Programming Languages and Compilers (CS 421)

Elsa L Gunter

2112 SC, UIUC

http://courses.engr.illinois.edu/cs421

Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

Untyped λ-Calculus

- How do you compute with the λ-calculus?
- Roughly speaking, by substitution:
- $(\lambda\ x.\ e_1)\ e_2 \Rightarrow^* e_1\ [e_2 / x]$
- * Modulo all kinds of subtleties to avoid free variable capture

Transition Semantics for λ-Calculus

$$\frac{E \dashrightarrow E''}{E\,E' \dashrightarrow E''\,E'}$$

- Application (version 1 - Lazy Evaluation)

$$(\lambda\, x\,.\, E)\ E' \dashrightarrow E[E'/x]$$

- Application (version 2 - Eager Evaluation)

$$\frac{E' \dashrightarrow E''}{(\lambda\, x\,.\, E)\ E' \dashrightarrow (\lambda\, x\,.\, E)\ E''}$$

$$\frac{}{(\lambda\, x\,.\, E)\ V \dashrightarrow E[V/x]}$$

V - variable or abstraction (value)

How Powerful is the Untyped λ-Calculus?

- The untyped λ-calculus is Turing Complete
  - Can express any sequential computation
- Problems:
  - How to express basic data: booleans, integers, etc?
  - How to express recursion?
  - Constants, if_then_else, etc, are conveniences; can be added as syntactic sugar

Typed vs Untyped λ-Calculus

- The *pure* λ-calculus has no notion of type: (f f) is a legal expression
- Types restrict which applications are valid
- Types are not syntactic sugar! They disallow some terms
- Simply typed λ-calculus is less powerful than the untyped λ-Calculus: NOT Turing Complete (no general recursion)

α Conversion

- α-conversion:

  λ x. exp --α--> λ y. (exp [y/x])

- Provided that

1. y is not free in exp
2. No free occurrence of x in exp becomes bound in exp when replaced by y

λ x. x (λ y. x y) -X-> λ y. y(λ y.y y)

α Conversion Non-Examples

1. Error: y is not free in termsecond

   λ x. x y --α--> λ y. y y

2. Error: free occurrence of x becomes bound in wrong way when replaced by y

   λ x. λ y. x y --α--> λ y. λ y. y y

   exp (λ y. x y) exp[y/x] (λ y. y y)

But λ x. (λ y. y) x --α--> λ y. (λ y. y) y

And λ y. (λ y. y) y --α--> λ x. (λ y. y) x

Congruence

- Let ~ be a relation on lambda terms. ~ is a congruence if:

1. It is an equivalence relation
2. If $e_1 \sim e_2$ then
   - $(e\ e_1) \sim (e\ e_2)$ and $(e_1 e) \sim (e_2\ e)$
   - $\lambda\ x.\ e_1 \sim \lambda\ x.\ e_2$

α Equivalence

- α equivalence is the smallest congruence containing α conversion

- One usually treats α-equivalent terms as equal - i.e. use α equivalence classes of terms

Example

Show: λ x. (λ y. y x) x ~α~ λ y. (λ x. x y) y

- λ x. (λ y. y x) x --α--> λ z. (λ y. y z) z so
  λ x. (λ y. y x) x ~α~ λ z. (λ y. y z) z
- (λ y. y z) --α--> (λ x. x z) so
  (λ y. y z) ~α~ (λ x. x z) so
  λ z. (λ y. y z) z ~α~ λ z. (λ x. x z) z
- λ z. (λ x. x z) z --α--> λ y. (λ x. x y) y so
  λ z. (λ x. x z) z ~α~ λ y. (λ x. x y) y
- λ x. (λ y. y x) x ~α~ λ y. (λ x. x y) y

Substitution

- Defined on α-equivalence classes of terms
- P [N / x] means replace every free occurrence of x in P by N
  - P called *redex*; N called *residue*
- Provided that no variable free in P becomes bound in P [N / x]
  - Rename bound variables in P to avoid capturing free variables of N

