## CS 498CG: Discrete & Computational Geometry, Spring 2023

Version: 3.0

Instructions: As in previous homeworks.

- **15** (100 PTS.) Luney lune. Do problem 9.14 in [dBCvKO08, page 217].
- **16** (100 PTS.) Report the points. Do problem 10.12 in [dBCvKO08, page 241].
- **17** (100 PTS.) Heavy time.

Let P be a set of n points in the plane in general position. A circumcircle C(p,q,r) of points  $p,q,r \in P$  is k-heavy if it contains k points of P in its interior. Let  $T_{\leq k}$  be the set of all circumcircles of points of P of weight at most k.

17.A. (30 PTS.)Let  $C = C(p, q, r) \in T_{\leq k}$  be a given circumcircle. Let R be a random sample from P of size n/k (without repetitions). Consider the event that the triangle  $\Delta = \Delta pqr$  is in the Delaunay triangulation of R (denoted by  $\mathcal{D}(R)$ ), and let

$$\alpha(C) = \mathbb{P}[\triangle pqr \in \mathcal{D}(R)]$$

be this probability. Prove that  $\alpha(C) \geq c/k^3$ , for some absolute constant c > 0.

- **17.B.** (30 PTS.) Prove that  $\sum_{C=C(p,q,r)\in T_{\leq k}} \alpha(C) = O(n/k).$
- **17.C.** (40 PTS.) Combining the above two parts, prove that  $|T_{\leq k}| = O(nk^2)$ .

## References

[dBCvKO08] Mark de Berg, Otfried Cheong, Marc J. van Kreveld, and Mark H. Overmars. *Computational Geometry: Algorithms and Applications*. Springer, Santa Clara, CA, USA, 3rd edition, 2008.