CS 598 TH1

(Recent) Advances in Theory CS

Instructor: Ruta Mehta TA: None

Theory CS (TCS) How to compute? And fast?



Too vast to even touch every topic.

Our goal:

Together learn about some interesting new results within TCS and attempt to solve open problems. Fuzzy! Pre-req: CS 473. Advanced theory courses like approximation/randomized algorithms would help further.

Course Format: Seminar course

- □ I/you/guest lectures on an interesting topic, and discuss open problems around it.
- \Box An open problem session.
- □Repeat.

Project:

□ Pick an open problem and try to solve it by yourself. You have the entire semester to try.

What is expected from you

- □ Active participation in the class
- □ Present a paper of your choice, and discuss open problems around it.
- □ Work on an open problem, and present findings at the end of the semester.
- **Grading:** (all good as far as you are sincere and put in hard work)
 - \Box Class participation (5%)
 - □ Paper presentation (50%)
 - \Box Project + presentation (45%)

Administrativia

Webpage:

https://courses.engr.illinois.edu/cs598th1/sp2 023

Slack: CS598 TH1 – Advances in TCS (registered students should have received an invite to join)

 \Box Others, please sign up through the link on the webpage.

Issue about the classroom!

Check webpage/slack at least twice a week for the updates.

Tentative Topics

- TFNP classes and (open) problems therein
 Query, communication, crypto complexity
- Discrete optimization through continuous opt.
- Fast algorithms to solve LP
- Beyond worst-case
- Expander decomposition and apps. to graph algo.
- Massively Parallel Algorithms
- Discrepancy theory and Quantum connection?
- ML: differential privacy, fairness, learning augmented algorithms, GANs for counter examples

Topic 1: Total Function NP (TFNP) zoo

- What is TFNP? And why?
- Classes within TFNP: Intuition, formal description, complete problems w/ interesting existence theorems.
- Relation between the classes.

Computable? (Decidable?)

Halting problem: Given a program and its input, can one/computer decide if the program will halt?

Decidable:

NO! Un-decidable!!!





Ja sul? NP YES onsuens are poly-time verifiable. 3-SAI Ø on nvariables zi Eai est, F3 \$(21,..., xn) is True.

co-NP NO is easy 3-5AT Ø is of tantology





Integer tactorization: Given the integers. n, k s.t. n>k, deck is I a prime tactor Bon loss Harrock.

PLS (Polynomial Local Search)



Intuition:
$$f: [n] \rightarrow [n]$$
 then
 $eithen \exists x \in [n] \quad s.t. \quad f(x) = 0$
or $\exists x, y \quad s.t. \quad f(x) = f(x)$
Formul:
 $fo, |z^n \rightarrow [c] \rightarrow \{o, |z^n \quad eithen \quad fi \rightarrow d \neq x \quad s.t. \quad c(x) = 0^n$
 $or \quad x, y \quad s.t. \quad c(x) = 0^n$
 $or \quad x, y \quad s.t. \quad c(x) = (y)$

- Equal-sum:
Given nintegas the,
$$a_{1},...,a_{n}$$
,
 $\sum_{i=1}^{n} a_{i} < 2^{n-1}$, fixed $S \neq T \leq \tilde{z}_{1}, \tilde{z}_{1}$
 $i < \tilde{z}_{1}$
 $s.t. \quad \tilde{z}_{1} < \tilde{z}_{1}$
 $i < \tilde{z}_{1} < \tilde{z}_{1}$

Mikowski: Given ratix Amen s.t. [det (A)] < I, find the int. linear combination B A's hours s.t. its low notion < I. - Richlet: Given n rational no. a....and int. N. Find integes 9, Pro..., Pon s.b. Vi, [ai-Pi] < 1/2 N, Find integes 9, Pro..., Pon s.b. Vi, [ai-Pi] < 1/2 N = 1 < 9 < Nⁿ 9 / 9 N