Substitution

- $x\ [N / x] = N$
- $y\ [N / x] = y$ if $y \neq x$
- $(e_1\ e_2)\ [N / x] = ((e_1\ [N / x]\ )\ (e_2\ [N / x]\ ))$
- $(\lambda\ x.\ e)\ [N / x] = (\lambda\ x.\ e)$
- $(\lambda\ y.\ e)\ [N / x] = \lambda\ y.\ (e\ [N / x]\ )$ provided $y \neq x$ and y not free in N
  - Rename y in redex if necessary

Example

$$(\lambda\ y.\ y\ z)\ [(\lambda\ x.\ x\ y)\ /\ z] = ?$$

- Problems?
  - z in redex in scope of y binding
  - y free in the residue
- $(\lambda\ y.\ y\ z)\ [(\lambda\ x.\ x\ y)\ /\ z]$ --α--> $(\lambda\ w.w\ z)\ [(\lambda\ x.\ x\ y)\ /\ z] =$ $\lambda\ w.\ w\ (\lambda\ x.\ x\ y)$

Example

- Only replace free occurrences
- $(\lambda\ y.\ y\ z\ (\lambda\ z.\ z))\ [(\lambda\ x.\ x)\ /\ z] =$

$$\lambda\ y.\ y\ (\lambda\ x.\ x)\ (\lambda\ z.\ z)$$

Not

$$\lambda\ y.\ y\ (\lambda\ x.\ x)\ (\lambda\ z.\ (\lambda\ x.\ x))$$

β reduction

- β Rule: $(\lambda\ x.\ P)\ N$ --β--> $P\ [N\ /x]$

- Essence of computation in the lambda calculus
- Usually defined on α-equivalence classes of terms

Example

- (λ z. (λ x. x y) z) (λ y. y z)
  --β--> (λ x. x y) (λ y. y z)
  --β--> (λ y. y z) y --β--> y z

- (λ x. x x) (λ x. x x)
  --β--> (λ x. x x) (λ x. x x)
  --β--> (λ x. x x) (λ x. x x) --β--> ....

α β Equivalence

- α β equivalence is the smallest congruence containing α equivalence and β reduction
- A term is in *normal form* if no subterm is α equivalent to a term that can be β reduced
- Hard fact (Church-Rosser): if e_1 and e_2 are $\alpha\beta$-equivalent and both are normal forms, then they are α equivalent

Order of Evaluation

- Not all terms reduce to normal forms
- Not all reduction strategies will produce a normal form if one exists

Lazy evaluation:

- Always reduce the left-most application in a top-most series of applications (i.e. Do not perform reduction inside an abstraction)
- Stop when term is not an application, or left-most application is not an application of an abstraction to a term

Example 1

- (λ z. (λ x. x)) ((λ y. y y) (λ y. y y))
- Lazy evaluation:
- Reduce the left-most application:
- (λ z. (λ x. x)) ((λ y. y y) (λ y. y y))
  --β--> (λ x. x)

Eager evaluation

- (Eagerly) reduce left of top application to an abstraction
- Then (eagerly) reduce argument
- Then β-reduce the application

Example 1

- (λ z. (λ x. x))((λ y. y y) (λ y. y y))
- Eager evaluation:
- Reduce the rator of the top-most application to an abstraction: Done.
- Reduce the argument:
- (λ z. (λ x. x))((λ y. y y) (λ y. y y))

--β--> (λ z. (λ x. x))((λ y. y y) (λ y. y y))

--β--> (λ z. (λ x. x))((λ y. y y) (λ y. y y))…

Example 2

- $(\lambda\ x.\ x\ x)((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$
- Lazy evaluation:

$(\lambda\ x.\ x\ \ x\ )((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$ --β-->

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

Example 2

- $(\lambda\ x.\ x\ x)((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$
- Lazy evaluation:

$(\lambda\ x.\ x\ \ x\ )((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$ --β-->

