Algorithms & Models of Computation CS/ECE 374, Fall 2020

## 3.1.1

## Graphical representation of $\mathrm{DFA}$

### Graphical Representation/State Machine



- Directed graph with nodes representing states and edge/arcs representing transitions labeled by symbols in  $\Sigma$
- For each state (vertex) *q* and symbol *a* ∈ Σ there is <u>exactly</u> one outgoing edge labeled by *a*
- Initial/start state has a pointer (or labeled as *s*, *q*<sub>0</sub> or "start")
- Some states with double circles labeled as accepting/final states Har-Peled (UUC) CS374 9



#### • Where does 001 lead?

- Where does 10010 lead?
- Which strings end up in accepting state?
- Can you prove it?
- Every string w has a unique walk that it follows from a given state q by reading one letter of w from left to right.

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#### Definition

A DFA M accepts a string w iff the unique walk starting at the start state and spelling out w ends in an accepting state.





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It means that M accepts each string in L and no others. Equivalently M accepts each string in L and does not accept/rejects strings in  $\Sigma^* \setminus L$ .

*M* "recognizes" *L* is a better term but "accepts" is widely accepted (and recognized) (joke attributed to Lenny Pitt)

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## THE END

# (for now)

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