Programming Languages and Compilers (CS 421)

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Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

Save the Environment!

A closure is a pair of an environment and an association of a pattern (e.g. (v1,...,vn) giving the input variables) with an expression (the function body), written:

< (v1,...,vn) \rightarrow exp, ρ >

 Where p is the environment in effect when the function is defined (for a simple function)

Evaluating declarations

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 - Evaluate expression e in p to value v
 - Update ρ with $x \rightarrow v$: $\{x \rightarrow v\} + \rho$

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- Update: ρ₁+ ρ₂ has all the bindings in ρ₁ and all those in ρ₂ that are not rebound in ρ₁
 {x → 2, y → 3, a → "hi"} + {y → 100, b → 6}
 {x → 2, y → 3, a → "hi", b → 6}

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- To evaluate a tuple (e₁,...,e_n),
 - Evaluate each e_i to v_i, right to left for Ocaml
 - Then make value (v₁,...,v_n)

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- To evaluate a local dec: let x = e1 in e2
 Eval e1 to v, then eval e2 using {x → v} + ρ
- To evaluate a conditional expression: if b then e1 else e2
 - Evaluate b to a value v
 - If v is True, evaluate e1
 - If v is False, evaluate e2

Evaluation of Application with Closures

- Given application expression f e
- In Ocaml, evaluate e to value v
- In environment ρ , evaluate left term to closure, c = <(x₁,...,x_n) → b, ρ' >
 - (x₁,...,x_n) variables in (first) argument
 - v must have form (v₁,...,v_n)
- Update the environment ρ' to
 - $\rho'' = \{\mathbf{x}_1 \rightarrow \mathbf{v}_1, \dots, \mathbf{x}_n \rightarrow \mathbf{v}_n\} + \rho'$

Evaluate body **b** in environment ρ''

Recursive Functions

- # let rec factorial n =
 if n = 0 then 1 else n * factorial (n 1);;
 val factorial : int -> int = <fun>
 # factorial 5;;
 : int = 120
- # (* rec is needed for recursive function
 declarations *)

Recursion Example

Compute n^2 recursively using: $n^2 = (2 * n - 1) + (n - 1)^2$	
<pre># let rec nthsq n = (match n (* pat with 0 -> 0</pre>	<pre>* rec for recursion *) tern matching for cases *) (* base case *) (* recursive case *) * recursive call *)</pre>

Structure of recursion similar to inductive proof

Recursion and Induction

let rec nthsq n = match n with 0 -> 0 | n -> (2 * n - 1) + nthsq (n - 1) ;;

- Base case is the last case; it stops the computation
- Recursive call must be to arguments that are somehow smaller - must progress to base case
- if or match must contain base case
- Failure of these may cause failure of termination



List can take one of two forms:

- Empty list, written []
- Non-empty list, written x :: xs
 - x is head element, xs is tail list, :: called "cons"
- Syntactic sugar: [x] == x :: []
- [x1; x2; ...; xn] == x1 :: x2 :: ... :: xn :: []



let fib5 = [8;5;3;2;1;1];;val fib5 : int list = [8; 5; 3; 2; 1; 1]# let fib6 = 13 :: fib5;; val fib6 : int list = [13; 8; 5; 3; 2; 1; 1]# (8::5::3::2::1::1::[]) = fib5;;-: bool = true # fib5 @ fib6;; - : int list = [8; 5; 3; 2; 1; 1; 13; 8; 5; 3; 2; 1;1]

Lists are Homogeneous

This expression has type float but is here used with type int



- Which one of these lists is invalid?
- **1**. [2; 3; 4; 6]
- 2. [2,3; 4,5; 6,7]
- **3**. [(2.3,4); (3.2,5); (6,7.2)]
- 4. [["hi"; "there"]; ["wahcha"]; []; ["doin"]]



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- **3**. [(2.3,4); (3.2,5); (6,7.2)]
- 4. [["hi"; "there"]; ["wahcha"]; []; ["doin"]]
- 3 is invalid because of last pair

Functions Over Lists

let rec double up list = match list with $[] \rightarrow []$ (* pattern before ->, expression after *) (x :: xs) -> (x :: x :: double_up xs);; val double_up : 'a list -> 'a list = <fun> # let fib5 2 =double up fib5;; val fib5 2 : int list = [8; 8; 5; 5; 3; 3; 2; 2; 1;1; 1; 1]

Functions Over Lists

- # let silly = double_up ["hi"; "there"];; val silly : string list = ["hi"; "hi"; "there"; "there"] # let rec poor rev list = match list with [] -> [] | (x::xs) -> poor_rev xs @ [x];; val poor_rev : 'a list -> 'a list = <fun> # poor_rev silly;;
- : string list = ["there"; "there"; "hi"; "hi"]

Structural Recursion

- Functions on recursive datatypes (eg lists) tend to be recursive
- Recursion over recursive datatypes generally by structural recursion
 - Recursive calls made to components of structure of the same recursive type
 - Base cases of recursive types stop the recursion of the function

Problem: write code for the length of the list

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- let rec length list =

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match list with

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 What patterns should we match against?
 let rec length list =

 match list with

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 What result do we give when list is empty?
 let rec length list =

 match list with [] -> 0
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 match list with [] -> 0
 (a :: bs) -> 1 + length bs

let rec length list = match list with [] -> 0 (* Nil case *) | a :: bs -> 1 + length bs;; (* Cons case *) val length : 'a list -> int = <fun> # length [5; 4; 3; 2];;

- -: int = 4
- Nil case [] is base case

Cons case recurses on component list bs

Same Length

How can we efficiently answer if two lists have the same length?