$((\lambda\ y.\ y\ \ y\ )\ (\lambda\ z.\ z))\ ((\lambda\ y.\ y\ \ y\ )\ (\lambda\ z.\ z))$

Example 2

- $(\lambda\ x.\ x\ x)((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$
- Lazy evaluation:

$(\lambda\ x.\ x\ \ x\ )((\lambda\ y.\ y\ y)\ (\lambda\ z.\ z))$ --β-->

$((\lambda\ y.\ y\ \ y\ )\ (\lambda\ z.\ z))\ ((\lambda\ y.\ y\ \ y\ )\ (\lambda\ z.\ z)$

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

--β--> (λ z. z) ((λ y. y y) (λ z. z))

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

--β--> (λ z. z) ((λ y. y y) (λ z. z)) --β-->

(λ y. y y) (λ z. z)

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

(((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> (((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

--β--> (λ z. z) ((λ y. y y) (λ z. z)) --β-->

(λ y. y y) (λ z. z)

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Lazy evaluation:

(λ x. x x)((λ y. y y) (λ z. z)) --β-->

((λ y. y y) (λ z. z)) ((λ y. y y) (λ z. z))

--β--> ((λ z. z) (λ z. z))((λ y. y y) (λ z. z))

--β--> (λ z. z) ((λ y. y y) (λ z. z)) --β-->

(λ y. y y) (λ z. z) ~β~ λ z. z

Example 2

- (λ x. x x)((λ y. y y) (λ z. z))
- Eager evaluation:

(λ x. x x) ((λ y. y y) (λ z. z)) --β-->

(λ x. x x) ((λ z. z) (λ z. z)) --β-->

(λ x. x x) (λ z. z) --β-->

(λ z. z) (λ z. z) --β--> λ z. z

η (Eta) Reduction

- η Rule: λ x. f x --η--> f if x not free in f
  - Can be useful in each direction
  - Not valid in Ocaml
    - recall lambda-lifting and side effects
  - Not equivalent to (λ x. f) x --> f (inst of β)

- Example: λ x. (λ y. y) x --η--> λ y. y

Untyped λ-Calculus

- Only three kinds of expressions:
  - Variables: x, y, z, w, …
  - Abstraction: λ x. e
    (Function creation)
  - Application: e_1 e_2

How to Represent (Free) Data Structures (First Pass - Enumeration Types)

- Suppose τ is a type with *n* constructors: $C_1,\ldots,C_n$ (no arguments)
- Represent each term as an abstraction:
- Let $C_i \rightarrow \lambda\, x_1 \ldots x_n.\, x_i$
- Think: you give me what to return in each case (think match statement) and I'll return the case for the *i*th constructor

How to Represent Booleans

- bool = True | False
- True $\rightarrow \lambda\, x_1.\ \lambda\, x_2.\ x_1 \quad \equiv_\alpha \quad \lambda\, x.\ \lambda\, y.\ x$
- False $\rightarrow \lambda\, x_1.\ \lambda\, x_2.\ x_2 \quad \equiv_\alpha \quad \lambda\, x.\ \lambda\, y.\ y$

- Notation
  - Will write

  $\lambda\, x_1 \ldots x_n.\ e$ for $\lambda\, x_1.\ \ldots\ \lambda x_n.\ e$

  $e_1\, e_2 \ldots e_n$ for $(\ldots(e_1\, e_2)\ldots e_n)$

Functions over Enumeration Types

- Write a "match" function
- match e with C_1 -> x_1
  | …
  | C_n -> x_n

$\rightarrow \quad \lambda\, x_1 \ldots x_n\, e.\ e\ x_1 \ldots x_n$

- Think: give me what to do in each case and give me a case, and I'll apply that case

