You saw the following context-free grammars in class on Thursday; in each example, the grammar itself is on the left; the explanation for each non-terminal is on the right.

• Properly nested strings of parentheses.

$$S \to \varepsilon \mid S(S)$$
 properly nested parentheses

Here is a different grammar for the same language:

$$S \rightarrow \varepsilon \mid (S) \mid SS$$
 properly nested parentheses

•  $\{0^m1^n \mid m \neq n\}$ . This is the set of all binary strings composed of some number of 0s followed by a *different* number of 1s.

$$S \to A \mid B$$
 {0<sup>m</sup>1<sup>n</sup> | m \neq n}  
 $A \to 0A \mid 0C$  {0<sup>m</sup>1<sup>n</sup> | m > n}  
 $B \to B1 \mid C1$  {0<sup>m</sup>1<sup>n</sup> | m < n}  
 $C \to \varepsilon \mid 0C1$  {0<sup>m</sup>1<sup>n</sup> | m = n}

Give context-free grammars for each of the following languages over the alphabet  $\Sigma = \{0,1\}$ . For each grammar, describe the language for each non-terminal, either in English or using mathematical notation, as in the examples above. We probably won't finish all of these during the lab session.

- 1. All palindromes in  $\Sigma^*$
- 2. All palindromes in  $\Sigma^*$  that contain an even number of 1s
- 3. All palindromes in  $\Sigma^*$  that end with 1
- 4. All palindromes in  $\Sigma^{\ast}$  whose length is divisible by 3
- 5. All palindromes in  $\Sigma^*$  that do not contain the substring 00

## Harder problems to work on later:

6. 
$$\{0^{2n}1^n \mid n \ge 0\}$$

7. 
$$\{0^m 1^n \mid m \neq 2n\}$$

[Hint: If  $m \neq 2n$ , then either m < 2n or m > 2n. Extend the previous grammar, but pay attention to parity. This language contains the string 01.]

8. 
$$\{0,1\}^* \setminus \{0^{2n}1^n \mid n \ge 0\}$$

[Hint: Extend the previous grammar. What's missing?]

- 9.  $\{w \in \{0,1\}^* \mid \#(0,w) = 2 \cdot \#(1,w)\}$  Binary strings where the number of 0s is exactly twice the number of 1s.
- \*10.  $\{0,1\}^* \setminus \{ww \mid w \in \{0,1\}^*\}.$

[Anti-hint: The language  $\{ww \mid w \in 0, 1^*\}$  is **not** context-free. Thus, the complement of a context-free language is not necessarily context-free!]