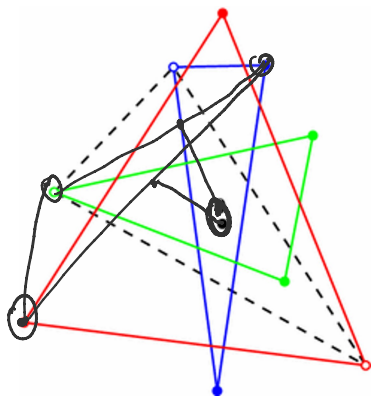


1. Colorful Caratheodory Theorem (CCT)



$(d+1)^{(d+1)}$ panchromatic simplices.

1. Is this CLS-hard?
OR is there a sub-exponential time algo?

2. For $d \geq 2$ we know there exists non-trivial soln. What is the complexity of finding say 2 solns?
what's the complexity of finding the # of solns?

2. Contraction Map (Banach Fixed Point)

Given: $f: [0,1]^n \rightarrow [0,1]^n$ given by an arithmetic circuit.
 $\lambda, \epsilon < 1$. p -norm
-m . $\|f(x) - f(y)\|_p \geq \lambda \|x - y\|_p$

Given: $\lambda, \epsilon < 1$. P -norm

Goal: either find $x, y \in [0, 1]^m$ s.t. $\|f(x) - f(y)\|_p > \lambda \|x - y\|_p$
 OR find $x \in [0, 1]^m$ s.t. $\|f(x) - x\|_p \leq \epsilon \left(\frac{1}{2^n}\right)$



$$\|f(x) - x^*\|_\infty \leq \lambda \|x - x^*\|_\infty$$

We know:

1. Unique fixed-point
2. For $\lambda = (1-\delta)$, it $\delta \leq \frac{1}{\text{poly}(m)}$ then easy (say $p=2$)

$$\|f^2(x) - f(x)\|_p \leq \lambda \|f(x) - x\|_p$$

$$\|f^3(x) - f^2(x)\|_p \leq \lambda \|f^2(x) - f(x)\|_p \leq \lambda^2 \|f(x) - x\|_p \leq \lambda^2$$

$$\|f^k(x) - f^{k-1}(x)\|_\infty \leq \lambda^k \|f(x) - x\|_p \leq \lambda^k \leq \epsilon$$

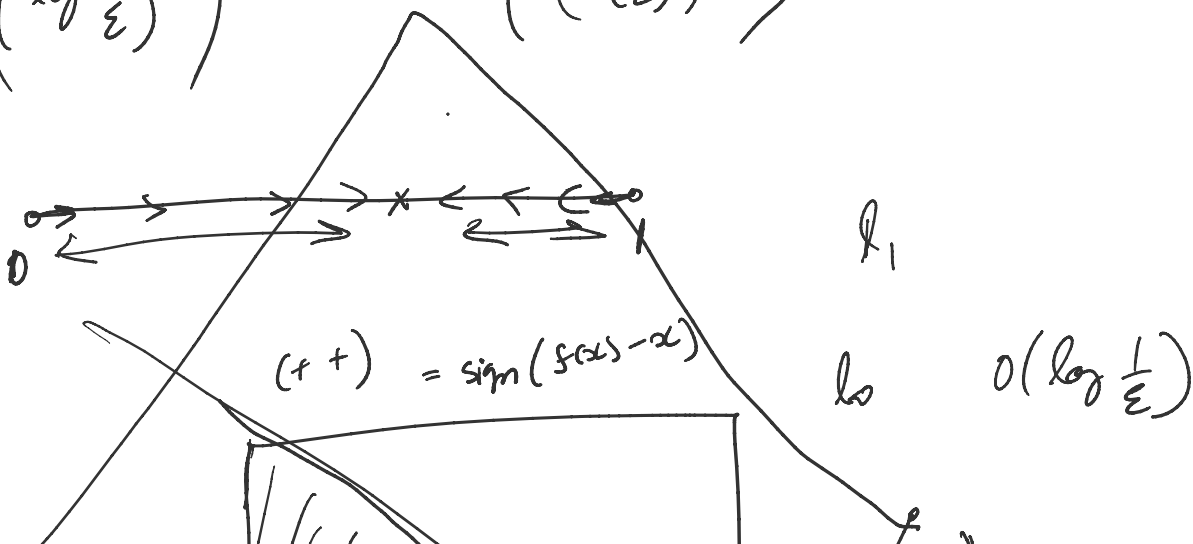
$$\Leftrightarrow (1-\delta)^k \leq \epsilon \Leftrightarrow \frac{1}{\epsilon} \leq (1-\delta)^k$$

