

CS 374 A

Induction/Recursion

Σ = alphabet = $\{0, 1\}$
any finite set

String = either ϵ empty
or $a \cdot x$ symbol $a \in \Sigma$
string $x \in \Sigma^*$

Σ^* = all strings over Σ

All strings are finite, but Σ^* is infinite

Sequence of $\langle \text{foo} \rangle$ s is ϵ
or $a \cdot x$ where a is $\langle \text{foo} \rangle$
 x is seq of $\langle \text{foos} \rangle$

STRING = S . TRING
= S . (T . RING)



length $|w| = \begin{cases} 0 & \text{if } w = \epsilon \\ 1 + |x| & \text{if } w = a \cdot x \end{cases}$

concatenation $w \cdot z = \begin{cases} z & w = \epsilon \\ a \cdot (x \cdot z) & w = a \cdot x \end{cases}$

NOW • HERE $\rightarrow$ NOWHERE

Lemma: For all strings w and z :

$$|w \circ z| = |w| + |z|$$

Proof:

Let w and z be arbitrary strings

Assume for all strings x shorter than w that

There are two cases: $|x \circ z| = |x| + |z|$

$$\begin{aligned} \bullet w = \epsilon & \quad |w \circ z| = |\epsilon \circ z| \\ & = |z| \\ & = 0 + |z| \\ & = |z| + |z| \\ & = |w| + |z| \end{aligned}$$

$w = \epsilon$
def $\circ$
duh
def $|z|$
def w

$$\bullet w = a \cdot x$$

$$\begin{aligned} |w \circ z| & = |a \cdot x \circ z| \\ & = |a(x \circ z)| \\ & = 1 + |\underline{x} \circ z| \\ & = 1 + |\underline{x}| + |z| \\ & = |ax| + |z| \\ & = |w| + |z| \end{aligned}$$

def w
def $\circ$
def $|$
IH!
def $|$
def w

Therefore, $|w \circ z| = |w| + |z|$.