

2.3: NP Completeness

P: Polynomial time problems.

Decision problems.

Circuit Satisfiability

Input: A circuit (boolean) C with n inputs.

Q: Is there a satisfying assignment?

Certificate is a "short" string that proves for a specific instance of a problem the answer is YES!

Verification is "easy"

$\Rightarrow \exists$ a polynomial time alg verify

verify(instance, certificate)
verifies in polynomial time that the answer is yes

$$p(x) = \sum_{i=0}^n \alpha_i x^i \quad O(x^n)$$

$$q(x) = \sum_i \beta_i x^i \quad p(q(x))$$

dup(x): print xxx

```

doit(x):
  y=x
  for i=1..n do
    y=dup(y)
  print y.
    
```

NP: Nondeterministic polynomial

class of all decision problems where there is a poly time verifier.

co-NP All decision problems s.t. if the answer is no, then there is a poly can "prove", and a poly time verifier.

P: Polynomial time

All decision problems solvable in poly time.

$$P \subseteq NP, P \subseteq coNP$$

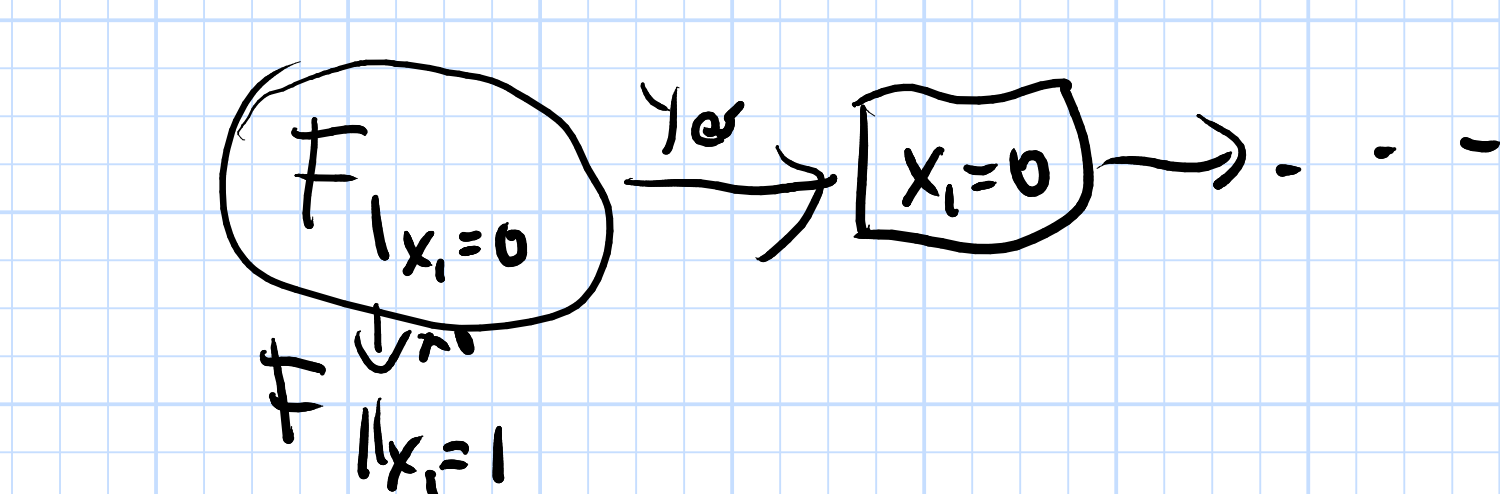
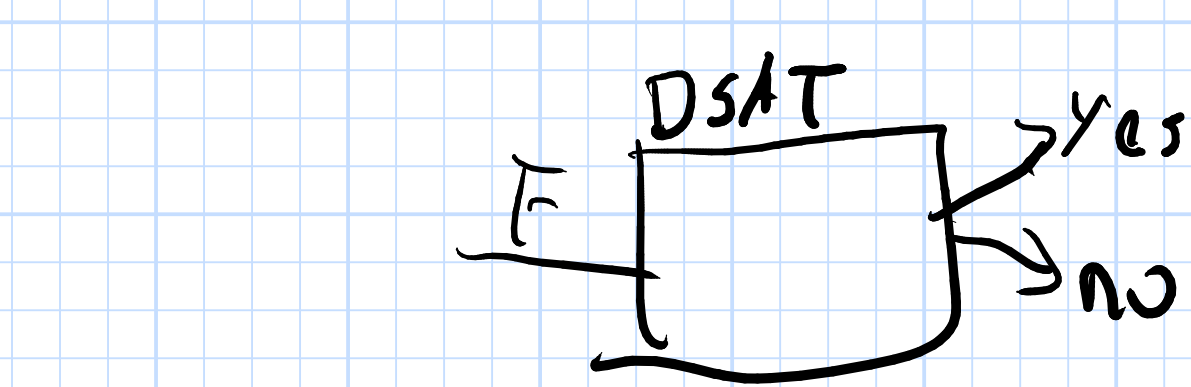
SAT

