUM for Reals

40 yrs later!

Grönlund-Pettre '14: 3SUM can be solved in $O(n^{1.5})$ comps

$a+b \leq a'+b'$

Warm-up with (min,+)-Convul or Convul3SUM:
↑
(BCDEHILPT'06)

divide $A = \overbrace{A_1 | A_2 | \dots}^{d \quad d} \quad \frac{n}{d} \text{ blocks}$
 $B = B_1 | B_2 | \dots$
 $C = C_1 | C_2 | \dots$

idea -

Fredman's trick

$$a_i + b_{k-i} \leq a_i + b_{k-i} \leftarrow$$

$$\Leftrightarrow a_i - a_i \leq b_{k-i} - b_{k-i} \leftarrow$$

$$\Leftrightarrow a_i - a_{i'} \stackrel{?}{\leq} b_{k-i'} - b_{k-i} \leftarrow$$

Preproc: Sort $\{ a_i - a_{i'} : i, i' \text{ in same block} \}$
 $\cup \{ b_j - b_{j'} : j, j' \text{ in same block} \}$

in $\tilde{O}\left(\frac{n}{d} \cdot d^2\right) = \tilde{O}(dn)$ time

Afterwards, can compute $\min(a_i + b_{k-i})$
 in each block with no real comp.

for (min,+)-convol:

\Rightarrow can compute c_k in $O\left(\frac{n}{d}\right)$ comps.

\Rightarrow total $\tilde{O}\left(dn + \frac{n^2}{d}\right)$

choose $d = \sqrt{n} \Rightarrow \tilde{O}(n^{3/2})$

for convol3SUM, to be cont'd ...