Algorithms \& Models of Computation

## 7.8

Supplemental: Why $a^{n} b^{n} c^{n}$ is not CFL

## You are bound to repeat yourself...

$L=\left\{a^{n} b^{n} c^{n} \mid n \geq 0\right\}$.
(1) For the sake of contradiction assume that there exists a grammar: $G$ a CFG for $L$.
(2) $T_{i}$ : minimal parse tree in $G$ for $a^{i} b^{i} \boldsymbol{c}^{i}$.
© $h_{i}=\operatorname{height}\left(T_{i}\right)$ : Length of longest path from root to leaf in $T_{i}$

- For any integer $t$, there must exist an index $j(t)$, such that $h_{j(t)}>t$
- There an index $\boldsymbol{i}$, such that $\boldsymbol{h}_{\boldsymbol{j}}>(2 * \#$ variables in $\boldsymbol{G})$.


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(1) For any integer $\boldsymbol{t}$, there must exist an index $\boldsymbol{j}(\boldsymbol{t})$, such that $\boldsymbol{h}_{j(t)}>\boldsymbol{t}$.

- There an index $\boldsymbol{j}$, such that $\boldsymbol{h}_{\boldsymbol{j}}>(2 * \#$ variables in $\boldsymbol{G})$.


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$x y z v w=a^{j} b^{j} c^{j} \Longrightarrow x y^{2} z v^{2} w \in L$

- We know:
$x y z v w=a^{j} b^{j} c^{j}$
$|y|+|v|>0$.
- We proved that $\boldsymbol{\tau}=x y^{2} z \boldsymbol{v}^{2} w \in L$.
- If $y$ contains both $a$ and $b$, then, $\tau=\ldots . . . . . b . . a . . . b . .$.

Impossible, since $\tau \in L=\left\{a^{n} b^{n} c^{n} \mid n \geq 0\right\}$

- Similarly, not possible that $\boldsymbol{y}$ contains both $\boldsymbol{b}$ and c
- Similarly, not possible that $v$ contains both $a$ and $b$
- Similarly, not possible that $v$ contains both $b$ and $c$
- If $\boldsymbol{y}$ contains only as, and $\boldsymbol{v}$ contains only $\boldsymbol{b s}$, then... \#(a) $(\tau) \neq \#_{(c)}(\tau)$. Not possible.
- Similarly, not possible that $y$ contains only as, and $v$ contains only cs. Similarly, not possible that $y$ contains only $b s$, and $v$ contains only $c s$.
- Must be that $\tau \notin L$. A contradiction.
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- Similarly, not possible that $\boldsymbol{y}$ contains both $b$ and $c$
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- If $\boldsymbol{y}$ contains only $a s$, and $v$ contains only $b s$, then... $\#_{(a)}(\tau) \neq \#_{(c)}(\tau)$ Not possible
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## We conclude...

## Lemma

The language $L=\left\{a^{n} b^{n} c^{n} \mid n \geq 0\right\}$ is not CFL (i.e., there is no CFG for it).

