

# The Memory Performance of DSS Commercial Workloads in Shared- Memory Multiprocessors

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# Motivation

- Database applications: common workload;
- Memory performance studies: very few;
- Memory resident: cheaper RAM;

What is the shared-memory multiprocessor memory hierarchy behavior for complex database queries ?

# Outline

- Background information;
- Memory access pattern;
- Experimental evaluation;
- Conclusions;
- Future work.

# Basic Operations

DeptLoc

Dept	Location
1	Illinois
2	Texas

Personnel

Name	Dept
Helen	1
John	2
Julie	2
Peter	1

## ■ Select:

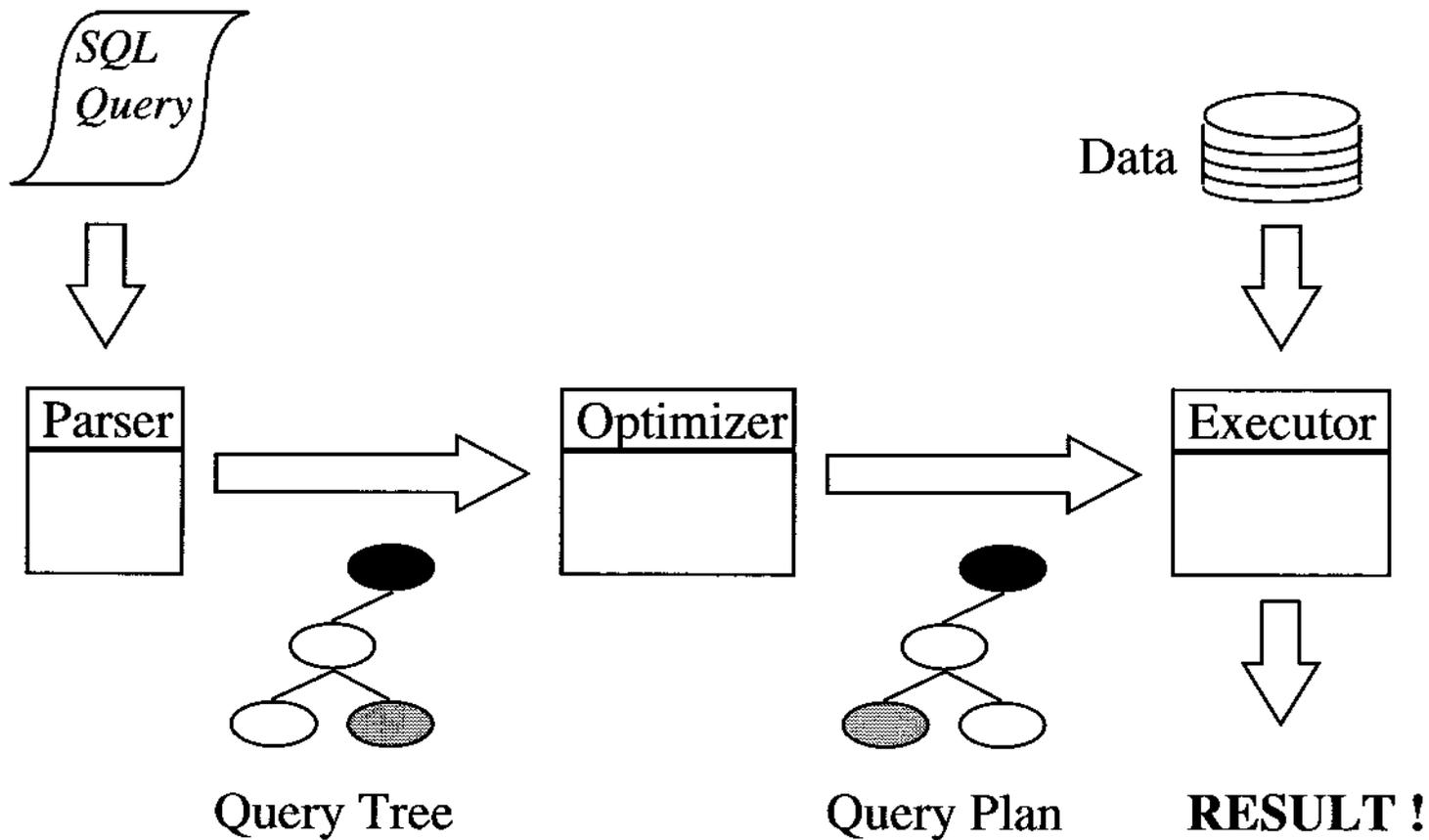
- Example: *Who works in Dept 1 ?*
- Implementation:
  - » Index Scan

