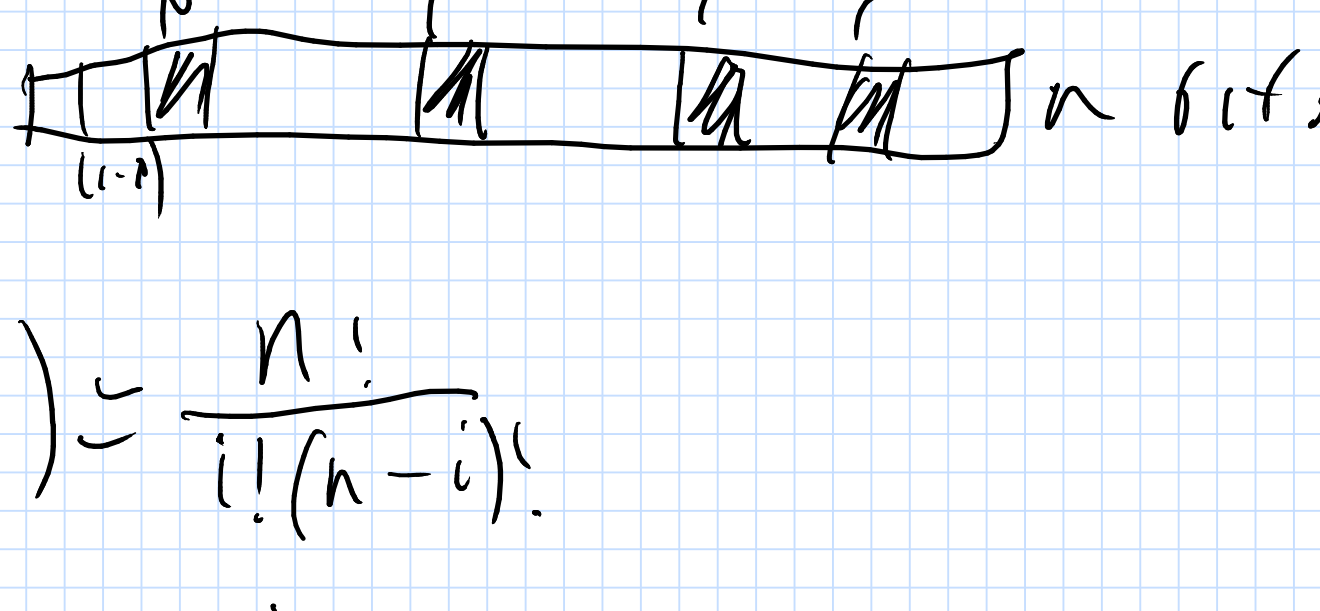


LL: Randomized algorithms

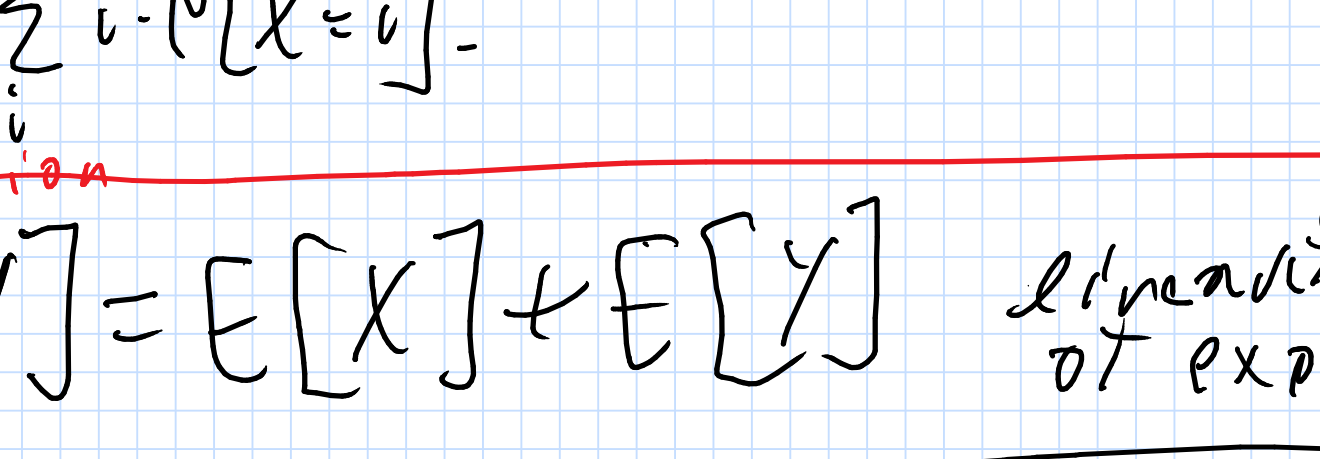
$X \sim \text{Bin}(n, p)$ $n = \#$ of coin flips
 $p =$ probability for heads
 $P[X=i] = \binom{n}{i} p^i (1-p)^{n-i}$