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How can we efficiently answer if two lists have the same length? let rec same length list1 list2 = match list1 with [] -> (match list2 with [] -> true $|(y::ys) \rightarrow false)$ (x::xs) -> (match list2 with [] -> false (y::ys) -> same_length xs ys)

Your turn: doubleList : int list -> int list

Write a function that takes a list of int and returns a list of the same length, where each element has been multiplied by 2

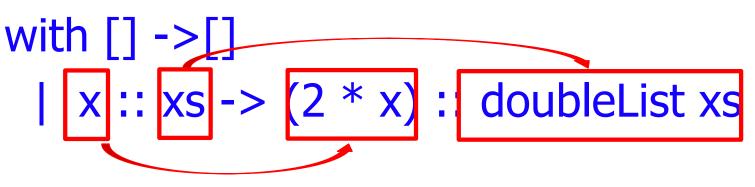
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- let rec doubleList list =
 - match list



Higher-Order Functions Over Lists

let rec map f list = match list with [] -> [] |(h::t) -> (f h) :: (map f t);;val map : ('a -> 'b) -> 'a list -> 'b list = $\langle fun \rangle$ # map plus_two fib5;; - : int list = [10; 7; 5; 4; 3; 3]# map (fun x -> x - 1) fib6;; : int list = [12; 7; 4; 2; 1; 0; 0]

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Mapping Recursion

Can use the higher-order recursive map function instead of direct recursion

let doubleList list =
 List.map (fun x -> 2 * x) list;;
val doubleList : int list -> int list = <fun>
doubleList [2;3;4];;

-: int list = [4; 6; 8]

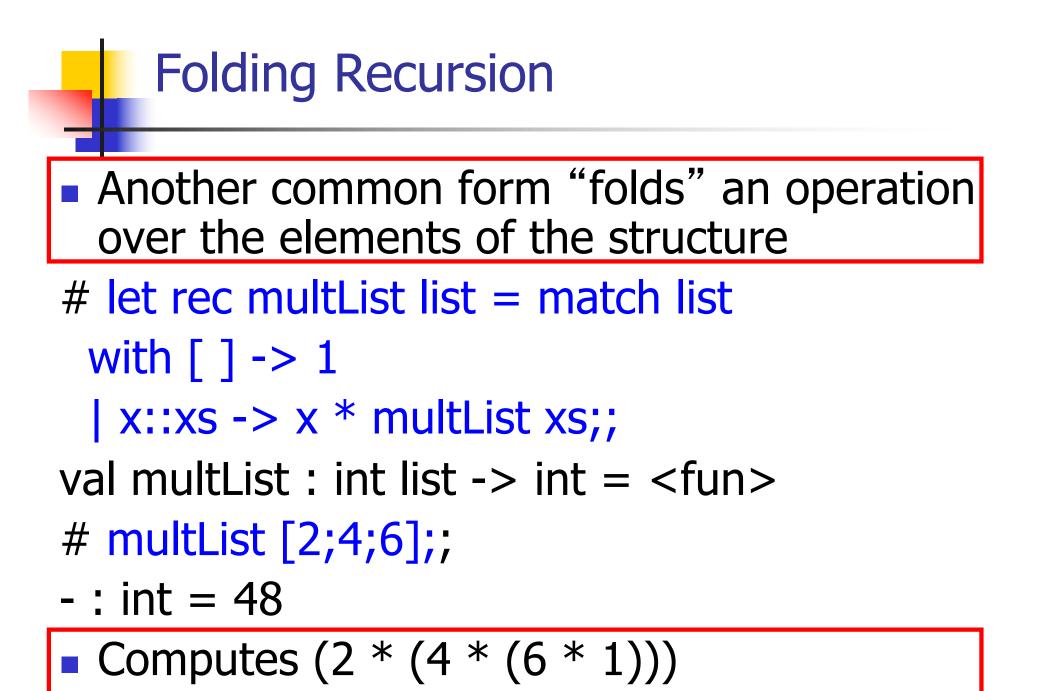
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Same function, but no explicit recursion



Folding Recursion : Length Example

let rec length list = match list with [] -> 0 (* Nil case *) | a :: bs -> 1 + length bs;; (* Cons case *) val length : 'a list -> int = <fun> # length [5; 4; 3; 2];;

- : int = 4
- Nil case [] is base case, 0 is the base value
- Cons case recurses on component list bs
- What do multList and length have in common?