## ■ Join:

- Example: *Who works in Texas ?*
- Implementation:

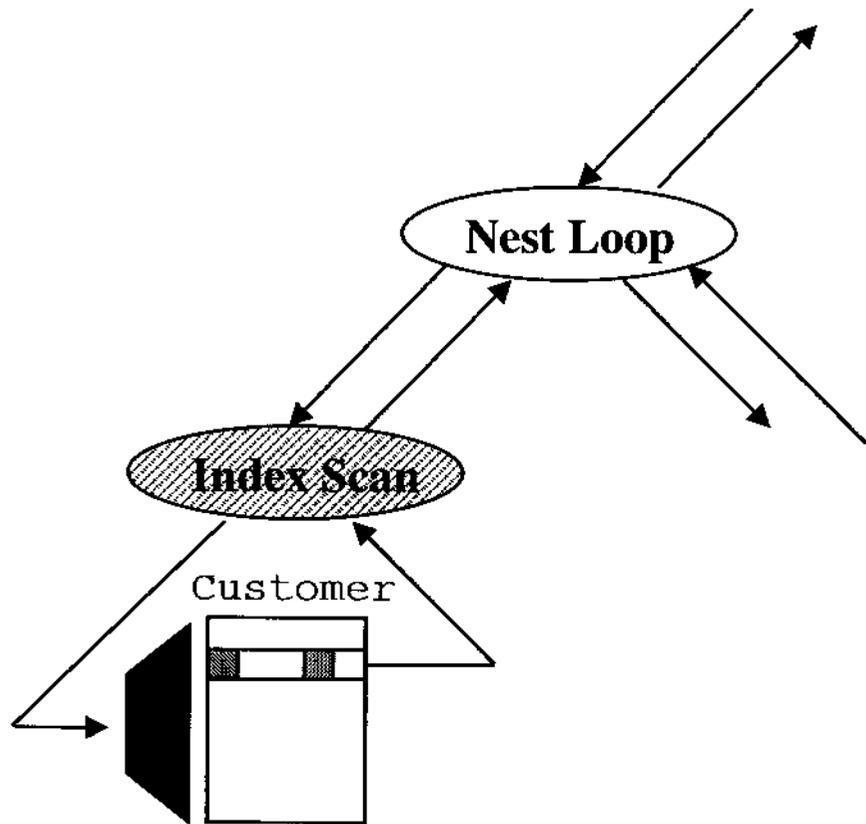
## ■ Other: Sort, Group, Aggregate;

