Algorithms \& Models of Computation
CS/ECE 374, Fall 2020
14.2

Edit Distance and Sequence Alignment

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14.2.1

Problem definition and background

## Spell Checking Problem

Given a string "exponen" that is not in the dictionary, how should a spell checker suggest a nearby string?

What does nearness mean?

Question: Given two strings $x_{1} x_{2} \ldots x_{n}$ and $y_{1} y_{2} \ldots y_{m}$ what is a distance between them?

Edit Distance: minimum number of "edits" to transform $\boldsymbol{x}$ into $\boldsymbol{y}$

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## Edit Distance

## Definition 14.1.

Edit distance between two words $\boldsymbol{X}$ and $\boldsymbol{Y}$ is the number of letter insertions, letter deletions and letter substitutions required to obtain $\boldsymbol{Y}$ from $\boldsymbol{X}$.

## Example 14.2.

The edit distance between FOOD and MONEY is at most 4:

$$
\underline{\text { FOOD }} \rightarrow \text { MOODD } \rightarrow \text { MONOD } \rightarrow \text { MONED } \rightarrow \text { MONEY }
$$

## Edit Distance: Alternate View

## Alignment

Place words one on top of the other, with gaps in the first word indicating insertions, and gaps in the second word indicating deletions.

| F | $\mathbf{O}$ | $\mathbf{O}$ |  | D |
| :---: | :---: | :---: | :---: | :---: |
| M | $\mathbf{O}$ | $\mathbf{N}$ | E | $\mathbf{Y}$ |

Formally, an alignment is a set $\boldsymbol{M}$ of pairs $(\boldsymbol{i}, \boldsymbol{j})$ such that each index appears at most once, and there is no "crossing": $\boldsymbol{i}<\boldsymbol{i}^{\prime}$ and $\boldsymbol{i}$ is matched to $\boldsymbol{j}$ implies $\boldsymbol{i}^{\prime}$ is matched to $j^{\prime}>j$. In the above example, this is $M=\{(1,1),(2,2),(3,3),(4,5)\}$. Cost of an alignment is the number of mismatched columns plus number of unmatched indices in both strings.

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## Edit Distance Problem

## Problem

Given two words, find the edit distance between them, i.e., an alignment of smallest cost.

## Applications

(1) Spell-checkers and Dictionaries
(2) Unix diff
(3) DNA sequence alignment . . . but, we need a new metric

## Similarity Metric

## Definition 14.3.

For two strings $\boldsymbol{X}$ and $\boldsymbol{Y}$, the cost of alignment $\boldsymbol{M}$ is
(1) [Gap penalty] For each gap in the alignment, we incur a cost $\boldsymbol{\delta}$.
(2) [Mismatch cost] For each pair $\boldsymbol{p}$ and $\boldsymbol{q}$ that have been matched in $\boldsymbol{M}$, we incur cost $\alpha_{p q}$; typically $\alpha_{p p}=\mathbf{0}$.

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Edit distance is special case when $\delta=\alpha_{p q}=\mathbf{1}$.

## THE END

(for now)

