

## Programming Languages and Compilers (CS 421)

Elsa L Gunter

2112 SC, UIUC



<https://courses.engr.illinois.edu/cs421/sp2023>

Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

2/8/23

1

## Terms

- A function is in **Direct Style** when it returns its result back to the caller.
- A function is in **Continuation Passing Style** when it, and every function call in it, passes its result to another function.
- Instead of returning the result to the caller, we pass it forward to another function giving the computation after the call.

2/8/23

3

## Simple Functions Taking Continuations

- Given a primitive operation, can convert it to pass its result forward to a continuation

- Examples:

```
# let subk (x, y) k = k(x - y);;
```

```
val subk : int * int -> (int -> 'a) -> 'a = <fun>
```

```
# let eqk (x, y) k = k(x = y);;
```

```
val eqk : 'a * 'a -> (bool -> 'b) -> 'b = <fun>
```

```
# let timesk (x, y) k = k(x * y);;
```

```
val timesk : int * int -> (int -> 'a) -> 'a = <fun>
```

2/8/23

4

## Nesting Continuations

```
# let add_triple (x, y, z) = (x + y) + z;;
```

```
val add_triple : int * int * int -> int = <fun>
```

```
# let add_triple (x,y,z)=let p = x + y in p + z;;
```

```
val add_triple : int * int * int -> int = <fun>
```

```
# let add_triple_k (x, y, z) k =
```

```
  addk (x, y) (fun p -> addk (p, z) k);;
```

```
val add_triple_k: int * int * int -> (int -> 'a) -> 'a = <fun>
```

2/8/23

5

## add\_three: a different order

```
■ # let add_triple (x, y, z) = x + (y + z);;
```

- How do we write `add_triple_k` to use a different order?

```
■ let add_triple_k (x, y, z) k =
```

2/8/23

6

## add\_three: a different order

```
■ # let add_triple (x, y, z) = x + (y + z);;
```

- How do we write `add_triple_k` to use a different order?

```
■ let add_triple_k (x, y, z) k =  
  addk (y,z) (fun r -> addk(x,r) k)
```

2/8/23

7

## add\_three: a different order

- # let add\_triple (x, y, z) = x + (y + z);;
- How do we write add\_triple\_k to use a different order?
- let add\_triple\_k (x, y, z) k =  
addk (y,z) (fun r -> addk(x,r) k)

2/8/23

8

## add\_three: a different order

- # let add\_triple (x, y, z) = x + (y + z);;
- How do we write add\_triple\_k to use a different order?
- let add\_triple\_k (x, y, z) k =  
addk (y,z) (fun r -> addk(x,r) k)

2/8/23

9

## Recursive Functions

### Recall:

- # let rec factorial n =  
if n = 0 then 1 else n \* factorial (n - 1);;
- val factorial : int -> int = <fun>
- # factorial 5;;
- - : int = 120

2/8/23

10

## Recursive Functions

- # let rec factorial n =  
let b = (n = 0) in (\* First computation \*)  
if b then 1 (\* Returned value \*)  
else let s = n - 1 in (\* Second computation \*)  
let r = factorial s in (\* Third computation \*)  
n \* r (\* Returned value \*) ;;
- val factorial : int -> int = <fun>
- # factorial 5;;
- - : int = 120

2/8/23

11

## Recursive Functions

- # let rec factorialk n k =  
eqk (n, 0)  
(fun b -> (\* First computation \*)  
if b then k 1 (\* Passed value \*)  
else subk (n, 1) (\* Second computation \*)  
(fun s -> factorialk s (\* Third computation \*)  
(fun r -> timesk (n, r) k))) (\* Passed value \*)
- val factorialk : int -> (int -> 'a) -> 'a = <fun>
- # factorialk 5 report;;
- 120
- - : unit = ()

2/8/23

12

## Recursive Functions

- To make recursive call, must build intermediate continuation to
  - take recursive value: r
  - build it to final result: n \* r
  - And pass it to final continuation:
    - times (n, r) k = k (n \* r)

2/8/23

13

## Recursive Functions

```
# let rec factorialk n k =
  eqk (n, 0)
  (fun b -> (* First computation *)
    if b then k 1 (* Passed value *)
    else subk (n, 1) (* Second computation *)
    (fun s -> factorialk s (* Third computation *)
      (fun r -> timesk (n, r) k))) (* Passed value *)
val factorialk : int -> (int -> 'a) -> 'a = <fun>
# factorialk 5 report;;
120
- : unit = ()
```

2/8/23

14

## Example: CPS for length

```
let rec length list = match list with [] -> 0
  | (a :: bs) -> 1 + length bs
```

What is the let-expanded version of this?