# Query Execution



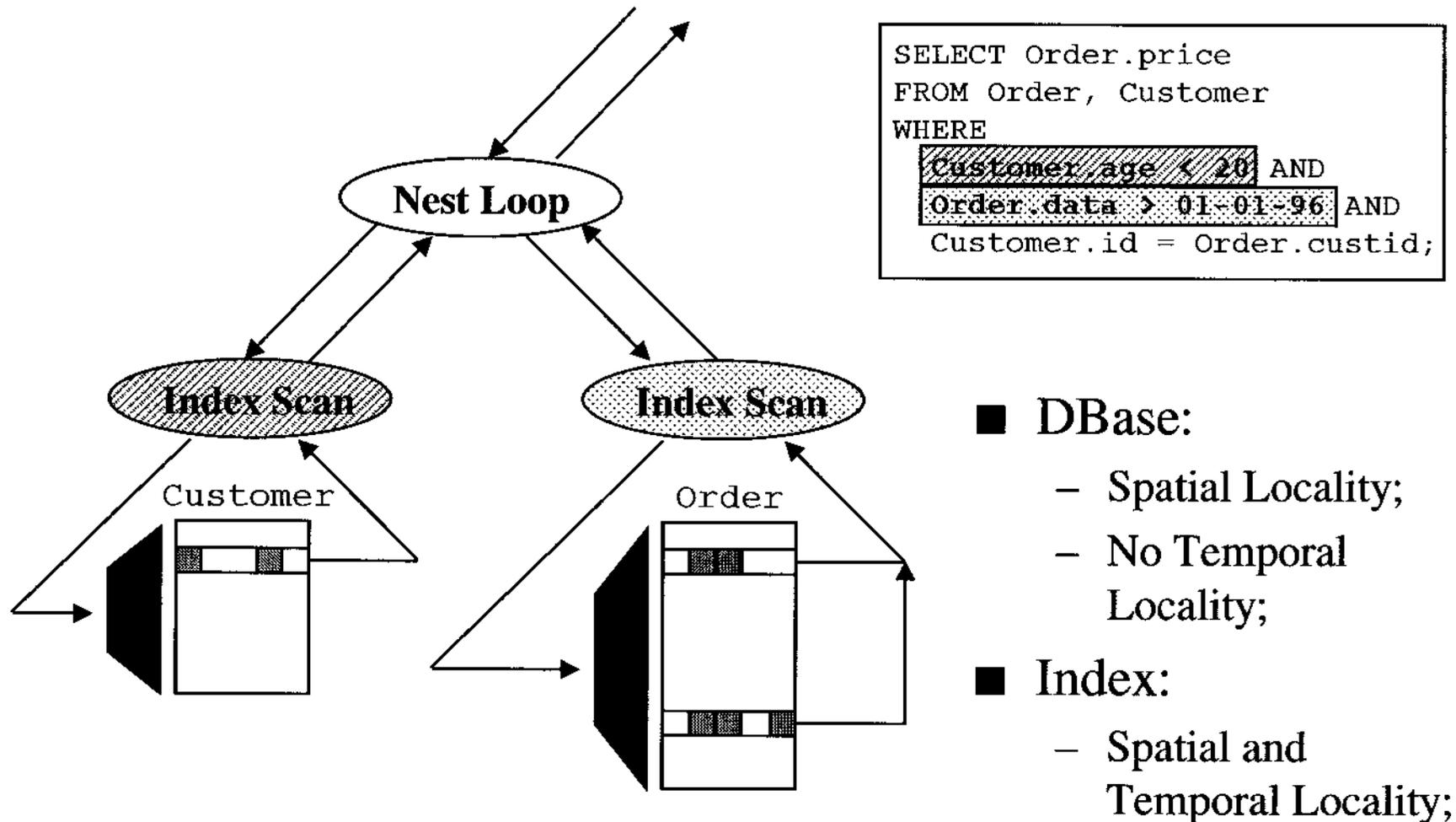


# Memory Access Pattern

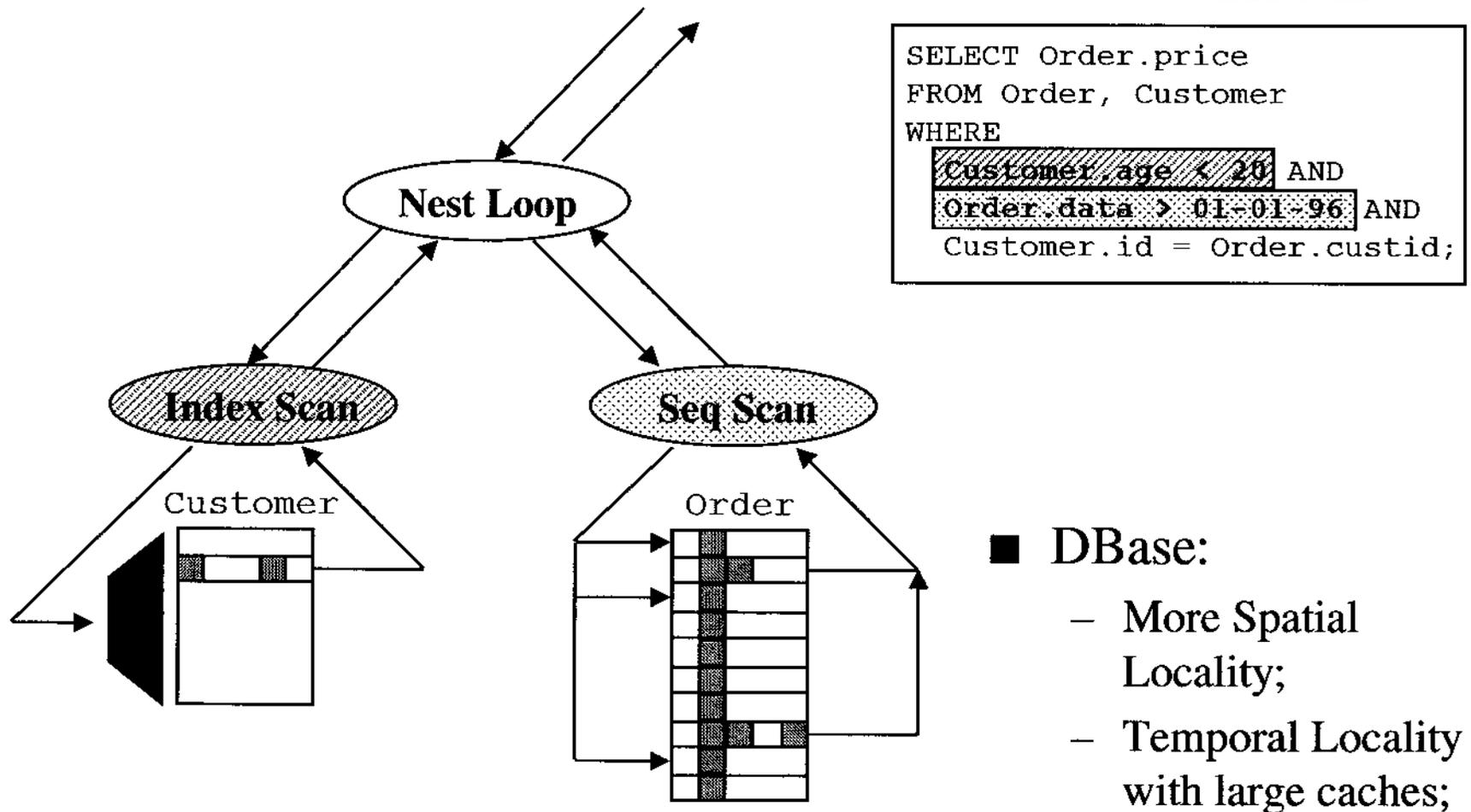


```
SELECT Order.price  
FROM Order, Customer  
WHERE  
Customer.age < 20 AND  
Order.data > 01-01-96 AND  
Customer.id = Order.custid;
```

# Memory Access Pattern



# Memory Access Pattern



# Workload

## ■ TPC-D:

- Decision Support System Benchmark;
- 17 read-only and 2 update queries;

## ■ Representative queries:

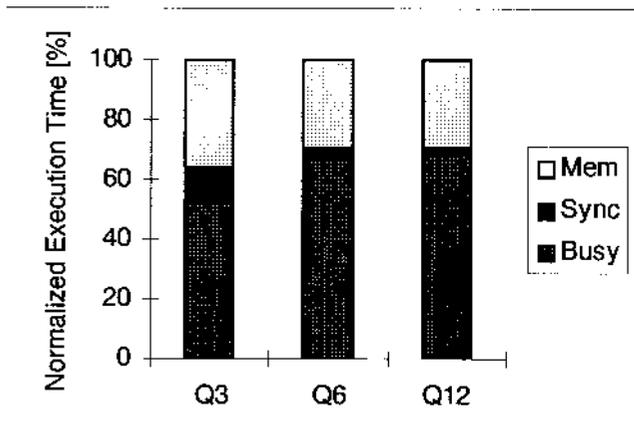
- Index Scan Select: Q3;
- Sequential Scan Select: Q6;
- Index and Sequential Scan Select: Q12;