$\binom{n}{i} = \frac{n!}{i!(n-i)!}$

$X \sim \text{Bin}(500, \frac{1}{2})$

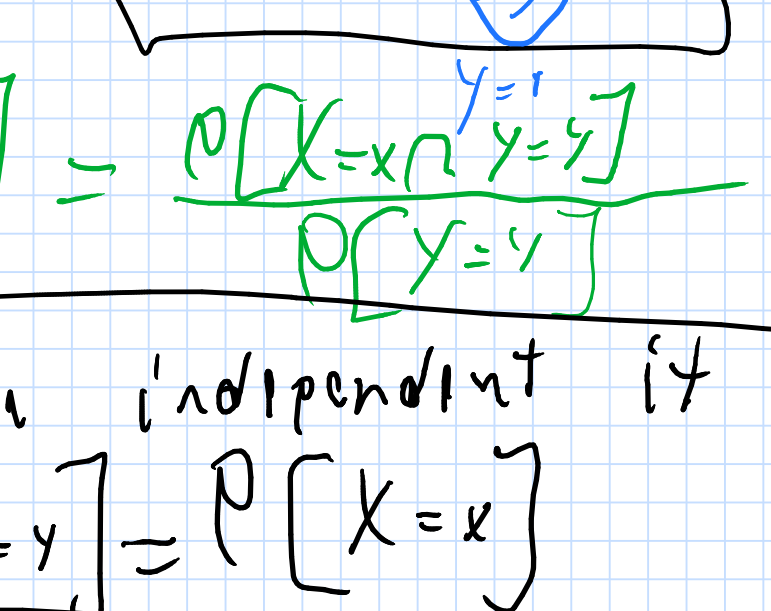
$P[X=i] =$ prob of X get value i .



$E[X] = \sum i \cdot P[X=i]$
expectation
 $E[X+Y] = E[X] + E[Y]$ *linearity of expectation*

Conditional probability

$P[X=x | Y=y] \equiv$ probability of $X=x$ given that $Y=y$.