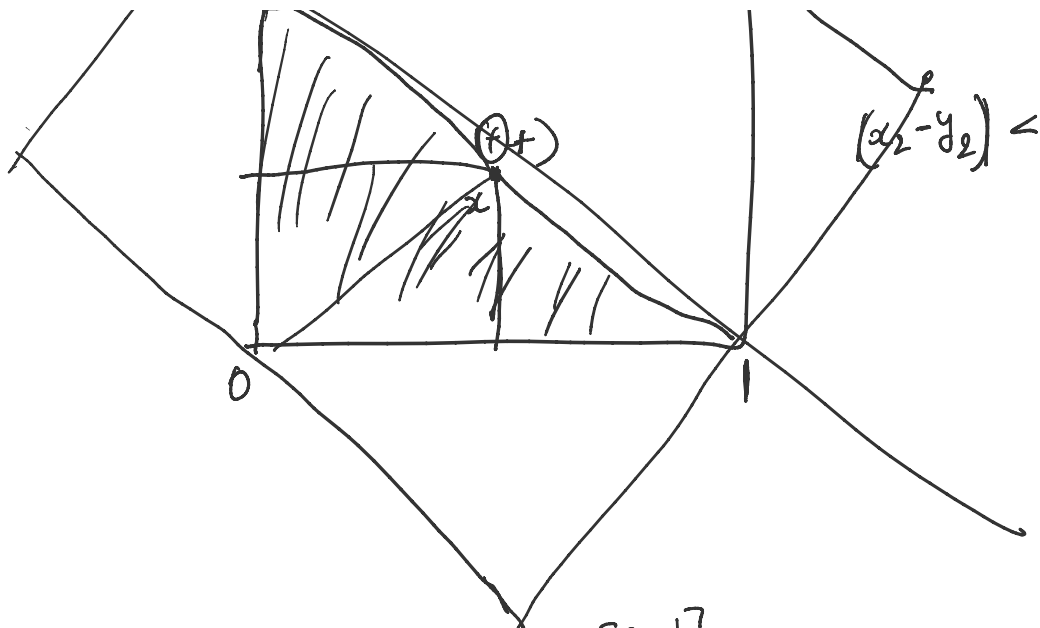
$$\Leftrightarrow \log \frac{1}{\epsilon} \leq k \log(1-\delta) \leq k \delta$$

$$\Leftrightarrow k \geq \left(\frac{1}{\delta} \log \frac{1}{\epsilon}\right)$$

3. $P=2$ poly-time algo.

$$h. \left(O\left(\log \frac{1}{\epsilon}\right)^m\right) \rightarrow O\left(\left(\log \left(\frac{1}{\epsilon}\right)\right)^{m/2}\right)$$





$$f: [0, 1]^n \rightarrow [0, 1]^n$$

$$g: [0, 1]^{n-1} \rightarrow [0, 1]^{n-1}$$

$$a \in [0, 1]$$

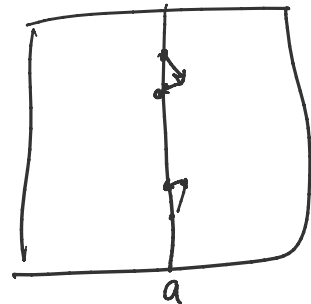
$$g(x) = f(a, x) \Big|_{\{2, \dots, n\}}$$

$$x, y \in [0, 1]^{n-1}$$

$$\|g(x) - g(y)\|_0 = \left\| \begin{matrix} f(a, x) \\ \{2, \dots, n\} \end{matrix} - \begin{matrix} f(a, y) \\ \{2, \dots, n\} \end{matrix} \right\|_0$$

$$\leq \|f(a, x) - f(a, y)\|_0$$

$$\leq \lambda \|x - y\|_0$$

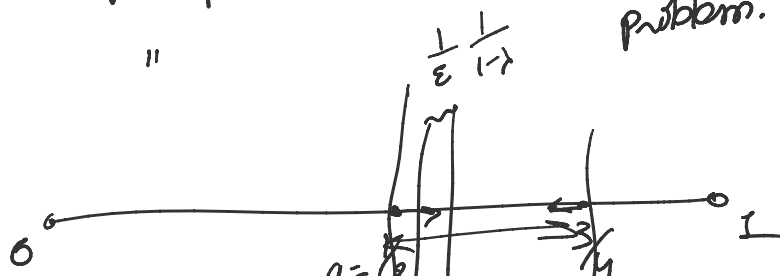


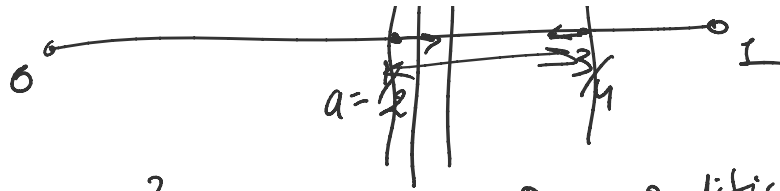
$$g(x^*) = x^* \quad x^* \in [0, 1]^{n-1}$$

