CS/ECE 374 A Lab 1b — August 29 Fall 2025

Recall from lecture that a *regular expression* is compact notation for a language (that is, a set of strings). Formally, a regular language is one of the following:

- The symbol Ø (representing the empty set)
- Any string (representing the set containing only that string)
- R + S for some regular expressions R and S (representing alternation / union)
- $R \cdot S$  or RS for some regular expressions R and S (representing concatenation)
- R\* for some regular expression R (representing Kleene closure / unbounded repetition)

In the absence of parentheses, Kleene closure has highest precedence, followed by concatenation. For example:

Give regular expressions for each of the following languages over the binary alphabet  $\{0, 1\}$ . (For extra practice, find multiple regular expressions for each language.)

- o. All strings.
- 1. All strings containing the substring 000.
- 2. All strings *not* containing the substring 000.
- 3. All strings in which every run of 0s has length at least 3.
- 4. All strings in which the *last* 1 appears before the *first* substring 000.
- 5. All strings containing at least three 0s.
- 6. Every string except 000. [Hint: Don't try to be clever.]

## More difficult problems to work on later:

- 7. All strings w such that in every prefix of w, the number of 0s and 1s differ by at most 1.
- 8. All strings containing at least two 0s and at least one 1.
- 9. All strings in which every run has odd length. (For example, 0001 and 100000111 and the empty string  $\varepsilon$  are in this language, but 000000 and 001000 are not.)
- ★10. All strings w such that in every prefix of w, the number of 0s and 1s differ by at most 2.
- \*11. All strings in which the substring 000 appears an even number of times. (For example, 01100 and 000000 and the empty string  $\varepsilon$  are in this language, but 00000 and 001000 are not.)