## Algorithms \& Models of Computation <br> CS/ECE 374, Fall 2020 <br> 17.4 <br> Shortest path trees and variants

## Algorithms \& Models of Computation

## CS/ECE 374, Fall 2020

17.4.1

Shortest Path Tree

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\(Q=\) make \(P Q()\)
insert ( \(Q,(s, 0)\) )
\(\operatorname{prev}(s) \leftarrow\) null
for each node \(\boldsymbol{u} \neq \boldsymbol{s}\) do
        insert \((\boldsymbol{Q},(\boldsymbol{u}, \infty))\)
        \(\operatorname{prev}(\boldsymbol{u}) \leftarrow\) null
\(\boldsymbol{X}=\emptyset\)
for \(\boldsymbol{i}=\mathbf{1}\) to \(|\boldsymbol{V}|\) do
    \((v, \operatorname{dist}(s, v))=\) extractMin \((Q)\)
    \(\boldsymbol{X}=\boldsymbol{X} \cup\{v\}\)
    for each \(u\) in \(\operatorname{Adj}(v)\) do
            if \((\operatorname{dist}(s, v)+\ell(v, u)<\operatorname{dist}(s, u))\) then
                \(\operatorname{decreaseKey}(\boldsymbol{Q},(\boldsymbol{u}, \operatorname{dist}(\boldsymbol{s}, \boldsymbol{v})+\ell(\boldsymbol{v}, \boldsymbol{u})))\)
                \(\operatorname{prev}(u)=\boldsymbol{v}\)
```


## Shortest Path Tree

## Lemma

The edge set $(\boldsymbol{u}, \operatorname{prev}(\boldsymbol{u}))$ is the reverse of a shortest path tree rooted at $\boldsymbol{s}$. For each $\boldsymbol{u}$, the reverse of the path from $\boldsymbol{u}$ to $\boldsymbol{s}$ in the tree is a shortest path from $\boldsymbol{s}$ to $\boldsymbol{u}$.

## Proof Sketch.

(1) The edge set $\{(\boldsymbol{u}, \operatorname{prev}(\boldsymbol{u})) \mid \boldsymbol{u} \in \boldsymbol{V}\}$ induces a directed in-tree rooted at $\boldsymbol{s}$ (Why?)
(2) Use induction on $|\boldsymbol{X}|$ to argue that the tree is a shortest path tree for nodes in $\boldsymbol{V}$.

## Shortest paths to $\mathbf{s}$

Dijkstra's algorithm gives shortest paths from $\boldsymbol{s}$ to all nodes in $\boldsymbol{V}$. How do we find shortest paths from all of $\boldsymbol{V}$ to $\boldsymbol{s}$ ?
(1) In undirected graphs shortest path from $s$ to $u$ is a shortest path from $u$ to $s$ so
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## THE END

(for now)