$$f(z^*) = z^* \quad \|f(z^*) - z^*\|_p = \|f(x^*) - f(x^*)\|_p \leq \lambda \|x^* - x^*\|_p$$

$$f(a, x^*) \Big|_{\{2, \dots, n\}} = x^* \quad \text{if } f(a, x^*)_1 > a \text{ then}$$

1-D fixed point and oracle for solving $(n-1)$ dim problem.





Lemma: $\{1, \dots, m\}$ D_1, D_2, \dots, D_K partition
 \downarrow \downarrow \dots \downarrow
 $[0, 1]^{D_1}$ P_2 \dots P_K
 \downarrow \downarrow \dots \downarrow
 P_1 P_2 \dots P_K
 $P = P_1 \times P_2 \times \dots \times P_K$

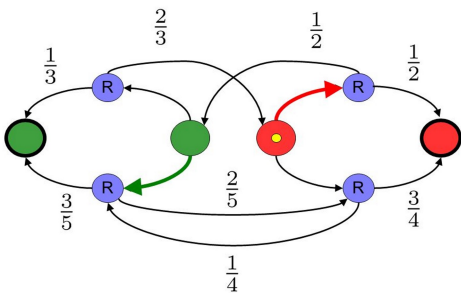
$1 \leq i \leq n/2$ $D_i = \{2i-1, 2i\} \rightarrow \log(1/\epsilon)$
 $P = \prod_{i=1}^{n/2} \text{time}(P_i) = \prod_{i=1}^{n/2} \log(1/\epsilon) = \left(\log \frac{1}{\epsilon}\right)^{n/2}$

Q: 1. Sub-exp? $\left(\log \frac{1}{\epsilon}\right)^{n^{1-\delta}}$ for $\delta > 0$, $\left(\log \frac{1}{\epsilon}\right)^{\frac{n}{\log n}}$?

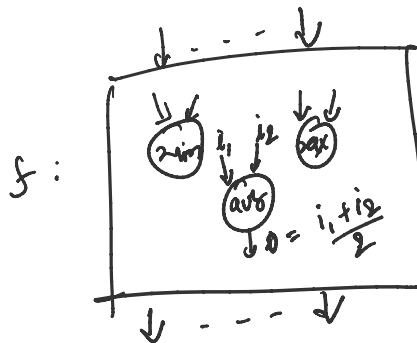
2. Query lowerbound: we know it's $n=3$ then req. $\left(\log \frac{1}{\epsilon}\right)^2$ queries.
 $\left(\log \frac{1}{\epsilon}\right)^{\frac{n}{\log n}}$ $\left(\log \frac{1}{\epsilon}\right)^{\lfloor \frac{n}{2} \rfloor}$?

3. Simple Stochastic Games

A simple Simple Stochastic Game



3 Randomized' Sub-exp. time algo.

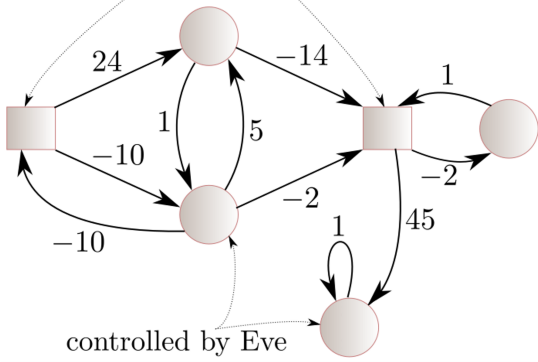




4. Mean Payoff Games & Parity Games

MPG (randomized sub-exp.)

controlled by Adam



controlled by Eve

quasi-poly

(STOC'17 Best Paper)

Parity

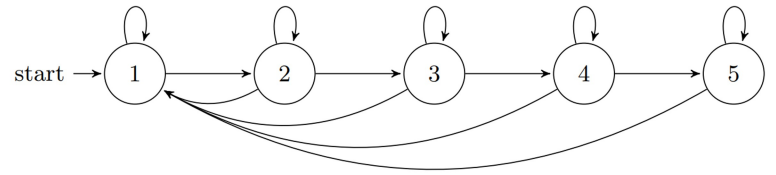


Figure 1.