Input: A boolean formula in CNF form

Q: Is there sat assignment?

$$(x_1 \vee \overline{x_2} \vee x_3) \wedge (x_7 \vee x_{11}) \dots$$

clause or



Self-reduction

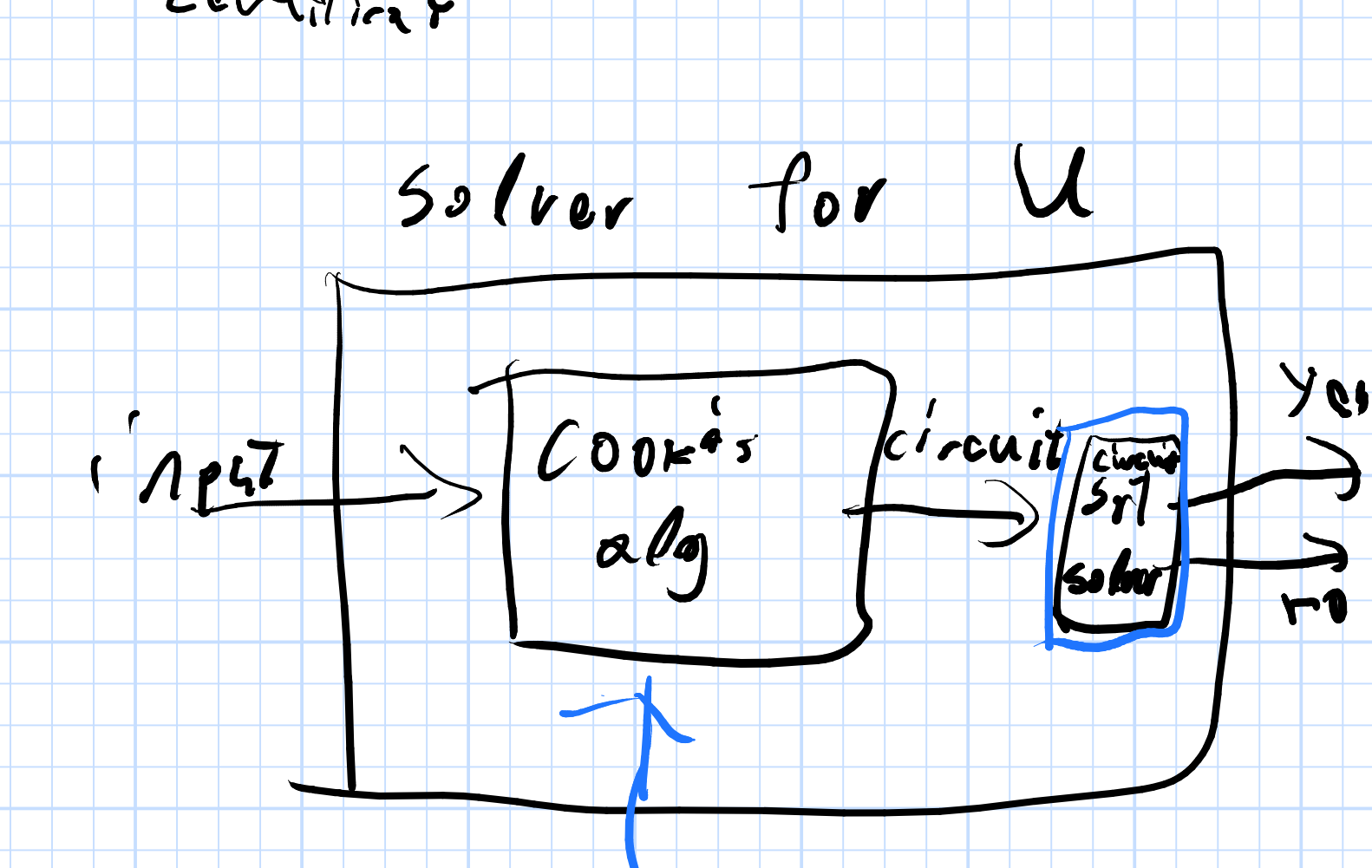
