

CS/ECE 374 A ✧ Fall 2023
☞ Practice Midterm 2 ☞
November 2, 2023

Name:	
NetID:	

-
- ***Don't panic!***
 - You have 120 minutes to answer five questions. The questions are described in more detail in a separate handout.
 - If you brought anything except your writing implements, your **hand-written** double-sided 8½" × 11" cheat sheet, and your university ID, please put it away for the duration of the exam. In particular, please turn off and put away *all* medically unnecessary electronic devices.
 - Please clearly print your name and your NetID in the boxes above.
 - Please also print your name at the top of every page of the answer booklet, except this cover page. We want to make sure that if a staple falls out, we can reassemble your answer booklet. (It doesn't happen often, but it does happen.)
 - **Do not write outside the black boxes on each page.** These indicate the area of the page that our scanner can actually see. Anything you write outside the boxes will be erased before we start grading.
 - If you run out of space for an answer, please use the scratch pages at the back of the answer booklet, but **please clearly indicate where we should look.** Please ask for more scratch paper if you need it.
 - Proofs or other justifications are required for full credit if and only if we explicitly ask for them, using the word ***prove*** or ***justify*** in bold italics.
 - Please return ***all*** paper with your answer booklet: your question sheet, your cheat sheet, and all scratch paper. **Please put all loose paper *inside* your answer booklet.**
-

- (a) Write the solution to each of the following recurrences in the box immediately below it. (Use the space below the boxes for scratch work.)

$$A(n) = 3A(n/2) + O(n^2)$$

$$B(n) = 7B(n/2) + O(n^2)$$

$$C(n) = 4C(n/2) + O(n^2)$$

- (b) Draw a directed acyclic graph with at most ten vertices, exactly one source, exactly one sink, and more than one topological order.
- (c) Draw a directed graph with at most ten vertices, with distinct positive edge weights, that has more than one shortest path from some vertex s to some other vertex t .
- (d) Describe an appropriate memoization structure and evaluation order for the following (meaningless) recurrence, and give the running time of the resulting iterative algorithm to compute $Huh(1, n)$.

$$Huh(i, k) = \begin{cases} 0 & \text{if } i > n \text{ or } k < 0 \\ \min \left\{ \begin{array}{l} Huh(i+1, k-2) \\ Huh(i+2, k-1) \end{array} \right\} + A[i, k] & \text{if } A[i, k] \text{ is even} \\ \max \left\{ \begin{array}{l} Huh(i+1, k-2) \\ Huh(i+2, k-1) \end{array} \right\} - A[i, k] & \text{if } A[i, k] \text{ is odd} \end{cases}$$

See the question sheet for a detailed description of your game with Elmo.

- (a) **Prove** that you should not also use the greedy strategy. That is, show that there is a game that you can win, but only if you do *not* follow the same greedy strategy as Elmo. Assume Elmo plays first.
 - (b) Describe and analyze an algorithm to determine, given the initial sequence of cards, the maximum number of points that you can collect playing against Elmo.
-

Suppose you are given a directed graph $G = (V, E)$, whose vertices are either red, green, or blue. Edges in G do not have weights, and G is not necessarily a dag. The *remoteness* of a vertex v is the *maximum* of three shortest-path lengths:

- The length of a shortest path to v from the closest red vertex
- The length of a shortest path to v from the closest blue vertex
- The length of a shortest path to v from the closest green vertex

In particular, if v is not reachable from vertices of all three colors, then v is infinitely remote. Describe and analyze an algorithm to find a vertex of G with *minimum* remoteness.

Suppose you are given an array $A[1..n]$ of integers such that $A[i] + A[i + 1]$ is even for *exactly one* index i . In other words, the elements of A alternate between even and odd, except for exactly one adjacent pair that are either both even or both odd. Describe and analyze an efficient algorithm to find the unique index i such that $A[i] + A[i + 1]$ is even.

A *zigzag walk* in a directed graph G is a sequence of vertices connected by edges in G , but the edges alternately point forward and backward along the sequence. Specifically, the first edge points forward, the second edge points backward, and so on. The *length* of a zigzag walk is the sum of the weights of its edges, both forward and backward.

Suppose you are given a directed graph G with non-negatively weighted edges, along with two vertices s and t . Describe and analyze an algorithm to find the shortest zigzag walk from s to t in G .

(scratch paper)

(scratch paper)

(scratch paper)

(scratch paper)

(scratch paper)