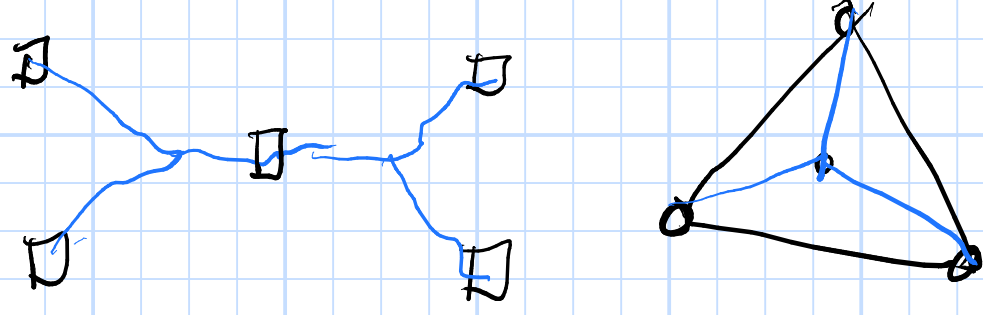


Steiner trees

$G=(V,E)$ undirected, positive weights on the edges.

$X \subseteq V$

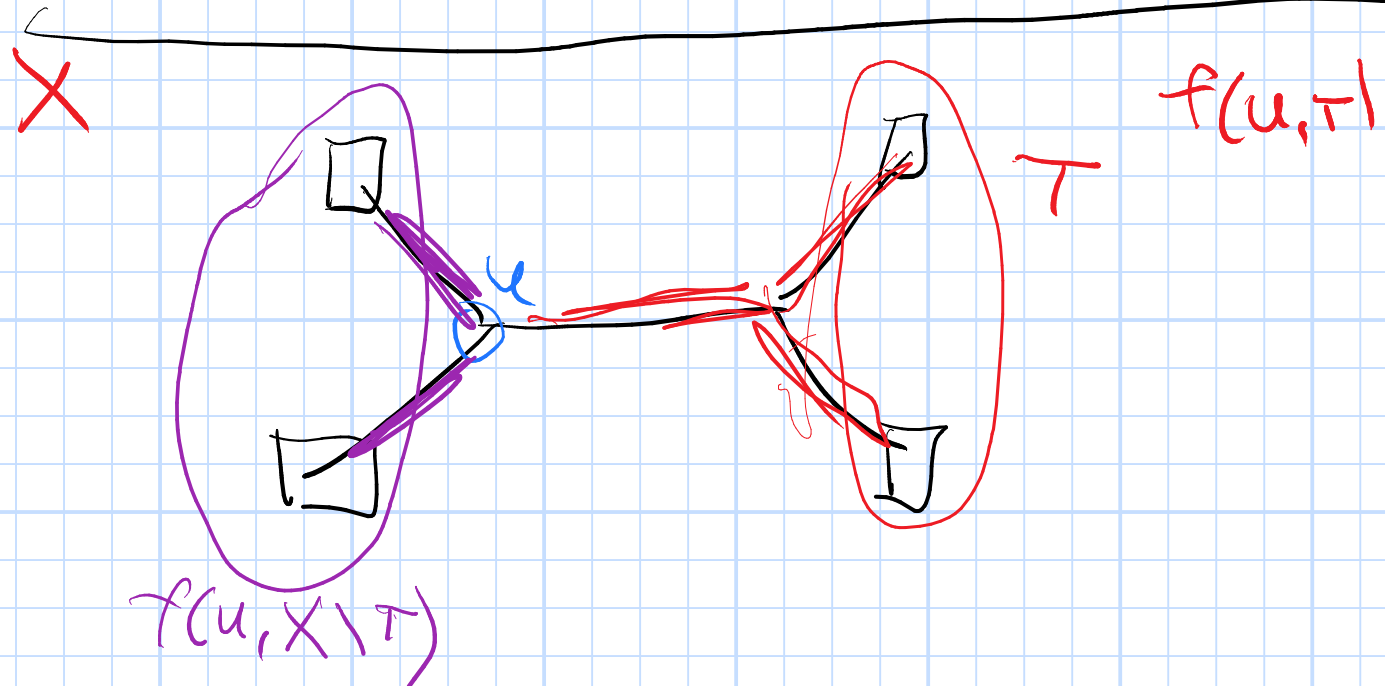
Q: Compute cheapest connected subgraph $H \subseteq G$ s.t. $X \subseteq V(H)$



Steiner tree is NP-Hard.

