

# CS374 A

## Induction/Recursion

$\Sigma = \text{alphabet} = \{\text{0, 1}\}$   
any finite set

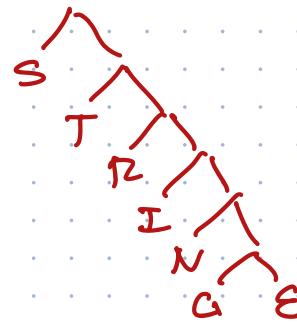
String = either  $\epsilon$  empty  
or  $a \cdot x$  symbol  $a \in \Sigma$   
string  $x \in \Sigma^*$

$\Sigma^*$  = all strings over  $\Sigma$

All strings are finite, but  $\Sigma^*$  is infinite

Sequence of  $\langle \text{FOO} \rangle$ s is  $\epsilon$   
or  $a \cdot x$  where  $a$  is  $\langle \text{FOO} \rangle$   
 $x$  is seq of  $\langle \text{FOO} \rangle$

STRING = S · T R I N G  
= S · (T · R I N G)



length  $|w| = \begin{cases} 0 & \text{if } w = \epsilon \\ 1 + |x| & \text{if } w = a \cdot x \end{cases}$

concatenation  $w \cdot z = \begin{cases} z & w = \epsilon \\ a \cdot (x \cdot z) & w = a \cdot x \end{cases}$

NOW · HERE  $\rightarrow$  NOWHERE

Lemma: For all strings  $w$  and  $z$ :

$$|w \circ z| = |w| + |z|$$

Proof:

Let  $w$  and  $z$  be arbitrary strings

Assume for all strings  $x$  shorter than  $w$  that

There are two cases:

$$\begin{aligned} \bullet w &= \epsilon & |w \circ z| &= |\epsilon \circ z| \\ &&&= |z| \\ &&&= 0 + |z| \\ &&&= |\epsilon| + |z| \\ &&&= |w| + |z| \end{aligned}$$

$w = \epsilon$   
def  $\bullet$   
duh  
def  $|c|$   
def  $w$

$$\bullet w = \alpha \cdot x$$

$$\begin{aligned} |w \circ z| &= |\alpha x \circ z| \\ &= |\alpha(x \circ z)| \\ &= 1 + |\underline{x \circ z}| \\ &= 1 + |\underline{x}| + |\underline{z}| \\ &= |\alpha x| + |z| \\ &= |w| + |z| \end{aligned}$$

def  $w$   
def  $\circ$   
def  $|1|$   
 $\text{IH!}$   
def  $|1|$   
def  $w$

Therefore,  $|w \circ z| = |w| + |z|$ .