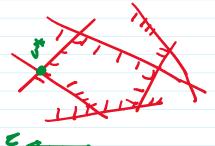
## 1 in 2D

Given n halfspaces in R<sup>2</sup> & vector &, find a point in infersection extreme along z.

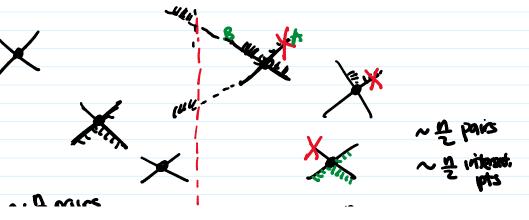


Prune - and - Search (Megiddo 83/ Dyer 83)

led. Lt = lines for upper halfplanes L = lines for lower halfplanes v\* = optimal sorn

repeat {

- 1. pair up lines in Lt, lines in I
- 2. Intersed each pair
- 3. Compute median x-coord xm of these n/z witersection Pts



~ # pairs ~ # pairs

remove line w. larger slope if from it

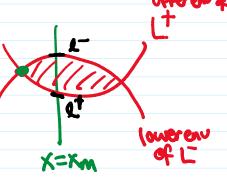
3 else similar

$$\Rightarrow T(n) \leq T\left(\frac{3}{4}n\right) + O(n)$$

How to decide in line 4?

Say objective is to min x

4.1. find highest the  $l \in L^{T}$  lowest the  $l \in L^{T}$  at  $x = x_{m}$ 



4.2. If I above  $l^{+}$  at  $x = x_{m}$  or slope of  $l < slope of <math>l^{+}$ 

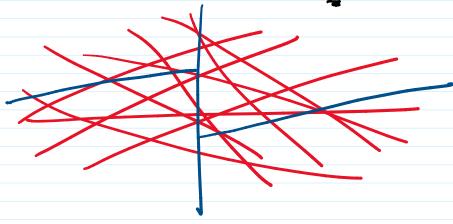
then left else right



## Extension to 3D (Megiddo'84)

Cutting Lemma Given n lines in  $\mathbb{R}^2$ , can divide  $\mathbb{R}^2$  into 4 regions

S.t. each region intersects  $\leq \frac{2}{8}n$  lines



## The algn:

## repeat (

- 1. pair up planes in Lt & planes in L
- 2. intersed each pair
- 3. apply cutting Lemma to xy-proj of these n/2 likes
- 4. decide which region ut is in
- 5. for each pair whose intersection (mè aussis the region remove one plane from pair





How to decide:

solve 2D LP ...

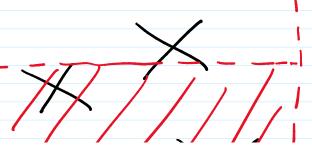
O(n) time



Pf of Cutting Lemma;

- 1. medran slope m
- 2. Pair lines of slope < m w. lines of slope > m
- 3. Intersect each par

5. on left side: when sheem the thru pts to the right



m=0

Megiddo  $O(2^{0(2^{k})}n)$ Dyer  $O(2^{0(4^{k})}n)$ Current best deterministic  $O(2^{0(4696)}n)$