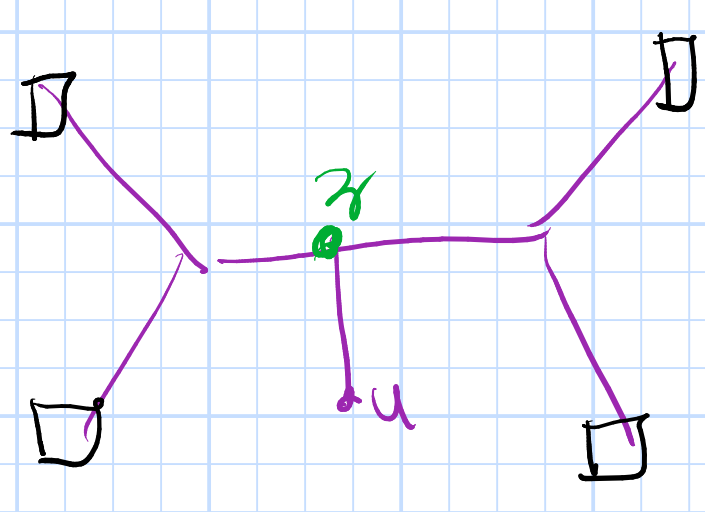
$k=|X|$: parameter of the problem

An algorithm A is FPT (fixed parameter tractable) if its running time is $f(k) \cdot n^{O(1)}$



$f(u, T) \equiv$ price of the cheapest Steiner tree spanning the terminals of T , and $\dots u$.

$$f(u, T) = \begin{cases} 0 & T = \emptyset \\ d_G(u, T) & T = \{u\} \\ f(u, T - u) & u \in T \\ \min_{z \in V} \min_{\substack{S \subseteq T \\ S \neq \emptyset \\ S \neq T}} (f(z, S) + f(z, T \setminus S) + d_G(z, u)) & \end{cases}$$



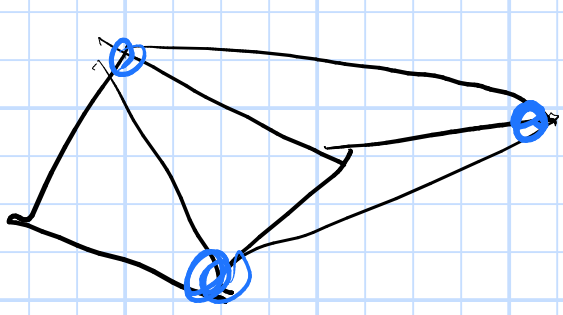
$X \subseteq V$: $k=|X|$ $n=|V(G)|$

$$\begin{matrix} f(u, T) \\ \nearrow \searrow \\ n \quad 2^k \end{matrix}$$

$2^k n$
 $O(2^k n)$ time to complete a recursive call ignoring time spent in recursion

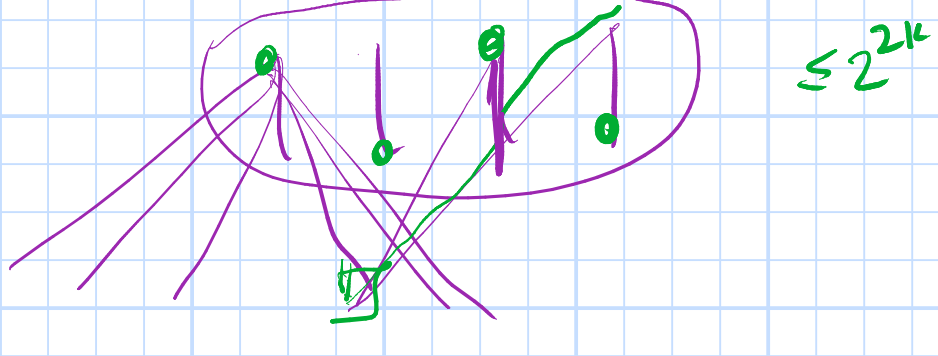
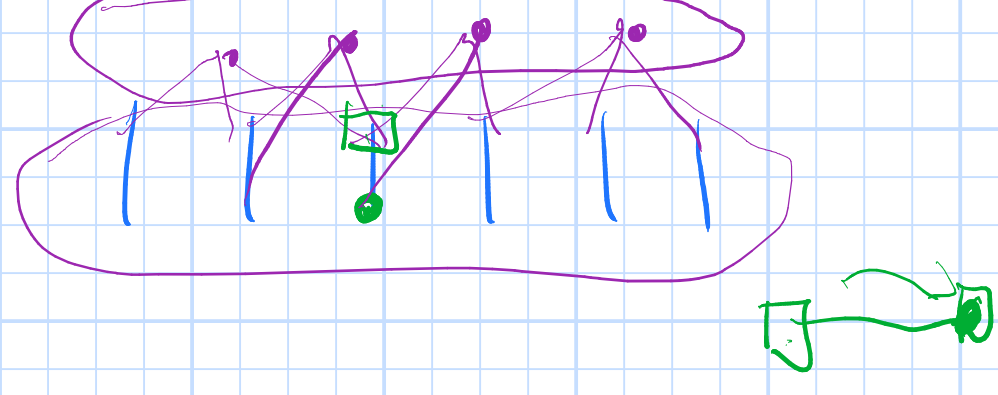
$O(4^k n^2 + n^3)$

FPT example VC: Vertex Cover



$e_1 \parallel e_2 \parallel \dots \parallel e_t$ $t \leq k$

Q: Does the graph has VC of size k ?



$O(n+m) \cdot 2^k$ greedy matching $O(n+m)$