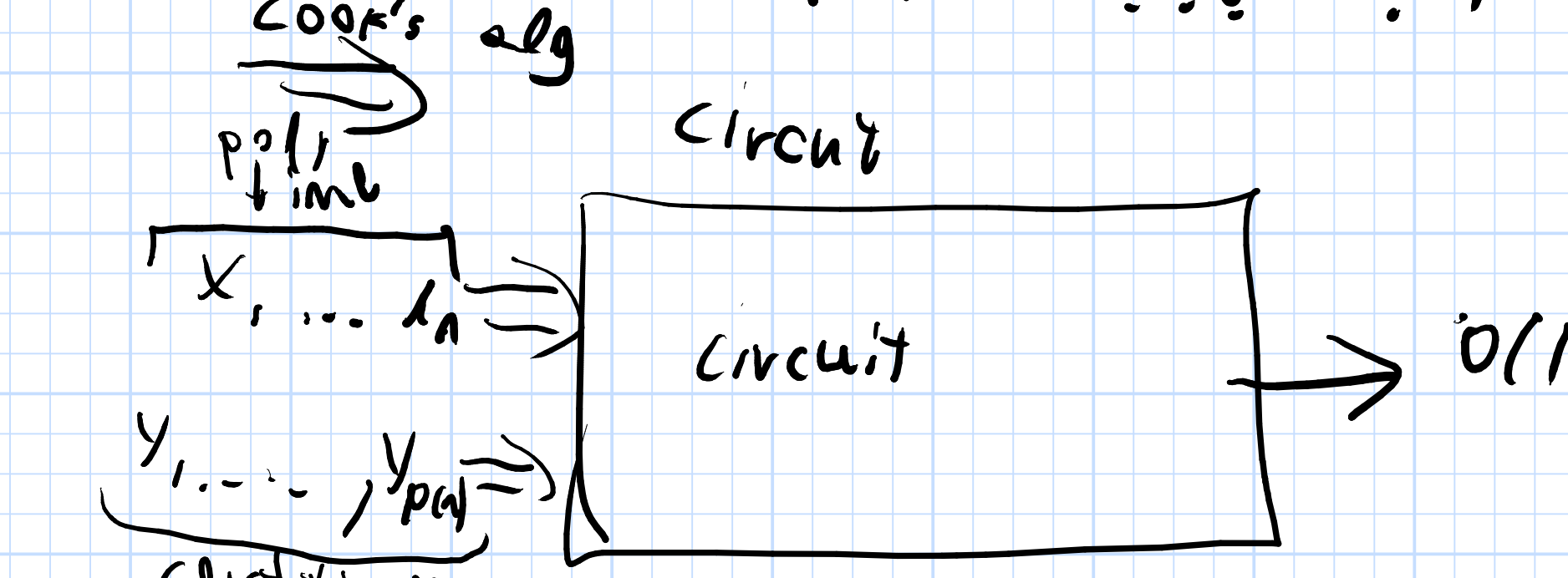
Cook-Levin Theorem 1973

If one can solve circuit SAT in poly time then all the problems in NP can be solved in polynomial time.