# Evaluation

## ■ Setup:

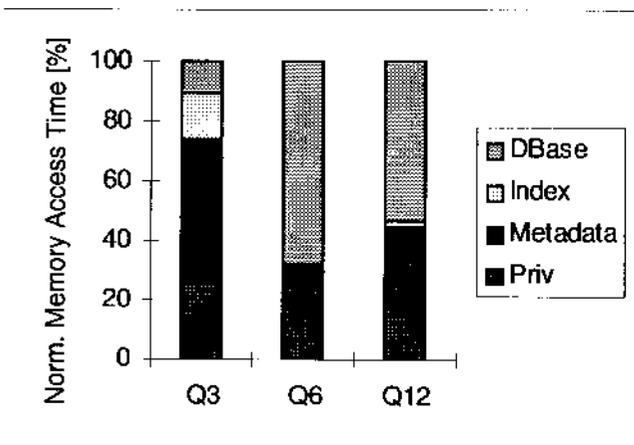
- System: 4-node directory-based cache-coherent NUMA shared-memory multiprocessor;
- Cache: 32-byte line 4-Kbyte primary and 64-byte line 128-Kbyte secondary;
- Simulator: Mint + Zibra (tracing shared and private data references);
- DBMS: Postgres95;
- Workload: TPC-D Q3, Q6, and Q12;
- Data set: data size aprox. 20 Mbyte.

# Overall Memory Behavior



## ■ Execution Time:

- Large *Busy*: 50-70%
- Significant *Mem*: 30-35%

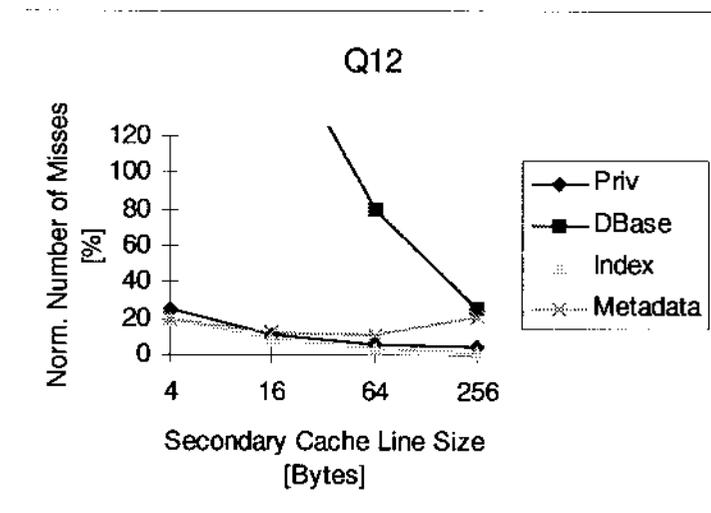
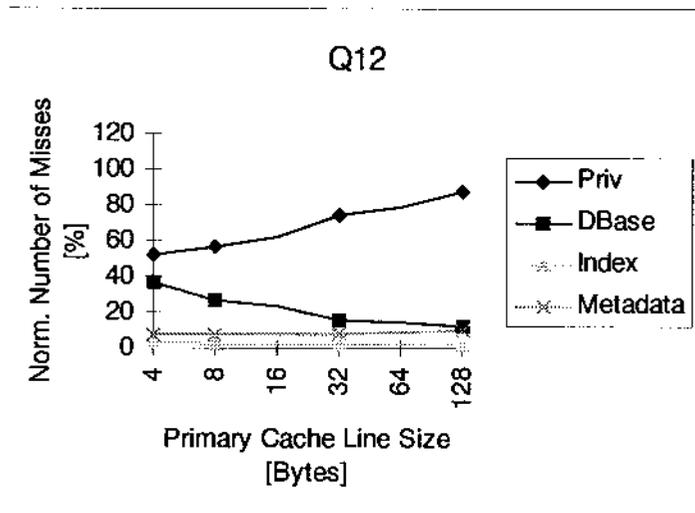


## ■ Memory Access Time:

- Q3: *Metadata+Index* dominate;
- Q6 and Q12: *DBase* dominates;