$P[X=x | Y=y] = \frac{P[X=x, Y=y]}{P[Y=y]}$

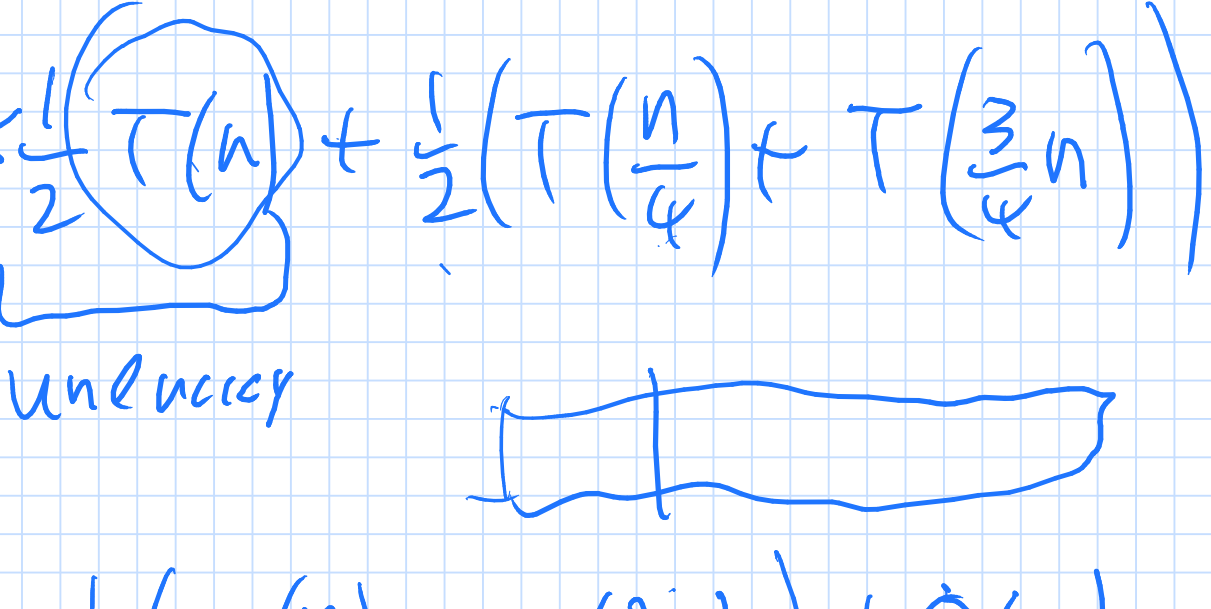
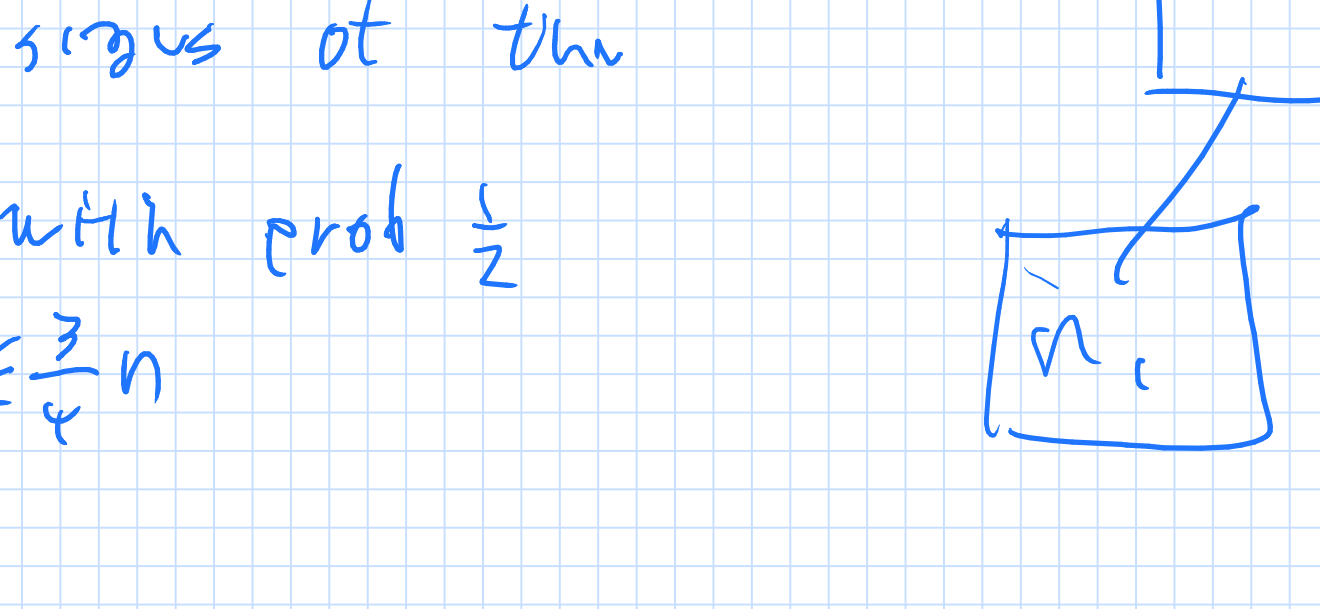
X and Y are independent if $P[X=x | Y=y] = P[X=x]$