Functions over Enumeration Types

- type $\tau = C_1 | \ldots | C_n$
- match e with C_1 -> x_1

  | …

  | C_n -> x_n

- $match\tau = \lambda\, x_1 \ldots x_n\, e.\ e\ x_1 \ldots x_n$

- e = expression (single constructor)

  x_i is returned if e = C_i

match for Booleans

- bool = True | False
- True $\rightarrow \lambda\ x_1\ x_2.\ x_1 \quad \equiv_\alpha \quad \lambda\ x\ y.\ x$
- False $\rightarrow \lambda\ x_1\ x_2.\ x_2 \quad \equiv_\alpha \quad \lambda\ x\ y.\ y$

- match$_{bool}$ = ?

match for Booleans

- bool = True | False
- True $\rightarrow \lambda\, x_1\, x_2.\, x_1 \quad \equiv_\alpha \quad \lambda\, x\, y.\, x$
- False $\rightarrow \lambda\, x_1\, x_2.\, x_2 \quad \equiv_\alpha \quad \lambda\, x\, y.\, y$

- $\text{match}_{bool} = \lambda\, x_1\, x_2\, e.\; e\; x_1\; x_2$

  $\equiv_\alpha \quad \lambda\, x\, y\, b.\; b\; x\; y$

How to Write Functions over Booleans

- if b then x_1 else x_2 $\rightarrow$
- if_then_else b x_1 x_2 = b x_1 x_2
- if_then_else $\equiv \lambda$ b x_1 x_2 . b x_1 x_2

How to Write Functions over Booleans

- Alternately:
- if b then x_1 else x_2 =
  match b with True -> x_1 | False -> x_2 $\rightarrow$
  $\text{match}_{\text{bool}}\ x_1\ x_2\ b$ =
  $(\lambda\ x_1\ x_2\ b\ .\ b\ x_1\ x_2\ )\ x_1\ x_2\ b = b\ x_1\ x_2$
- if_then_else
  $\equiv \lambda\ b\ x_1\ x_2.\ (\text{match}_{\text{bool}}\ x_1\ x_2\ b)$
  $= \lambda\ b\ x_1\ x_2.\ (\lambda\ x_1\ x_2\ b\ .\ b\ x_1\ x_2\ )\ x_1\ x_2\ b$
  $= \lambda\ b\ x_1\ x_2.\ b\ x_1\ x_2$

Example:

not b

= match b with True -> False | False -> True

$\rightarrow$ $(\text{match}_{\text{bool}})$ False True b

= $(\lambda\ x_1\ x_2\ b\ .\ b\ x_1\ x_2\ )\ (\lambda\ x\ y.\ y)\ (\lambda\ x\ y.\ x)\ b$

= $b\ (\lambda\ x\ y.\ y)(\lambda\ x\ y.\ x)$

- not $\equiv \lambda\ b.\ b\ (\lambda\ x\ y.\ y)(\lambda\ x\ y.\ x)$
- Try and, or

and or

How to Represent (Free) Data Structures (Second Pass - Union Types)

- Suppose τ is a type with *n* constructors:
  type $\tau \;= C_1\, t_{11} \ldots t_{1k} \mid \ldots \mid C_n\, t_{n1} \ldots t_{nm,}$
- Represent each term as an abstraction:
- $C_i\, t_{i1} \ldots t_{ij,} \rightarrow \lambda\, x_1 \ldots x_n.\, x_i\, t_{i1} \ldots t_{ij,}$
- $C_i \rightarrow \lambda\, t_{i1} \ldots t_{ij,}\, x_1 \ldots x_n\, .\, x_i\, t_{i1} \ldots t_{ij,}$
- Think: you need to give each constructor its arguments fisrt

How to Represent Pairs

- Pair has one constructor (comma) that takes two arguments
- type (α,β)pair = (,) α β
- (a , b) --> λ x . x a b
- (_ , _) --> λ a b x . x a b

Functions over Union Types

- Write a "match" function
- match e with $C_1\ y_1 \ldots y_{m1}$ -> $f_1\ y_1 \ldots y_{m1}$
  $| \ldots$
  $| \ C_n\ y_1 \ldots y_{mn}$ -> $f_n\ y_1 \ldots y_{mn}$
- $match\tau \rightarrow \ \lambda\, f_1 \ldots f_n\, e.\ e\ f_1 \ldots f_n$
- Think: give me a function for each case and give me a case, and I'll apply that case to the appropriate fucntion with the data in that case