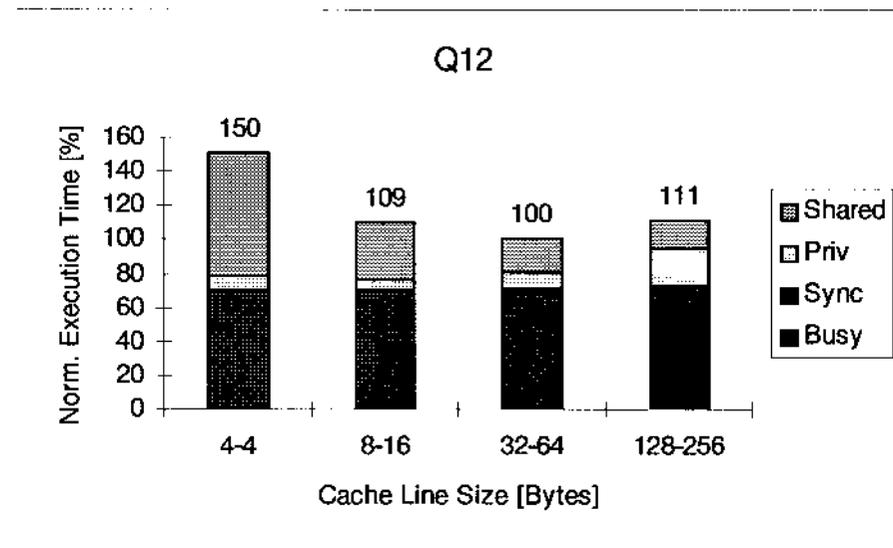
*Index and Sequential queries*

# Spatial Locality: Good



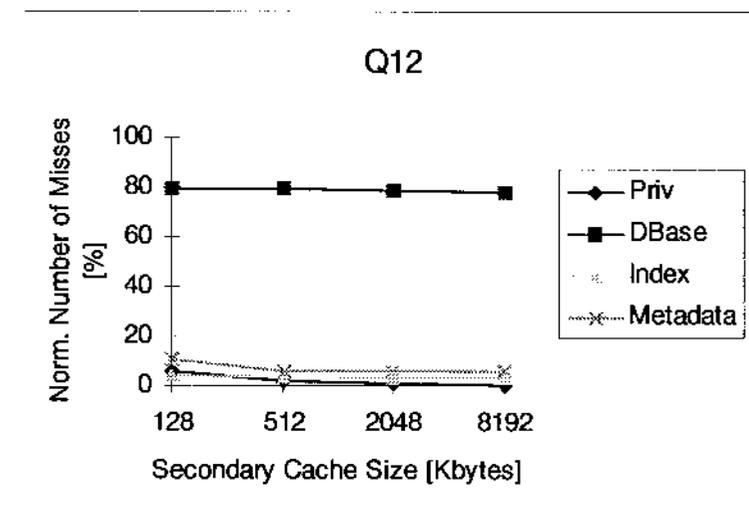
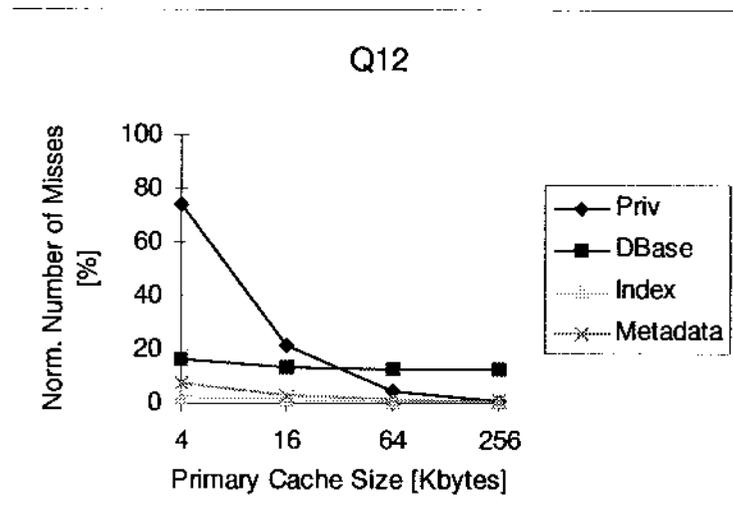
- *DBase* misses decrease with larger cache line size;