If X and Y are independent then $E[XY] = E[X]E[Y]$

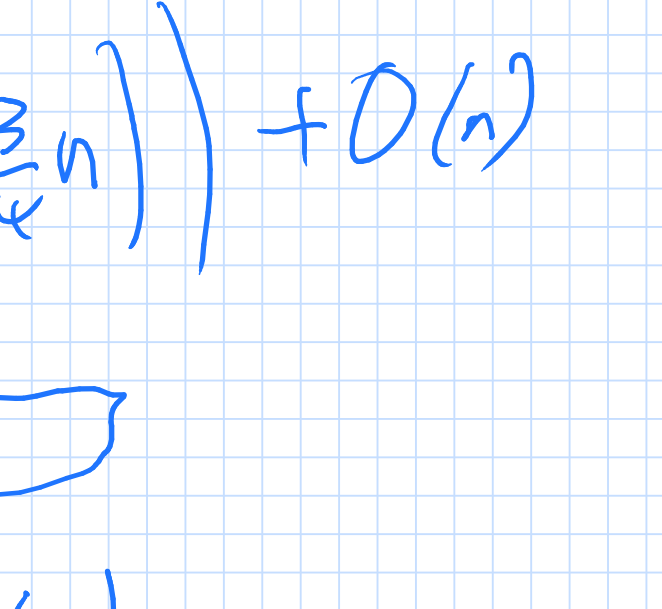
Nuts and bolts

Do quicksort!

$T(n)$: expected running time of QS for n numbers.



n_1, n_2 : sizes of the
 - Lucky with prob $\frac{1}{2}$
 $n_1, n_2 \leq \frac{3}{4}n$

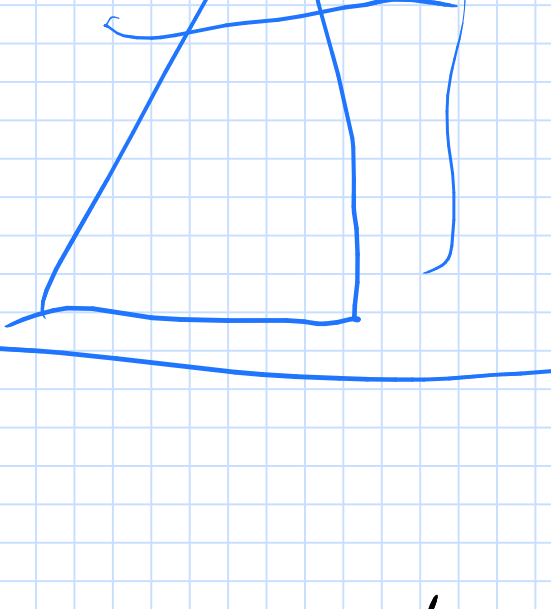


$T(n) \leq \frac{1}{2}T(n) + \frac{1}{2}(T(\frac{n}{4}) + T(\frac{3}{4}n)) + O(n)$
unlucky