Functions over Pairs

- $match_{pair} = \lambda$ f p. p f

- fst p = match p with (x,y) -> x
- fst $\to \lambda$ p. $match_{pair}$ (λ x y. x)

  = (λ f p. p f) (λ x y. x) = λ p. p (λ x y. x)

- snd $\to \lambda$ p. p (λ x y. y)

How to Represent (Free) Data Structures (Third Pass - Recursive Types)

- Suppose τ is a type with n constructors:

  type $\tau = C_1\, t_{11} \ldots t_{1k} \mid \ldots \mid C_n\, t_{n1} \ldots t_{nm},$
- Suppose $t_{ih} : \tau$ (ie. is recursive)
- In place of a value t_{ih} have a function to compute the recursive value $r_{ih}\, x_1 \ldots x_n$
- $C_i\, t_{i1} \ldots r_{ih} \ldots t_{ij} \rightarrow \lambda x_1 \ldots x_n .\, x_i\, t_{i1} \ldots (r_{ih}\, x_1 \ldots x_n) \ldots t_{ij}$
- $C_i \rightarrow \lambda\, t_{i1} \ldots r_{ih} \ldots t_{ij}\, x_1 \ldots x_n .x_i\, t_{i1} \ldots (r_{ih}\, x_1 \ldots x_n) \ldots t_{ij},$

How to Represent Natural Numbers

- nat = Suc nat | 0
- $\overline{\text{Suc}} = \lambda$ n f x. f (n f x)
- $\overline{\text{Suc n}} = \lambda$ f x. f (n f x)
- $\overline{0} = \lambda$ f x. x
- Such representation called *Church Numerals*

Some Church Numerals

- $\overline{\text{Suc } 0}$ = (λ n f x. f (n f x)) (λ f x. x) -->

λ f x. f ((λ f x. x) f x) -->

λ f x. f ((λ x. x) x) --> λ f x. f x

Apply a function to its argument once

Some Church Numerals

- $\overline{\text{Suc(Suc 0)}}$ = (λ n f x. f (n f x)) (Suc 0) -->
(λ n f x. f (n f x)) (λ f x. f x) -->
λ f x. f ((λ f x. f x) f x)) -->
λ f x. f ((λ x. f x) x)) --> λ f x. f (f x)

Apply a function twice

In general $\overline{n}$ = λ f x. f (… (f x)…) with n applications of f

Primitive Recursive Functions

- Write a "fold" function
- fold $f_1 \ldots f_n$ = match e
  with $C_1\ y_1 \ldots y_{m1}$ -> $f_1\ y_1 \ldots y_{m1}$
  | …
  | $Ci\ y_1 \ldots r_{ij} \ldots y_{in}$ -> $f_n\ y_1 \ldots (\text{fold } f_1 \ldots f_n\ r_{ij}) \ldots y_{mn}$
  | …
  | $C_n\ y_1 \ldots y_{mn}$ -> $f_n\ y_1 \ldots y_{mn}$
- *fold*$\tau \rightarrow \lambda f_1 \ldots f_n e.\ e\ f_1 \ldots f_n$
- Match in non recursive case a degenerate version of fold

Primitive Recursion over Nat

- fold f z n=
- match n with 0 -> z
-                 | Suc m -> f (fold f z m)
- $\overline{\text{fold}} \equiv \lambda\ f\ z\ n.\ n\ f\ z$
- $\overline{\text{is\_zero}}\ \ \overline{n} = \overline{\text{fold}}\ (\lambda\ r.\ \overline{\text{False}})\ \overline{\text{True}}\ \overline{n}$
- $= (\lambda\ f\ x.\ f^{n}\ x)\ (\lambda\ r.\ \text{False})\ \text{True}$
- $= ((\lambda\ r.\ \text{False})^{n})\ \text{True}$
- $\equiv$ if n = 0 then True else False

