# Strings and Languages 

Lecture 1
Tuesday, August 25, 2020

## Algorithms \& Models of Computation

## 1.1 <br> Strings

## Alphabet

An alphabet is a finite set of symbols.
Examples of alphabets:

- $\Sigma=\{0,1\}$,
- $\Sigma=\{a, b, c, \ldots, z\}$,
- ASCII.
- UTF8
- $\Sigma=\{\langle$ moveforward $\rangle,\langle$ moveback $\rangle\}$


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## String Definitions

## Definition

（1）A string／word over $\Sigma$ is a finite sequence of symbols over $\Sigma$ ．For example， ＇0101001＇，＇string＇，＇〈moveback〉〈rotate90〉＇
（2）$\epsilon$ is the empty string．
（0）The length of a string $w$（denoted by $|w|$ ）is the number of symbols in $w$ ．For example，$|101|=3,|\epsilon|=0$
－For integer $\boldsymbol{n} \geq 0, \Sigma^{\boldsymbol{n}}$ is set of all strings over $\Sigma$ of length $\boldsymbol{n}$ ．$\Sigma^{*}$ is the set of all strings over $\Sigma$ ．

## Inductive/recursive definition of strings

Formal definition of a string:

- $\epsilon$ is a string of length 0
- $a x$ is a string if $a \in \Sigma$ and $x$ is a string. The length of $a x$ is $1+|x|$

The above definition helps prove statements rigorously via induction.

- Alternative recursive definition useful in some proofs: $x a$ is a string if $a \in \Sigma$ and $x$ is a string. The length of $x a$ is $1+|x|$


## Convention

- $a, b, c, \ldots$ denote elements of $\Sigma$
- $w, x, y, z, \ldots$ denote strings
- $A, B, C, \ldots$ denote sets of strings


## Much ado about nothing

- $\epsilon$ is a string containing no symbols. It is not a set
- $\{\epsilon\}$ is a set containing one string: the empty string. It is a set, not a string.
- $\emptyset$ is the empty set. It contains no strings.
- $\{\emptyset\}$ is a set containing one element, which itself is a set that contains no elements.


## Concatenation and properties

- If $x$ and $y$ are strings then $x y$ denotes their concatenation.
- concatenation defined recursively :
- $x y=y$ if $x=\epsilon$
- $x y=a(w y)$ if $x=a w$
- xy sometimes written as $x \bullet y$
- concatenation is associative: $(u v) w=u(v w)$ hence write $u v w \equiv(u v) w=u(v w)$
- not commutative: uv not necessarily equal to vu
- The identity element is the empty string $\epsilon$


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$$
\boldsymbol{\epsilon} \boldsymbol{u}=\boldsymbol{u} \boldsymbol{\epsilon}=\boldsymbol{u}
$$

## Substrings, prefix, suffix

## Definition

$v$ is substring of $w \Longleftrightarrow$ there exist strings $x, y$ such that $w=x v y$.

- If $\boldsymbol{x}=\boldsymbol{\epsilon}$ then $\boldsymbol{v}$ is a prefix of $\boldsymbol{w}$
- If $\boldsymbol{y}=\boldsymbol{\epsilon}$ then $\boldsymbol{v}$ is a suffix of $\boldsymbol{w}$


## String exponents

## Definition

If $w$ is a string then $w^{\boldsymbol{n}}$ is defined inductively as follows:
$w^{\boldsymbol{n}}=\boldsymbol{\epsilon}$ if $\boldsymbol{n}=\mathbf{0}$
$w^{n}=w w^{n-1}$ if $n>0$

Example: $(\text { blah })^{4}=$ blahblahblahblah.

## Set Concatenation

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Given two sets $\boldsymbol{X}$ and $\boldsymbol{Y}$ of strings (over some common alphabet $\Sigma$ ) the concatenation of $\boldsymbol{X}$ and $\boldsymbol{Y}$ is

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X Y=\{x y \mid x \in X, y \in Y\}
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## Example

```
X = {fido, rover, spot },
Y ={fluffy, tabby }
XY = { fidofluffy, fidotabby, roverfluffy, ...}.
```


## $\sum^{*}$ and languages

## Definition

(1) $\Sigma^{n}$ is the set of all strings of length $n$. Defined inductively:
$\Sigma^{n}=\{\epsilon\}$ if $\boldsymbol{n}=0$
$\Sigma^{n}=\Sigma \Sigma^{n-1}$ if $n>0$
(2) $\Sigma^{*}=\cup_{n \geq 0} \Sigma^{n}$ is the set of all finite length strings
( $\Sigma^{+}=\cup_{n \geq 1} \Sigma^{n}$ is the set of non-empty strings.

## Definition

A language $L$ is a set of strings over $\Sigma$. In other words $L \subseteq \sum^{*}$.

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## Exercise

Answer the following questions taking $\Sigma=\{0,1\}$.
(1) What is $\Sigma^{0}$ ?
(2) How many elements are there in $\Sigma^{3}$ ?
(3) How many elements are there in $\sum^{n}$ ?
(4) What is the length of the longest string in $\Sigma$ ?
(5) Does $\Sigma^{*}$ have strings of infinite length?
(6) If $|\boldsymbol{u}|=2$ and $|\boldsymbol{v}|=3$ then what is $|\boldsymbol{u} \bullet \boldsymbol{v}|$ ?
(3) Let $\boldsymbol{u}$ be an arbitrary string in $\Sigma^{*}$. What is $\boldsymbol{\epsilon} \boldsymbol{u}$ ? What is $\boldsymbol{u} \boldsymbol{\epsilon}$ ?
(8) Is $u \boldsymbol{v}=\boldsymbol{v u}$ for every $u, v \in \Sigma^{*}$ ?
(0) Is $(u v) \boldsymbol{w}=\boldsymbol{u}(\boldsymbol{v} \boldsymbol{w})$ for every $u, v, w \in \Sigma^{*}$ ?

## THE END

## (for now)

