# Non-deterministic Finite Automata (NFAs)

Lecture 4

Thursday, September 3, 2020

LATEXed: July 22, 2020 21:54

### Algorithms & Models of Computation CS/ECE 374, Fall 2020

# **4.1** NFA Introduction

#### Non-deterministic Finite State Automata by example When you come to a fork in the road, take it.



### Non-deterministic Finite State Automata by example II ..but only if it is made out of silver.



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### Non-deterministic Finite State Automata (NFAs)



#### Differences from DFA

- From state q on same letter  $a \in \Sigma$  multiple possible states
- No transitions from *q* on some letters
- $\varepsilon$ -transitions!

Questions:

- Is this a "real" machine?
- What does it do?

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- From  $q_{\varepsilon}$  on 1
- From  $q_{\epsilon}$  on 0
- From  $q_0$  on  $\varepsilon$
- From  $q_{\varepsilon}$  on 01
- From *q*<sub>00</sub> on 00



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#### NFA acceptance: informal



**Informal definition:** An NFA N accepts a string w iff some accepting state is reached by N from the start state on input w.

The language accepted (or recognized) by a NFA N is denote by L(N) and defined as:  $L(N) = \{w \mid N \text{ accepts } w\}.$ 

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#### • Is 01 accepted?

- Is 001 accepted?
- Is 100 accepted?
- Are all strings in 1\*01 accepted?
- What is the language accepted by **N**?



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#### ababa.

Idea: Keep track of the states where the NFA might be at any given time.



Remaining input: *ababa*.



Remaining input: **baba**.



Remaining input: **baba**.



Remaining input: *aba*.



Remaining input: *aba*.



Remaining input: ba.



Remaining input: **ba**.



Remaining input: a.



Remaining input: a.

### Simulating NFAExample the first



Remaining input: *c*.



Remaining input:  $\varepsilon$ .

Accepts: *ababa*.

#### An exercise For you to think about...

A. What is the language that the following NFA accepts?



B. What is the minimal number of states in a DFA that recognizes the same language?

### THE END

### (for now)

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