Adding Church Numerals

- $\overline{n} \equiv \lambda f\, x.\ f^{\,n}\, x$ and $\overline{m} \equiv \lambda f\, x.\ f^{\,m}\, x$

- $\overline{n + m} = \lambda f\, x.\ f^{\,(n+m)}\, x$
  $= \lambda f\, x.\ f^{\,n}\, (f^{\,m}\, x) = \lambda f\, x.\ \overline{n}\, f\, (\overline{m}\, f\, x)$

- $\overline{+} \equiv \lambda\, n\, m\, f\, x.\ n\, f\, (m\, f\, x)$

- Subtraction is harder

Multiplying Church Numerals

- $\overline{n} \equiv \lambda f\, x.\ f^{\,n}\, x$ and $\overline{m} \equiv \lambda f\, x.\ f^{\,m}\, x$
- $\overline{n * m} = \lambda f\, x.\ (f^{\,n * m})\, x = \lambda f\, x.\ (f^{\,m})^n\, x$
  $= \lambda f\, x.\ \overline{n}\, (\overline{m}\, f)\, x$

$\overline{*} \equiv \lambda n\, m\, f\, x.\ n\, (m\, f)\, x$

Predecessor

- let pred_aux n =
 match n with 0 -> (0,0)
 | Suc m
-> (Suc(fst(pred_aux m)), fst(pred_aux m)
= fold (λ r. (Suc(fst r), fst r)) (0,0) n

- pred ≡ λ n. snd (pred_aux n) n =
λ n. snd (fold (λ r.(Suc(fst r), fst r)) (0,0) n)

Recursion

- Want a λ-term Y such that for all term R we have
- Y R = R (Y R)
- Y needs to have replication to "remember" a copy of R
- $Y = \lambda\ y.\ (\lambda\ x.\ y(x\ x))\ (\lambda\ x.\ y(x\ x))$
- $Y\ R = (\lambda\ x.\ R(x\ x))\ (\lambda\ x.\ R(x\ x))$

  $= R\ ((\lambda\ x.\ R(x\ x))\ (\lambda\ x.\ R(x\ x)))$
- Notice: Requires lazy evaluation

Factorial

- Let F = λ f n. if n = 0 then 1 else n * f (n - 1)

Y F 3 = F (Y F) 3

= if 3 = 0 then 1 else 3 * ((Y F)(3 - 1))

= 3 * (Y F) 2 = 3 * (F(Y F) 2)

= 3 * (if 2 = 0 then 1 else 2 * (Y F)(2 - 1))

= 3 * (2 * (Y F)(1)) = 3 * (2 * (F(Y F) 1)) =…

= 3 * 2 * 1 * (if 0 = 0 then 1 else 0*(Y F)(0 -1))

= 3 * 2 * 1 * 1 = 6

Y in OCaml

```
# let rec y f = f (y f);;
val y : ('a -> 'a) -> 'a = <fun>
# let mk_fact =
   fun f n -> if n = 0 then 1 else n * f(n-1);;
val mk_fact : (int -> int) -> int -> int = <fun>
# y mk_fact;;
Stack overflow during evaluation (looping
   recursion?).
```

Eager Eval Y in Ocaml

```
# let rec y f x = f (y f) x;;
val y : (('a -> 'b) -> 'a -> 'b) -> 'a -> 'b
   = <fun>
# y mk_fact;;
- : int -> int = <fun>
# y mk_fact 5;;
- : int = 120
```

- Use recursion to get recursion

Some Other Combinators

- For your general exposure

- $I = \lambda\, x\, .\, x$
- $K = \lambda\, x.\, \lambda\, y.\, x$
- $K_* = \lambda\, x.\, \lambda\, y.\, y$
- $S = \lambda\, x.\, \lambda\, y.\, \lambda\, z.\, x\, z\, (y\, z)$