$\frac{1}{2}T(n) \leq \frac{1}{2}(T(\frac{n}{4}) + T(\frac{3}{4}n)) + O(n)$

$T(n) \leq T(\frac{n}{4}) + T(\frac{3}{4}n) + O(n)$

$= O(n \log n)$



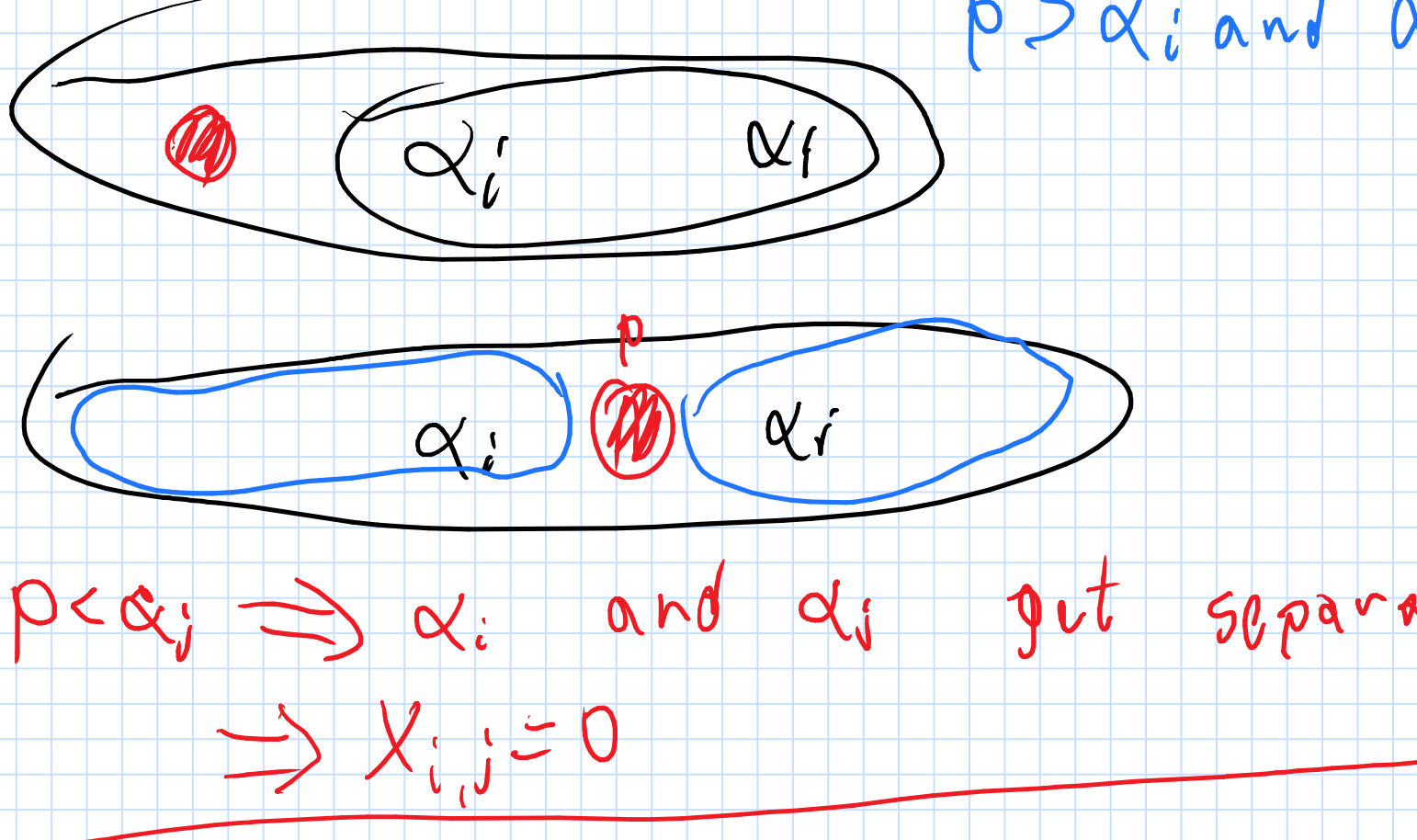
Quicksort Analysis

Indicator variable $X \in \{0,1\}$
 $X=1 \iff$ a certain event happened.

$X_{i,j} = 1 \iff$ QS compared the i th smallest element to the j th smallest element.

RT of QS $= O(\sum X_{i,j})$
 $E[RT] = O(E[\sum X_{i,j}])$
 $E[\sum X_{i,j}] = \sum E[X_{i,j}] = \sum P[X_{i,j}=1]$
 $E[X_{i,j}] = P[X_{i,j}=0] \cdot 0 + P[X_{i,j}=1] \cdot 1 = P[X_{i,j}=1]$

$P[X_{i,j}=1] = \frac{1}{j-i+1}$ $j=i+1$
 $j=i+\frac{n}{4}$



$a_i < p < a_j \implies a_i$ and a_j get separated $\implies X_{i,j}=0$

Run quicksort till a pivot is chosen from the set $[a_i, a_{i+1}, \dots, a_j]$ $a[i,j]$

$P[X_{i,j}=1] = P[\text{pivot} = a_i \text{ or } \text{pivot} = a_j \mid \text{this is the first pivot in } a[i,j]]$
 $= \frac{2}{|a[i,j]|} = \frac{2}{j-i+1}$

$\sum_{i < j} E[X_{i,j}] = \sum_{i < j} P[X_{i,j}]$
 $= \sum_{i < j} \frac{2}{j-i+1}$
 $\leq \sum_{i=0}^n \sum_{s=1}^n \frac{2}{j-i+1} = \sum_{i=1}^n \sum_{s=1}^n \frac{2}{s+i}$
 $= 2n \cdot \sum_{s=1}^n \frac{1}{s+1}$
 $\leq 2n \cdot \sum_{i=1}^n \frac{1}{i}$
 $= 2n H_n = O(n \log n)$
 $H_n = \ln n + O(1)$