# Spatial Locality: Good



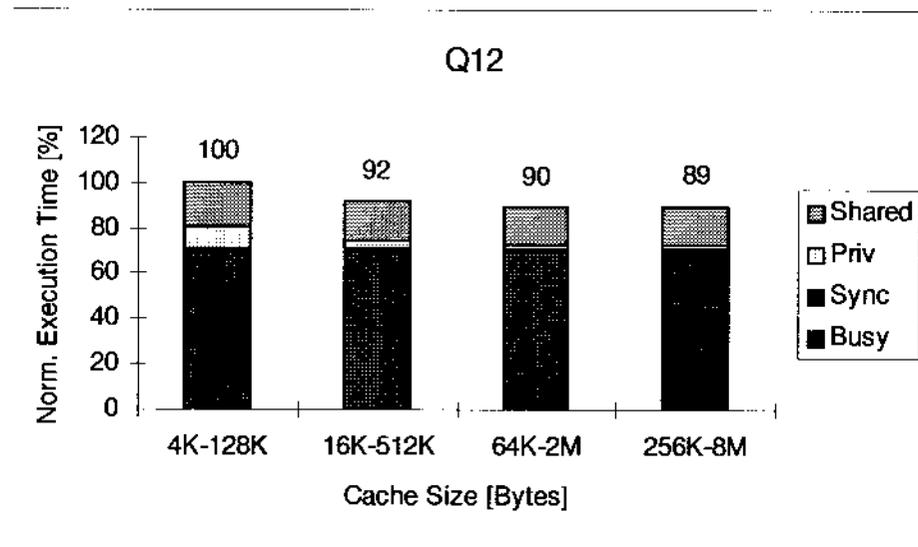
- With larger cache line size:
  - *Shared* decreases;
  - *Priv* increases;
- Best results: relatively long cache lines (64 bytes);

# Intra-Query Temporal Locality: Bad



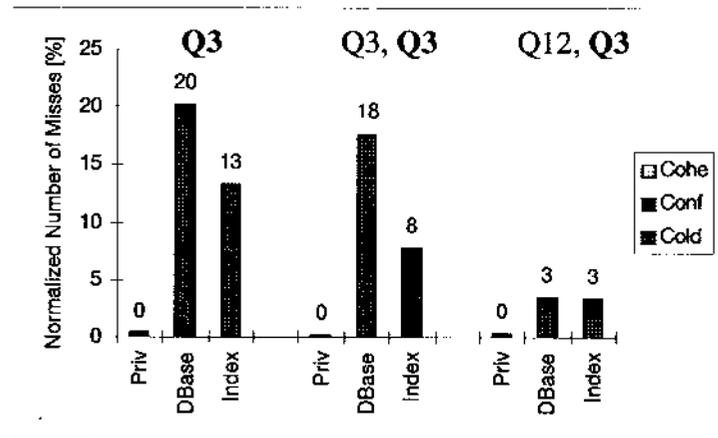
- *DBase* misses: maintain constant with different cache size;

# Intra-Query Temporal Locality: Bad



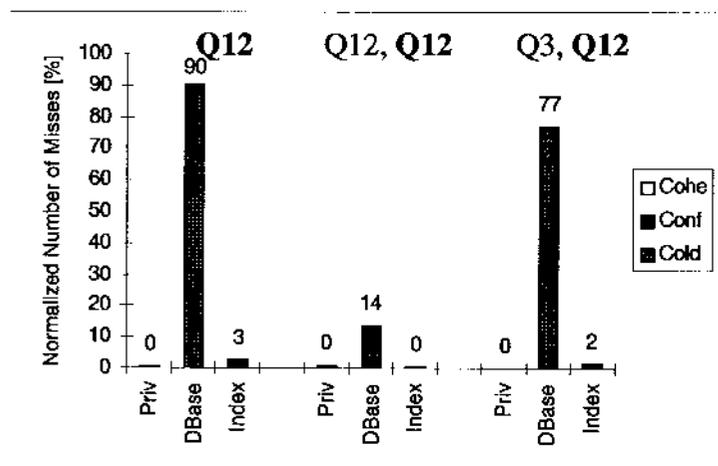
- With larger cache size:
  - *Shared* maintains the same;
  - *Priv* decreases;
- No change: no temporal locality;

# Inter-Query Temporal Locality: Maybe



## ■ Q3:

- Q3, Q3: reuse of *Index*;
- Q12, Q3: reuse *Index+DBase*;



## ■ Q12:

- Q12, Q12: large reuse of *DBase*;
- Q3, Q12: very little reuse;

# Conclusions

- Queries are different:
  - *Index* and *Sequential* queries;
  
- Caches can work:
  - Spatial locality: index and sequential queries;
  - Temporal locality: mostly across sequential queries;
  
- Simple optimizations:
  - Data prefetching for sequential queries;

## Future Work

- Further memory analysis:
  - More complex queries;
  
- More complex optimizations:
  - Achieve better resource usage.

More information:

<http://www.csrd.uiuc.edu/iacoma>