U some NP problem

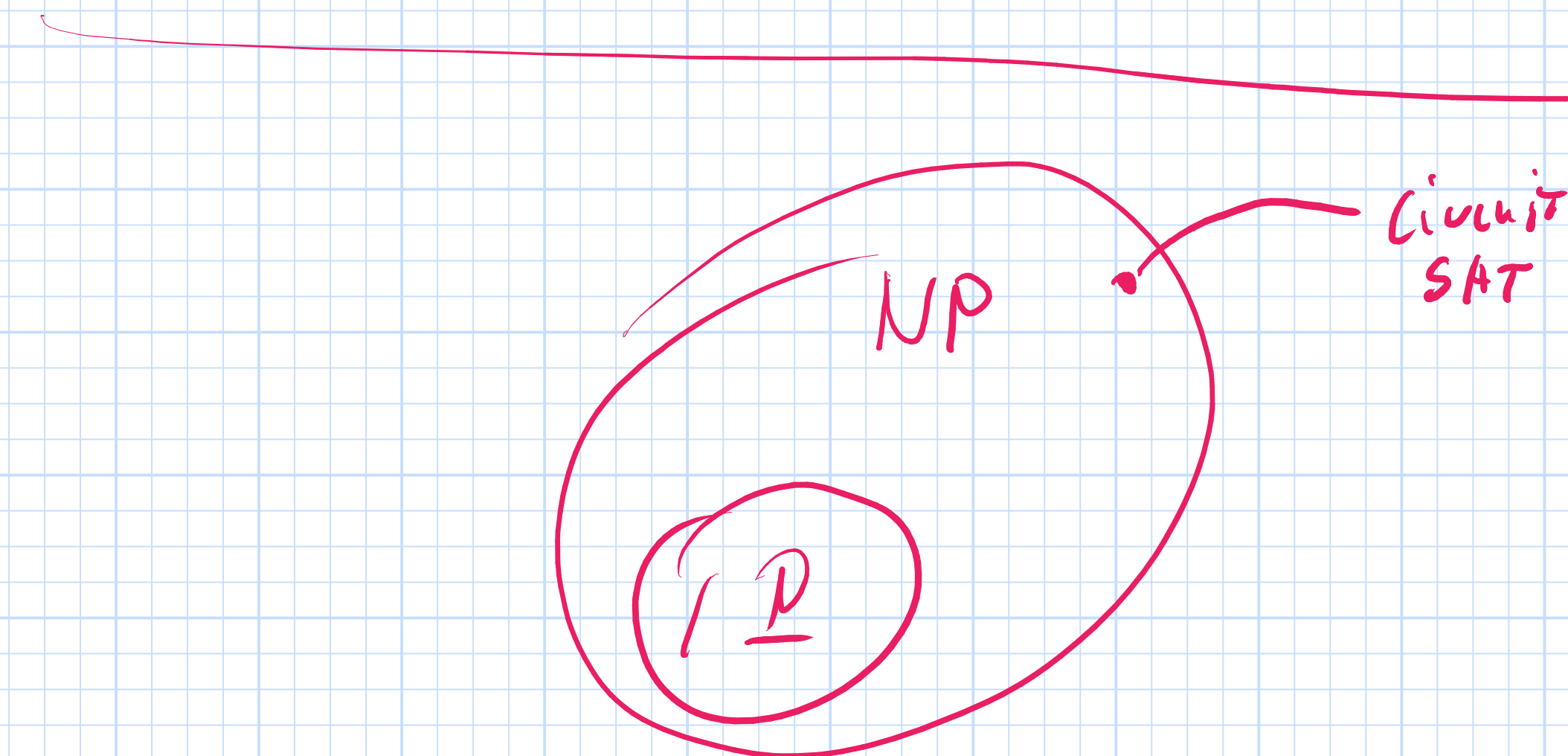
$\Rightarrow \exists$ poly time verifier