2/8/23

15

## Example: CPS for length

```
let rec length list = match list with [] -> 0
  | (a :: bs) -> 1 + length bs
```

What is the let-expanded version of this?

```
let rec length list = match list with [] -> 0
  | (a :: bs) -> let r1 = length bs in 1 + r1
```

2/8/23

17

## Example: CPS for length

```
#let rec length list = match list with [] -> 0
  | (a :: bs) -> let r1 = length bs in 1 + r1
```

What is the CSP version of this?

2/8/23

18

## Example: CPS for length

```
#let rec length list = match list with [] -> 0
  | (a :: bs) -> let r1 = length bs in 1 + r1
```

What is the CSP version of this?

```
#let rec lengthk list k = match list with [ ] -> k 0
  | x :: xs -> lengthk xs (fun r -> addk (r,1) k);;
```

```
val lengthk : 'a list -> (int -> 'b) -> 'b = <fun>
```

```
# lengthk [2;4;6;8] report;;
```

4

```
- : unit = ()
```

2/8/23

19

## CPS for sum

```
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> x + sum xs ;;
```

```
val sum : int list -> int = <fun>
```

2/8/23

20

## CPS for sum

```
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> x + sum xs ;;
val sum : int list -> int = <fun>
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> let r1 = sum xs in x + r1;;
```

2/8/23

21

## CPS for sum

```
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> x + sum xs ;;
val sum : int list -> int = <fun>
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> let r1 = sum xs in x + r1;;
val sum : int list -> int = <fun>
# let rec sumk list k = match list with [ ] -> k 0
  | x :: xs -> sumk xs (fun r1 -> addk x r1 k);;
```

2/8/23

22

## CPS for sum

```
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> x + sum xs ;;
val sum : int list -> int = <fun>
# let rec sum list = match list with [ ] -> 0
  | x :: xs -> let r1 = sum xs in x + r1;;
val sum : int list -> int = <fun>
# let rec sumk list k = match list with [ ] -> k 0
  | x :: xs -> sumk xs (fun r1 -> addk (x, r1) k);;
val sumk : int list -> (int -> 'a) -> 'a = <fun>
# sumk [2;4;6;8] report;;
20
- : unit = ()
```

2/8/23

23

## CPS for Higher Order Functions

- In CPS, every procedure / function takes a continuation to receive its result
- Procedures passed as arguments take continuations
- Procedures returned as results take continuations
- CPS version of higher-order functions must expect input procedures to take continuations

2/8/23

24

## Example: all

```
#let rec all (p, l) = match l with [ ] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
```

2/8/23

25

## Example: all

```
#let rec all (p, l) = match l with [ ] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k =
```

2/8/23

26

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> true
```

2/8/23

27

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> k true
```

2/8/23

28

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) ->
```

2/8/23

29

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) -> pk x
```

2/8/23

30

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) -> pk x
    (fun b -> if b then           else
    )
```

2/8/23

31

## Example: all

