

CS 498 TC: Computational Geometry

<http://courses.grainger.illinois.edu/cs498tcu>

Course Work:

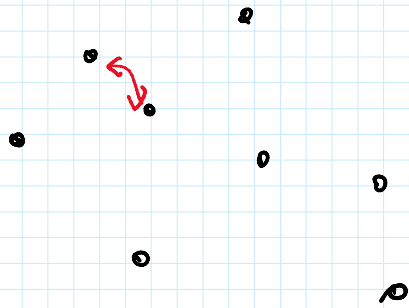
	TCU/TC3	TC4
5 HWs (groups ≤ 3)	40%	35%
Midterm (Oct 7 Mon 7p-9p)	20%	15%
Final	40%	35%
(Presentation for TC4 grad.)		15%

Prerequisite: CS374 or equiv.

What is CG?

algs for geom problems

e.g. closest pair



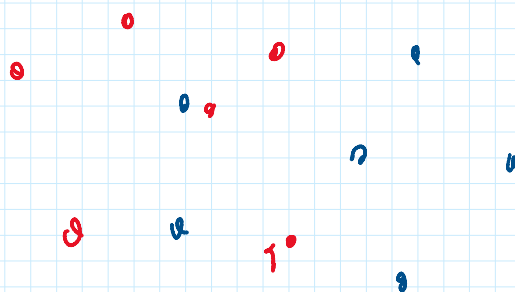
1D: $O(n \log n)$ time

2D: $O(n \log n)$ time

3D & beyond:
 $O(n \log n)$



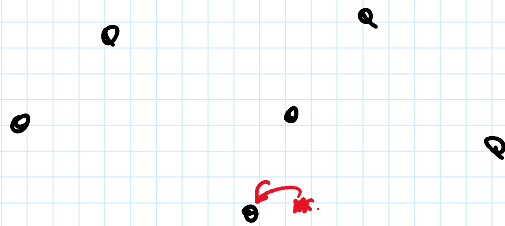
red-blue closest pair



2D: $O(n \log n)$

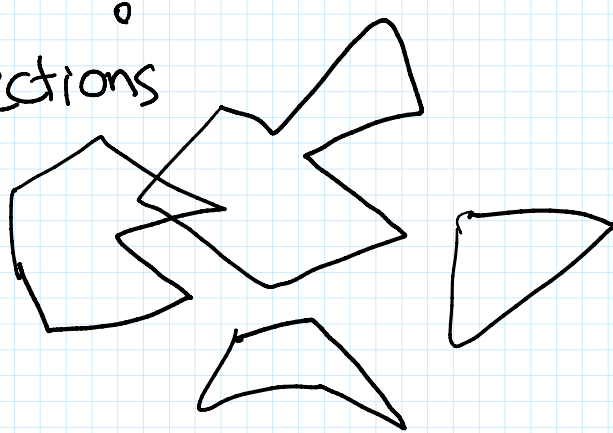
3D: $O(n^{4/3})$

nearest neighbor search



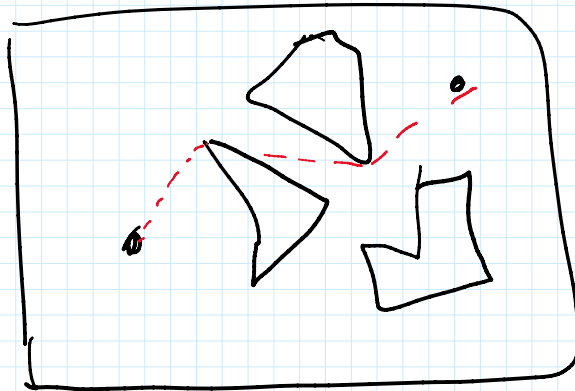
2D: $O(n)$ space
 $O(\log n)$ query time