verifier($\begin{matrix} x_1 \dots x_n \\ \text{input} \end{matrix}$, $\begin{matrix} y_1 \dots y_{poly} \\ \text{certificate} \end{matrix}$)



know how to do in poly time

Reduction from U to circuit SAT.



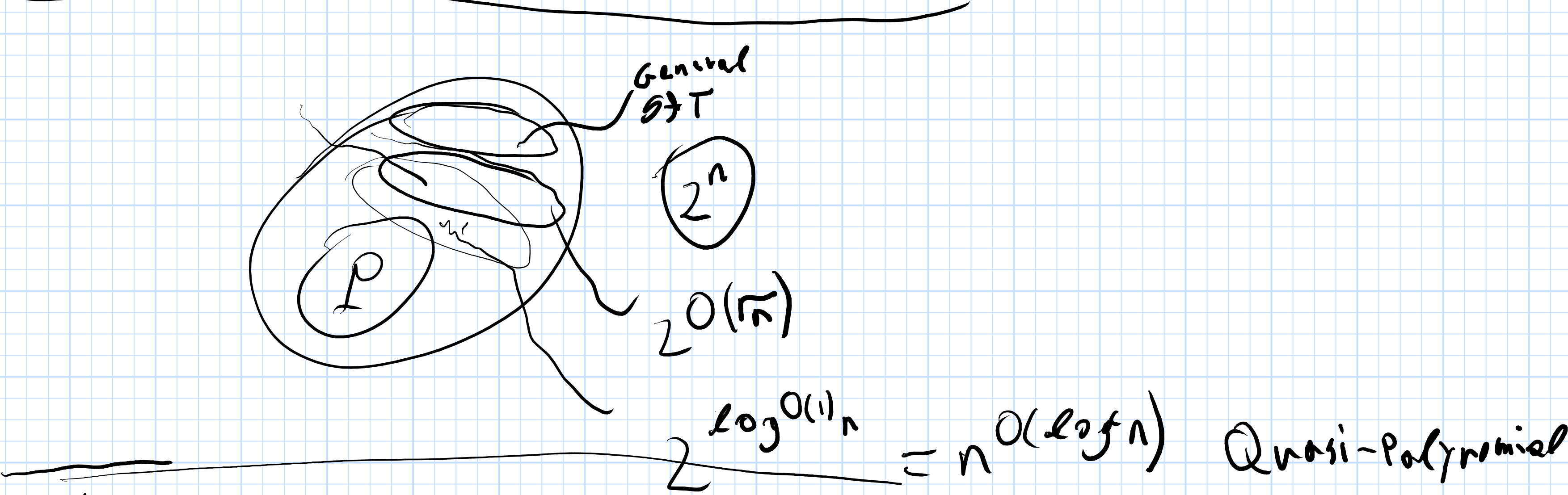
C SAT is poly solvable then $P=NP$.

NP-complete

All problems in NP s.t. if any of them are solvable in poly time then $\forall u \in NP$ are poly solvable.

NP-Hard

All the problems s.t. if any of them are poly solvable $\Rightarrow P=NP$.



Factoring
 n : number
 $n = \prod p_i^{x_i}$
 $\forall i: p_i$ primes.

$2^{O((\#bits \text{ of } n)^{1/3})}$
quantum computers poly time alg shor algorithm!

Problem X is polynomially equivalent to circuit SAT.

A decision problem X is in NPC

$\Leftrightarrow X$ is poly time equivalent to circuit SAT.