```
#let rec all (p, l) = match l with [] -> true
  | (x :: xs) -> let b = p x in
    if b then all (p, xs) else false
val all : ('a -> bool) -> 'a list -> bool = <fun>
■ What is the CPS version of this?
#let rec allk (pk, l) k = match l with [] -> k true
  | (x :: xs) -> pk x
    (fun b -> if b then allk (pk, xs) k else k
false)
val allk : ('a -> (bool -> 'b) -> 'b) * 'a list ->
(bool -> 'b) -> 'b = <fun>
```

2/8/23

32

## Terminology: Review

- A function is in **Direct Style** when it returns its result back to the caller.
- A function is in **Continuation Passing Style** when it, and every function call in it, passes its result to another function.
- A **Tail Call** occurs when a function returns the result of another function call without any more computations (eg tail recursion)
- Instead of returning the result to the caller, we pass it forward to another function giving the computation after the call.

2/8/23

33

## CPS Transformation

- Step 1: Add continuation argument to any function definition:
  - $\text{let } f \text{ arg} = e \Rightarrow \text{let } f \text{ arg } k = e$
  - Idea: Every function takes an extra parameter saying where the result goes
- Step 2: A simple expression in tail position should be passed to a continuation instead of returned:
  - $\text{return } a \Rightarrow k \ a$
  - Assuming  $a$  is a constant or variable.
  - “Simple” = “No available function calls.”

2/8/23

34

## CPS Transformation

- Step 3: Pass the current continuation to every function call in tail position
  - $\text{return } f \text{ arg} \Rightarrow f \text{ arg } k$
  - The function “isn’t going to return,” so we need to tell it where to put the result.

2/8/23

35

## CPS Transformation

- Step 4: Each function call not in tail position needs to be converted to take a new continuation (containing the old continuation as appropriate)
  - $\text{return } \text{op } (f \text{ arg}) \Rightarrow f \text{ arg } (\text{fun } r \rightarrow k(\text{op } r))$
  - $\text{op}$  represents a primitive operation
- $\text{return } g(f \text{ arg}) \Rightarrow f \text{ arg } (\text{fun } r \rightarrow g \ r \ k)$

2/8/23

36

## Example

**Before:**  
let rec mem (y,lst) =  
 match lst with  
 [ ] -> false  
 | x :: xs ->  
 if (x = y)  
 then true  
 else mem(y,xs);;

**After:**  
let rec memk (y,lst) k =  
 (\* rule 1 \*)

2/8/23

38

## Example

**Before:**  
let rec mem (y,lst) =  
 match lst with  
 [ ] -> false  
 | x :: xs ->  
 if (x = y)  
 then true  
 else mem(y,xs);;

**After:**  
let rec memk (y,lst) k =  
 (\* rule 1 \*)  
 [ ] -> false k false (\* rule 2 \*)  
 | x :: xs ->  
 if (x = y)  
 then true k true (\* rule 2 \*)  
 else mem(y,xs);;

2/8/23

39

## Example

### Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

### After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  k true (* rule 2 *)  
  memk (y, xs) k (* rule 3 *)
```

2/8/23

40

## Example

### Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

### After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  eqk (x, y)  
  (fun b -> b (* rule 4 *)  
   k true (* rule 2 *)  
   memk (y, xs) (* rule 3 *))
```

2/8/23

41

## Example

### Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

### After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  k false (* rule 2 *)  
  eqk (x, y)  
  (fun b -> if b (* rule 4 *)  
   then k true (* rule 2 *)  
   else memk (y, xs) (* rule 3 *))
```

2/8/23

42

## Example

### Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

### After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  match lst with  
  | [ ] -> k false (* rule 2 *)  
  | x :: xs ->  
    eqk (x, y)  
    (fun b -> if b (* rule 4 *)  
     then k true (* rule 2 *)  
     else memk (y, xs) k (* rule 3 *))
```

2/8/23

43

## Example

### Before:

```
let rec mem (y,lst) =  
  match lst with  
  [ ] -> false  
  | x :: xs ->  
    if (x = y)  
    then true  
    else mem(y,xs);;
```

### After:

```
let rec memk (y,lst) k =  
  (* rule 1 *)  
  match lst with  
  | [ ] -> k false (* rule 2 *)  
  | x :: xs ->  
    eqk (x, y)  
    (fun b -> if b (* rule 4 *)  
     then k true (* rule 2 *)  
     else memk (y, xs) k (* rule 3 *))
```

2/8/23

44

## Example

### Before:

```
let rec add_list lst =  
  match lst with  
  [ ] -> 0  
  | 0 :: xs -> add_list xs  
  | x :: xs -> (+) x  
    (add_list xs);;
```

### After:

```
let rec add_listk lst k =  
  (* rule 1 *)  
  match lst with  
  | [ ] -> k 0 (* rule 2 *)  
  | 0 :: xs -> add_listk xs k  
    (* rule 3 *)  
  | x :: xs -> add_listk xs  
    (fun r -> k ((+) x r));  
  (* rule 4 *)
```

2/8/23

45