intersections

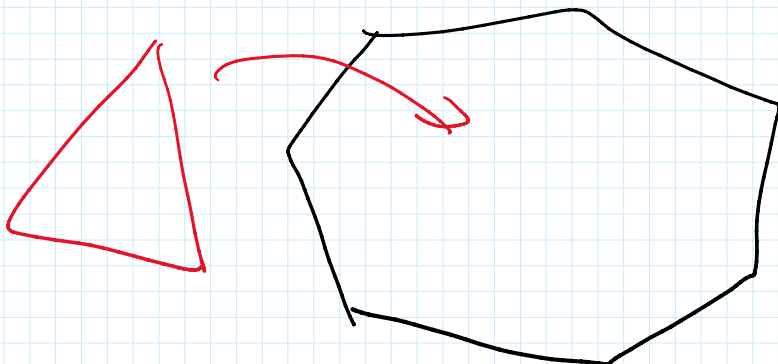


$O(n \log n)$

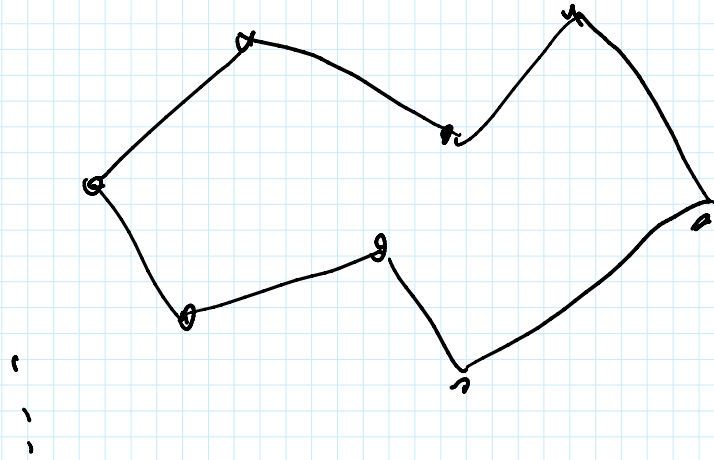
paths



polygon containment



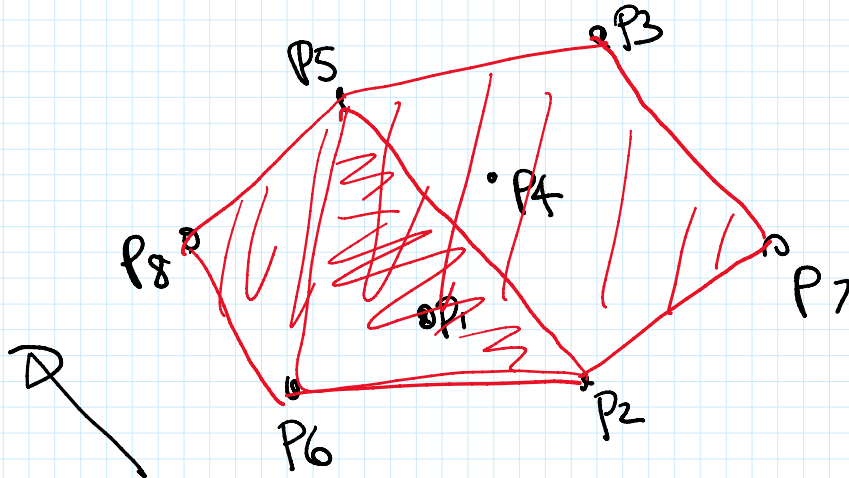
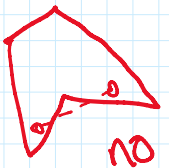
TSP



Convex Hull in 2D

$$P_i = (x_i, y_i)$$

Problem Given n points $P = \{P_1, \dots, P_n\} \subset \mathbb{R}^2$,
compute $CH(P) =$ Smallest convex set containing P \uparrow

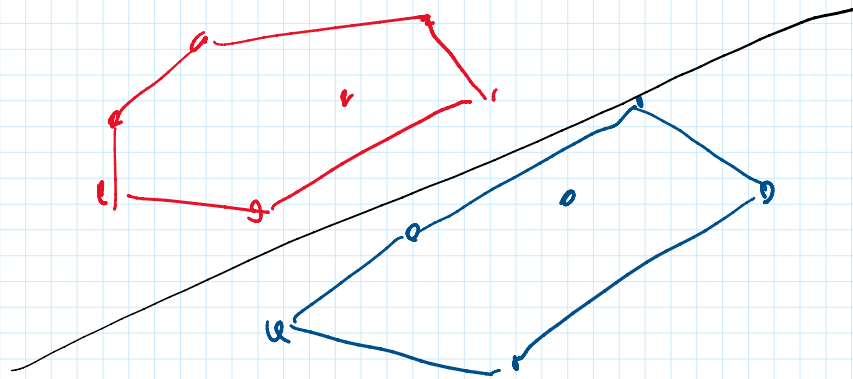


convex polygon
(all angles $\leq 180^\circ$)

Output: $P_8 P_6 P_2 P_7 P_3 P_5$ (in ccw order)

appl - find "shape" of point cloud

- bounding volume
- extreme pts, farthest pair
- line separation



Math Properties

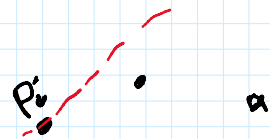
$$CH(P) = \bigcap_{\text{convex } S \supseteq P} S = \bigcap_{\text{halfplane } h \supseteq P} h$$

$$= \left\{ \text{all convex combinations of } P_1, \dots, P_n \right\}$$
$$\sum_{i=1}^n \alpha_i p_i \quad \text{for some } \alpha_i \geq 0 \quad \sum_{i=1}^n \alpha_i = 1.$$

$$= \bigcup_{i,j,k} \left\{ \text{all convex combs of } P_i, P_j, P_k \right\}$$

p_i is a vertex of $CH(P)$

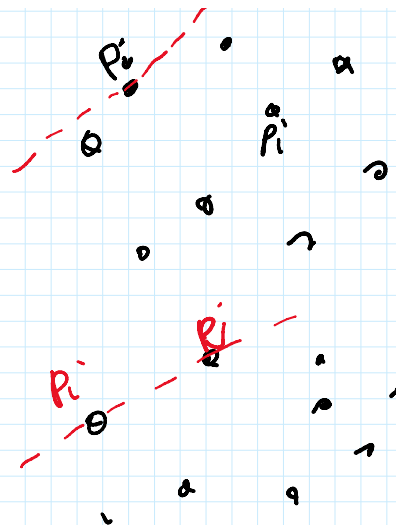
iff \exists line thru p_i s.t.
 \mathbb{D} lies on one side



iff \exists line thru p_i s.t.
 P lies on one side

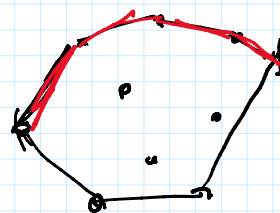
$p_i p_j$ is an edge of $CH(P)$

iff P lies on ~~one~~ side
of line $\overleftrightarrow{p_i p_j}$



Note - suffice to compute
upper hull (UH)

Alg'm O: Brute force



for $i = 1$ to n do

for $j = 1$ to n do {

flag = true

for $k = 1$ to n do $(k \neq i, j)$

if p_k above $\overleftrightarrow{p_i p_j}$ then flag = false

if flag then output $p_i p_j$

}

$O(n^3)$ time

Implementation issues:

(i) Primitive ops:
how to test p_k above $\overleftrightarrow{p_i p_j}$?

(ii) degeneracies

(ii) precision issues

: