

# Poisson Example: Genome Assembly

- **Goal:** DNA sequence (ACTG) of the entire genome
- **Problem:** Sequencers generate random short reads

Sequencer	Sanger 3730xl	454 GS	Ion Torrent	SOLiDv4	Illumina HiSeq 2000	Pac Bio
Mechanism	Dideoxy chain termination	Pyrosequencing	Detection of hydrogen ion	Ligation and two-base coding	Reversible Nucleotides	Single molecule real time
Read length	400-900 bp	700 bp	~400 bp	50 + 50 bp	100 bp PE	>10000 bp
Error Rate	0.001%	0.1%	2%	0.1%	2%	10-15%
Output data (per run)	100 KB	1 GB	100 GB	100 GB	1 TB	10 GB
Approx cost per GB		10,000	1000	100	10	1000

- **Solution:** assemble genome from short reads using computers. Whole Genome Shotgun Assembly.

Table from the course EE 372 taught by David Tse at Stanford

# Current sequencing technologies

	Second gen. (Illumina)	Oxford Nanopore (MinIon)	PacBio
read length (bases)	100-500	10K-100K	10K-20K
error rates	< 1%	10-15%	10-15%
speed (time/base)	6 mins/base/strand	250 bases/s	3 bases/s
# of reads in parallel	$10^9$	2000	150K
throughput (total # of bases/s)	3M	500K	450K

Table from the course EE 372: Data Science for High-Throughput Sequencing.  
taught by David Tse at Stanford



MinION, a palm-sized gene sequencer made by UK-based Oxford Nanopore Technologies

# Where is the Poisson?

- $G$  - genome length (in bp)
- $L$  - short read average length
- $N$  - number of short read sequenced
- $\lambda$  - sequencing coverage redundancy =  $LN/G$
- $x$  - number of short reads covering a given site on the genome

$$P(x) = \frac{\lambda^x e^{-\lambda}}{x!}$$

Poisson as a limit of Binomial: For a given site on the genome for each short read Prob(site covered):  $p=L/G$  is very small. Number of attempts (short reads):  $N$  is very large. Their product (sequencing redundancy):  $\lambda = NL/G$  is  $O(1)$ .



What fraction of genome is covered?

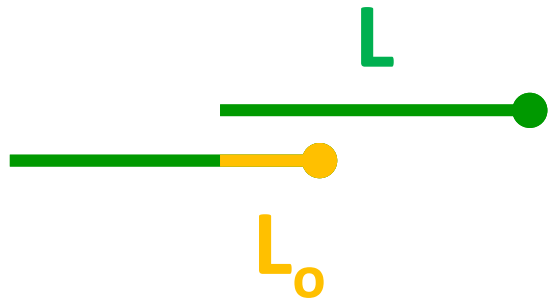
# What fraction of genome is covered?

- Coverage:  $\lambda = NL/G$ ,  
*X* – random variable equal to the number of times a given site is covered by short reads.  
Poisson:  $P(X=x) = \lambda^x \exp(-\lambda) / x!$   
 $P(X=0) = \exp(-\lambda)$ ,  $P(X>0) = 1 - \exp(-\lambda)$
- Total length covered:  $G * [1 - \exp(-\lambda)]$

$\lambda$	2	4	6	8	10	12
Mean proportion of genome covered	.864665	.981684	.997521	.999665	.999955	.999994

Table 5.1. The mean proportion of the genome covered for different values of  $\lambda$

# How long should the overlap be to connect two short reads?



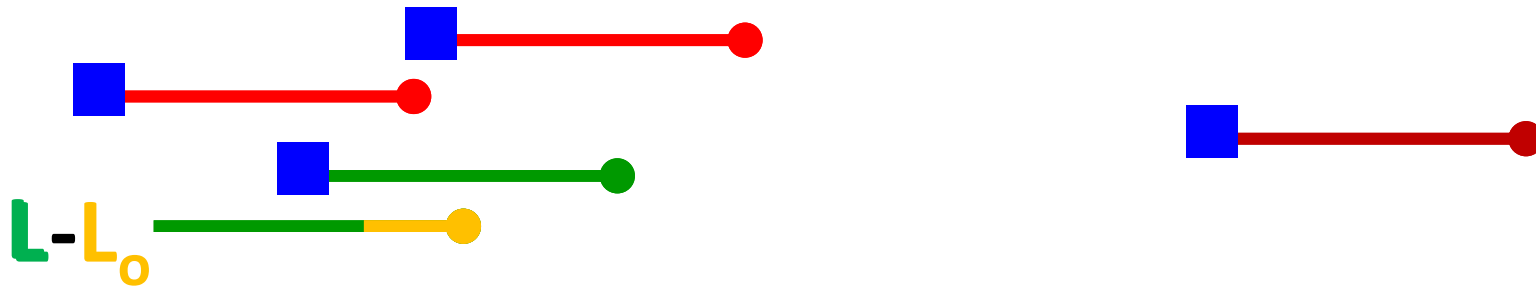
If DNA was a random chain with  $p_A = p_C = p_G = p_T = 1/4$

$L_0 \sim 16-20$  would be enough

$$2 \cdot G \cdot 4^{-L_0} = 2 \cdot 3 \times 10^9 \cdot 4^{-16} = 1.4$$

$$2 \cdot 3 \times 10^9 \cdot 4^{-20} = 0.0055 \ll 1$$

# How many contigs?

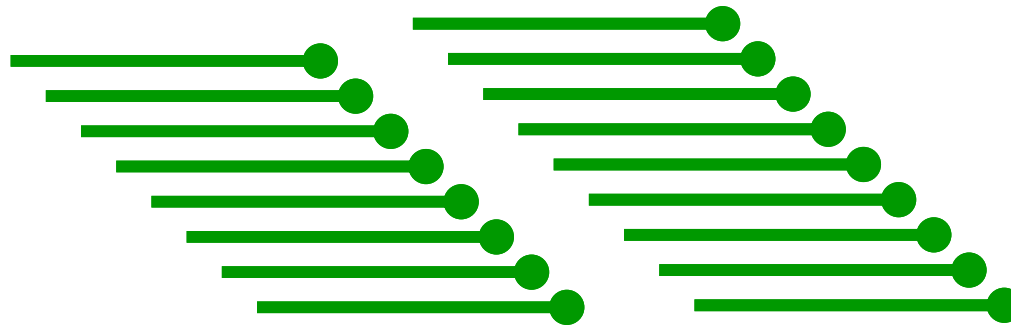


**G**

$$P(\text{short read can be extended by another short read}) = \frac{L - L_0}{G} = p$$

$$P(\text{short read cannot be extended by any short reads}) = e^{-pN} \approx Ne^{-\lambda}$$

$$\text{number of contigs} = Ne^{-pN} \approx Ne^{-\lambda}$$





# How many contigs?

- A given short read is the right end of a contig if and only if no left ends of other short reads fall within first  $L-L_{overlap}$  base pairs
- The left end of another short read has the probability  $p=(L-L_{overlap})/G$  to fall within a given read. There are  $N-1$  other reads.
- The expected number of left ends inside a given short read is  $p \cdot (N-1)=(N-1) \cdot (L-L_{overlap})/G \approx \lambda$  (if  $L \gg L_{overlap}$ )
- Probability that no left ends fall inside a given short read is  $\exp(-\lambda)$ . Thus, the Number of contigs is  $N_{contigs} = Ne^{-\lambda}$ :

$\lambda$	0.5	0.75	1	1.5	2	3	4	5	6	7
Mean number of contigs	60.7	70.8	73.6	66.9	54.1	29.9	14.7	6.7	3.0	1.3

Table 5.2. The mean number of contigs for different levels of coverage, with  $G = 100,000$  and  $L = 500$ .

# Average length of a contig?

- Length of a genome covered:

$$G_{covered} = G \cdot P(X > 0) = G \cdot (1 - \exp(-\lambda))$$

- Number of contigs  $N_{contigs} = N \cdot e^{-\lambda}$

- Average length of a contig =

$$\langle L \rangle = \sum_i L_i / N_{contigs} = G_{covered} / N_{contigs} =$$

$$G \cdot (1 - \exp(-\lambda)) / N \cdot e^{-\lambda} = L \cdot (1 - \exp(-\lambda)) / \lambda \cdot e^{-\lambda}$$

$\lambda$	2	4	6	8	10
Mean contig size	1,600	6,700	33,500	186,000	1,100,000

Table 5.3. The mean contig size for different values of  $a$  for the case  $L = 500$ .

# Estimate

- Human genome is  $3 \times 10^9$  bp long
- Chromosome 1 is about  $G = 0.25 \times 10^9$  bp
- Illumina generates short reads  $L = 100$  bp long
- What number of reads  $N$  are needed to completely assemble the 1<sup>st</sup> chromosome?
- The formula to use is:  $1 = N_{contigs} = N e^{-\lambda} = N e^{-NL/G}$
- Answer:  $N = 4.4 \times 10^7$  short (100bp) reads  
Test:  $4.4e7 * \exp(-4.4e7 * 100 / 0.25e9) = 0.99997$
- What coverage redundancy  $\lambda$  will it be?  
Answer:  $\lambda = NL/G = 17.6$  coverage redundancy

# How much would it cost to assemble human genome now?

- Human Genome Project: **\$2.7 billion** in 1991 dollars.
- Now a **de novo full assembly** of the whole human genome would now cost  $3 \times 10^9 \times 17.6 / 10^6 \times 0.1\$/\text{MB} = \$5300$
- **2<sup>nd</sup> genome** (and after) would be **even cheaper** as we would already have a **reference genome** to which we can **map short reads**. (Puzzle: picture on the box)
- But this is a **naïve estimate**. In reality, there are complications. See next slides:

# What spoils these estimates?

```
>gi|224514922|ref|NT_024477.14| Homo sapiens chromosome 12 genomic
contig, GRCh37.p13 Primary Assembly (displaying 3' end)
CGGGAAATCAAAAGCCCCTCTGAATCCTGCGCACCGAGATTCTCCCCAGCCAAGGTGAGGCGGCAGCAGT
GGGAGATCCACACCGTAGCATTGGAACACAAATGCAGCATTACAAATGCAGACATGACACCGAAAATATA
ACACACCCCATTGCTCATGTAACAAGCACCTGTAATGCTAATGCACTGCCTCAAAACAAAATATTAATAT
AAGATCGGCAATCCGCACACTGCCGTGCAGTGCTAAGACAGCAATGAAAATAGTCAACATAATAACCCTA
ATAGTGTTAGGGTTAGGGTCAGGGTCCCGGTCCGGGTCCGGGTCCGGGTCCGGGTCCGGGTCCGGGTCA
GGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGT
TAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGG
GTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTA
GGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGT
TAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGGGTTAGG
GTTAGGGTTAGGGTTAGGGTTAG
```

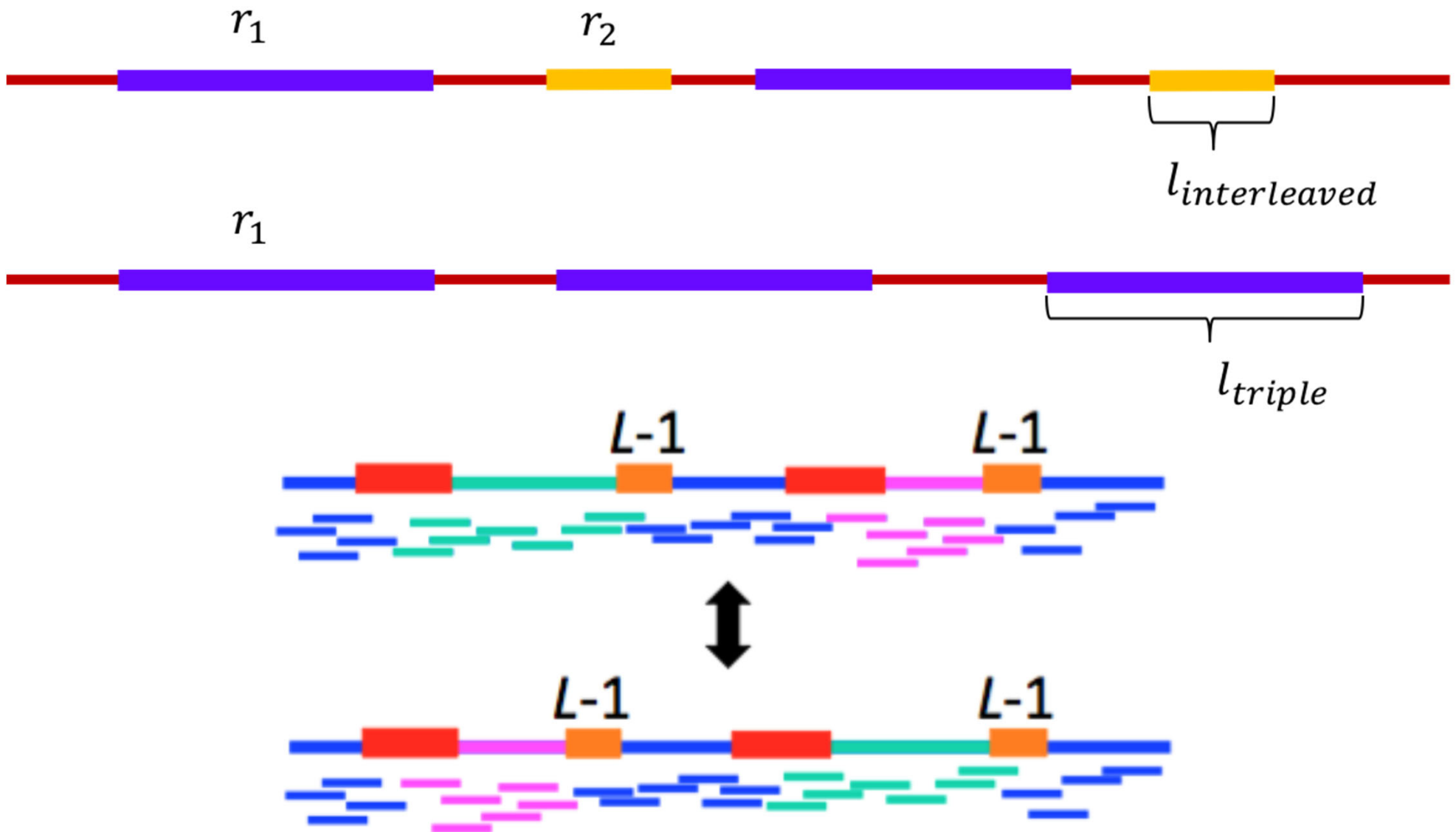
**FIGURE 8.11** A BLASTN search of the human genome (all assemblies) database was performed at the NCBI website using **TTAGGGTTAGGGTTAGGG** as query (i.e., three TTAGGG repeats). There were matches to hundreds of genomic scaffolds. This figure shows an example (NT\_024477.14) assigned to the **telomere of chromosome 12q having many dozens of TTAGGG repeats.** These occurred at the 3' end of the genomic contig sequence.

There were **100s of matches** while **one expects  $\ll 1$  match:**

$$2 \cdot 3 \times 10^9 \cdot 4^{-18} = 0.08 \ll 1$$

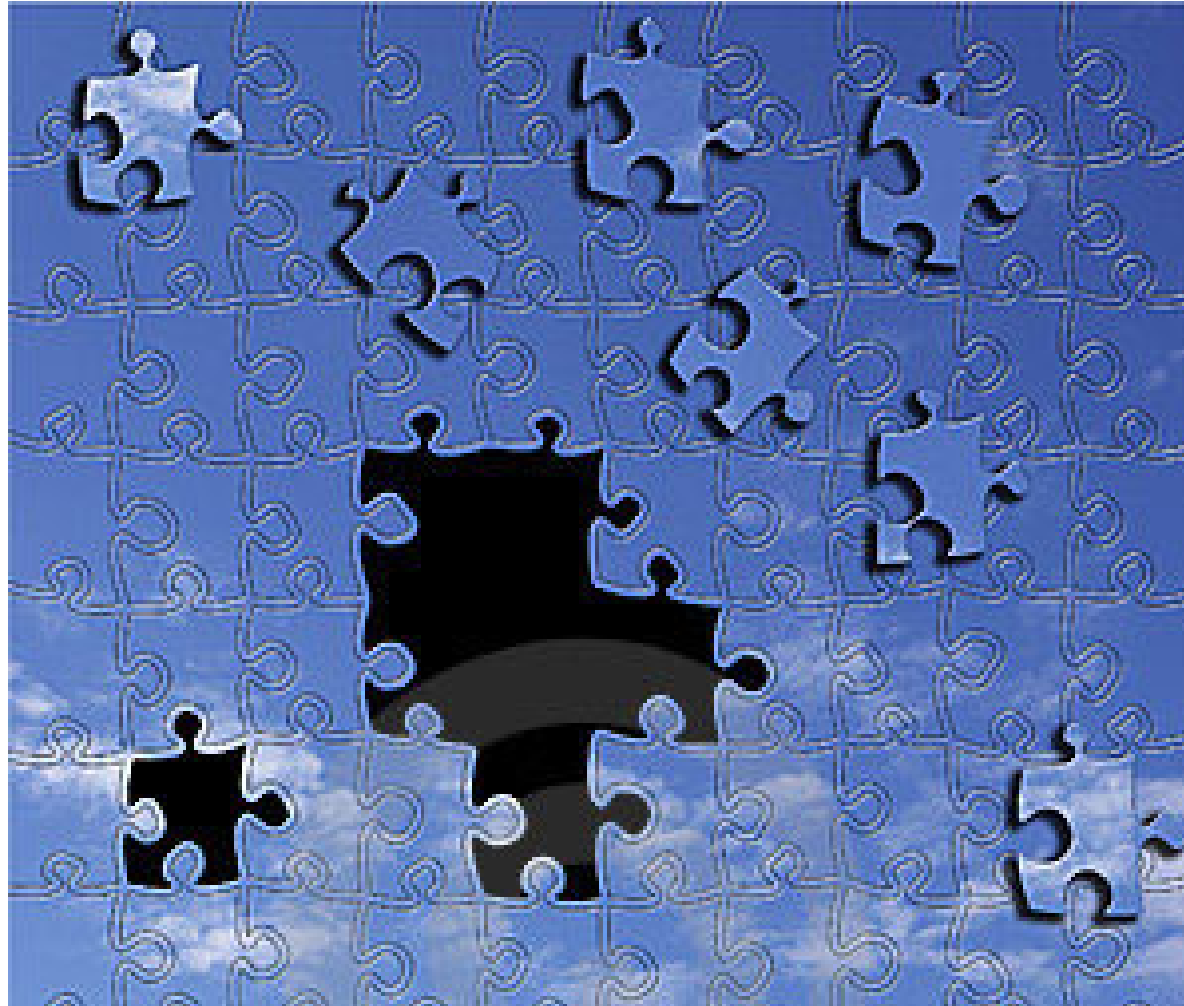
**DNA repeats** make assembly difficult

# Why repeats make assembly difficult?



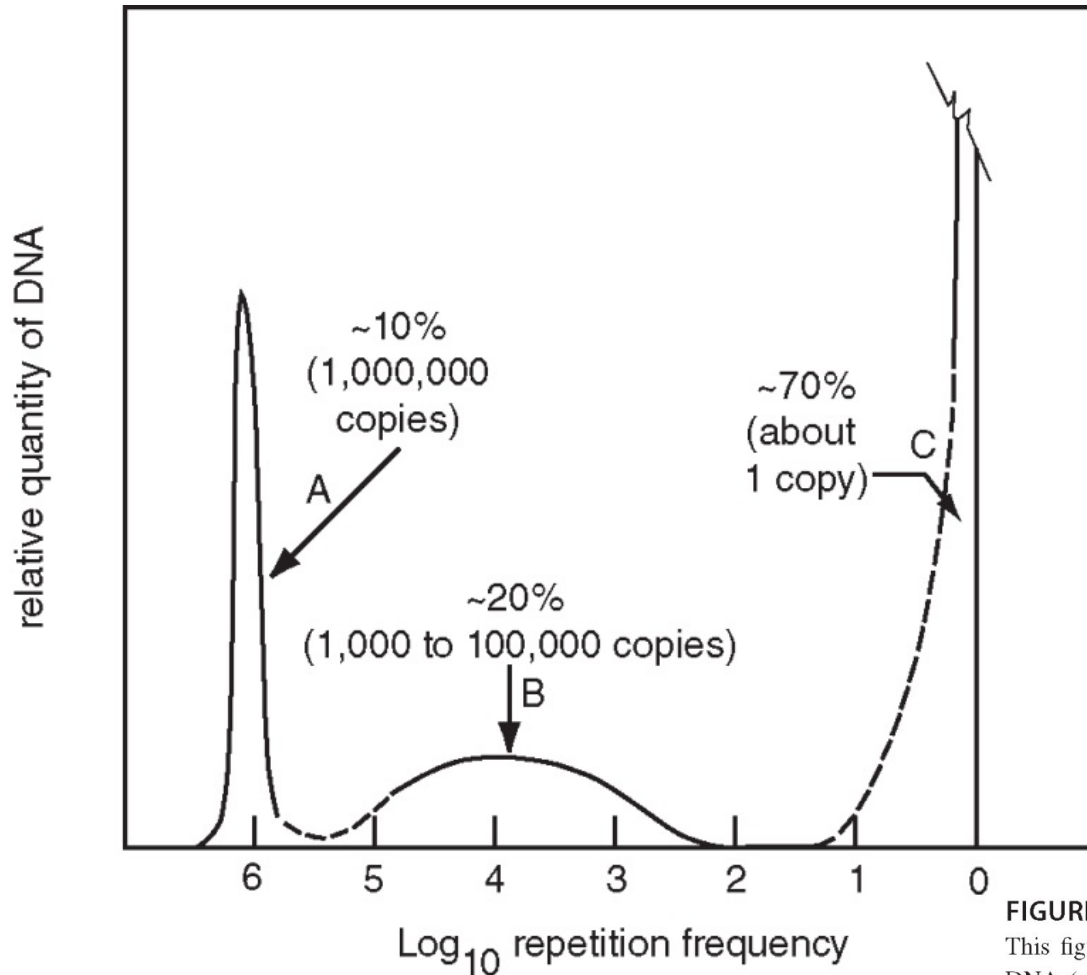
Images from the course EE 372: Data Science for High-Throughput Sequencing.  
taught by David Tse at Stanford

**Repeats** are like sky puzzle pieces





# How many repeats are in eukaryotic genomes?

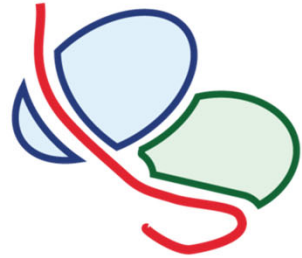


Data for **mouse genome** obtained in 1961 (sic!) using DNA denaturation and renaturation curves

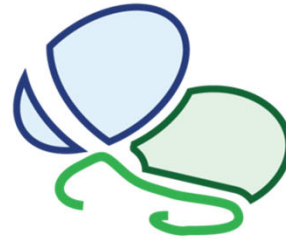
**FIGURE 8.6** The complexity of genomic DNA can be estimated by denaturing then renaturing DNA. This figure (redrawn from Britten and Kohne, 1968) depicts the relative quantity of mouse genomic DNA (y axis) versus the logarithm of the frequency with which the DNA is repeated. The data are derived from a  $C_0 t_{1/2}$  curve, which describes the percent of genomic DNA that reassociates at particular times and DNA concentrations. A large  $C_0 t_{1/2}$  value implies a slower reassociation reaction. Three classes are apparent. The fast component accounts for 10% of mouse genomic DNA (arrow A), and represents highly repetitive satellite DNA. An intermediate component accounts for about 20% of mouse genomic DNA and contains repeats having from 1000 to 100,000 copies. The slowly reassociating component, comprising 70% of the mouse genome, corresponds to unique, single-copy DNA. Britten and Kohne (1968) obtained similar profiles from other eukaryotes, although distinct differences were evident between species. Used with permission.



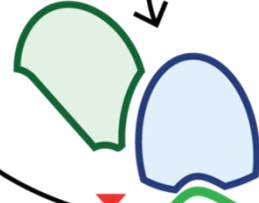
Formation of  
Ribonucleoprotein complexes



Reverse  
Transcription



Integration

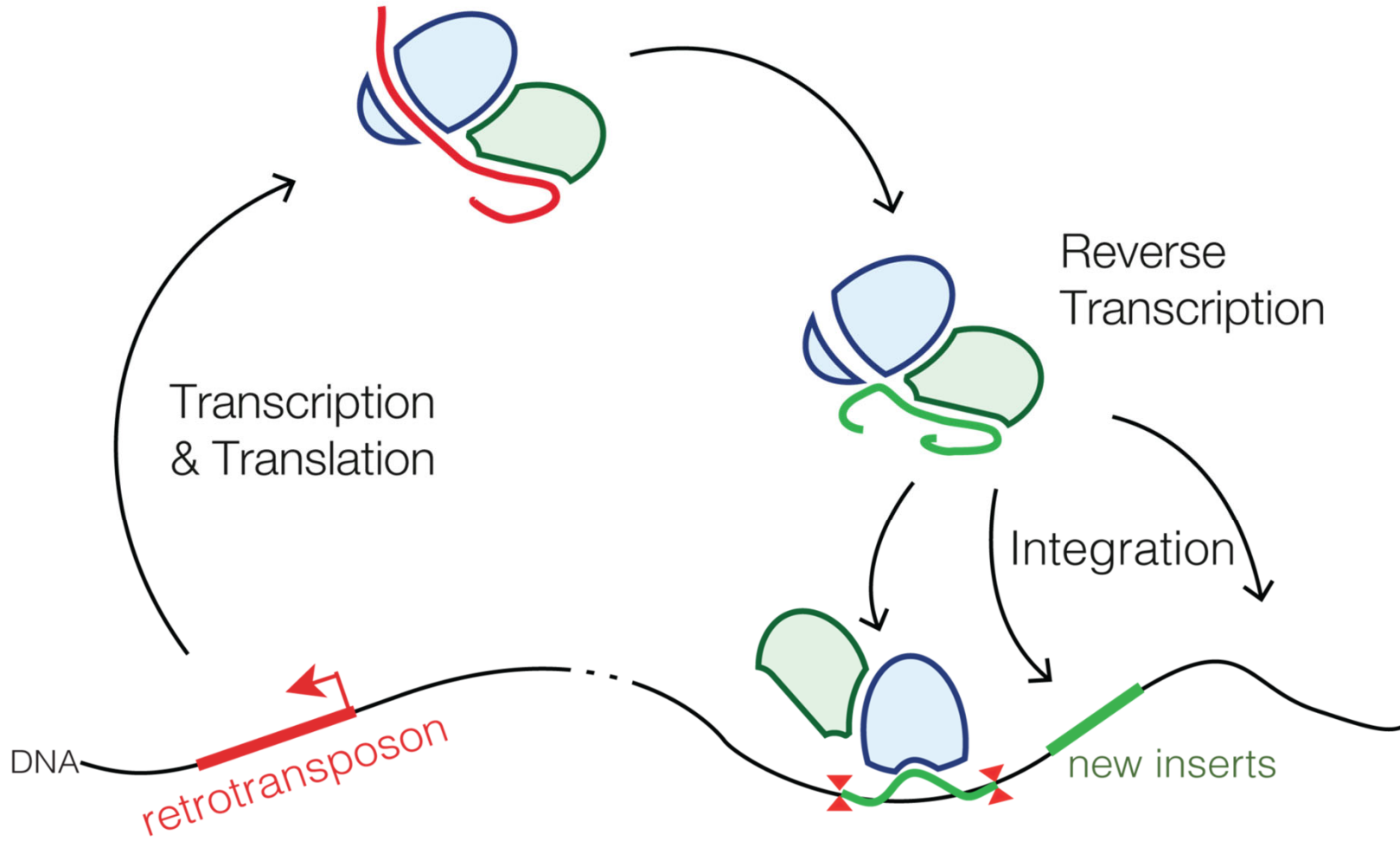


new inserts

Transcription  
& Translation


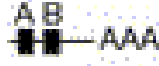




DNA

retrotransposon

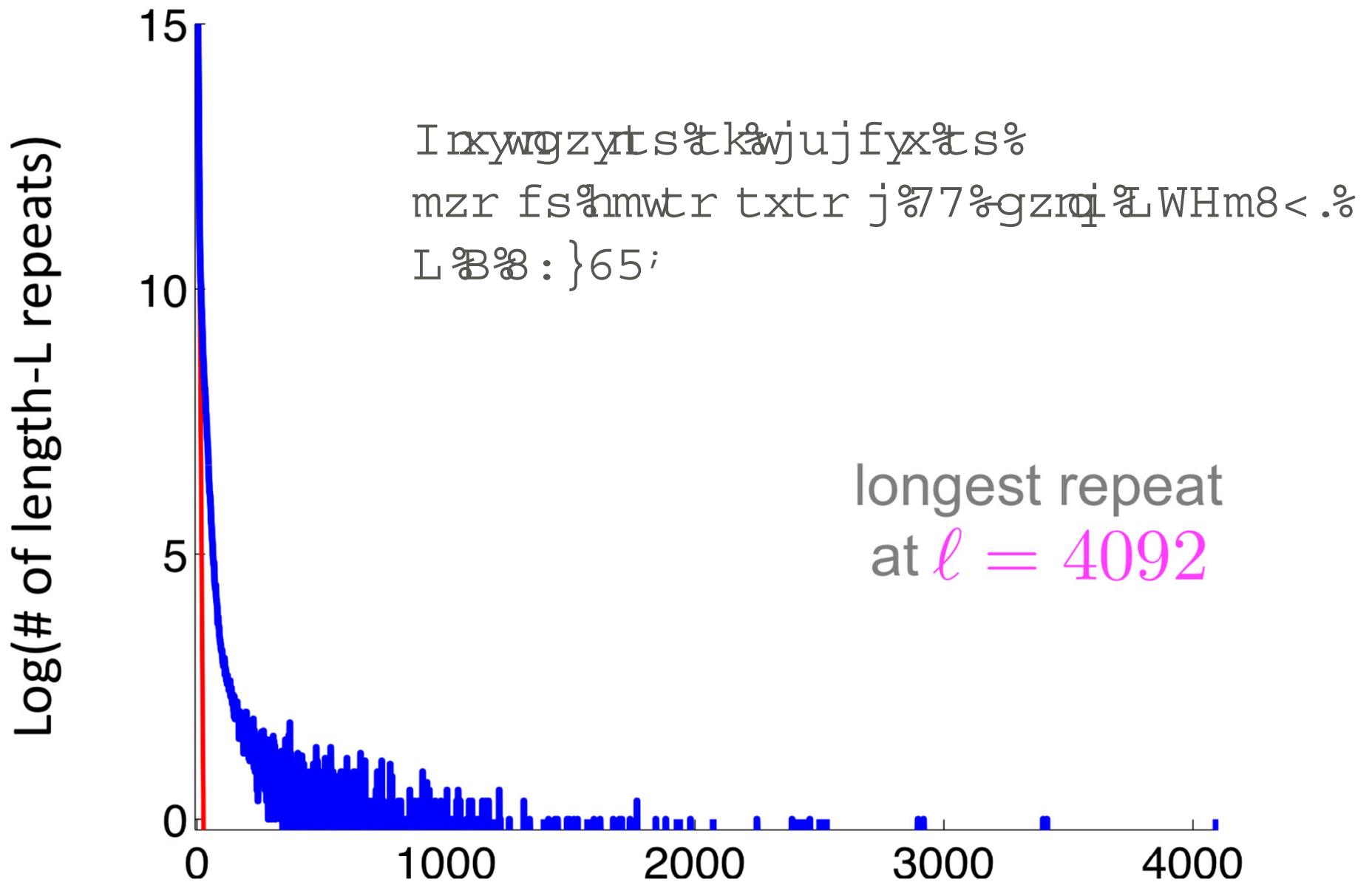


# Almost all transposable elements in mammals fall into one of four classes

Classes of interspersed repeat in the human genome

			Length	Copy number	Fraction of genome
LINEs	Autonomous		6–8 kb	850,000	21%
	Non-autonomous		100–300 bp		
Retrovirus-like elements	Autonomous		6–11 kb	450,000	8%
	Non-autonomous		1.5–3 kb		
DNA transposon fossils	Autonomous		2–3 kb	300,000	3%
	Non-autonomous		80–3,000 bp		

Slide by Ross Hardison, Penn State U.



# How to assemble a real genome with repeats?

Here we assume a “de novo” assembly  
without help from the previously  
assembled genomes



Nicolaas Govert de Bruijn (1918 – 2012) was a Dutch mathematician, noted for his many contributions in the fields of **graph theory**, analysis, number theory, combinatorics and logic

Courtesy of [Ben Langmead](#). Used with permission.

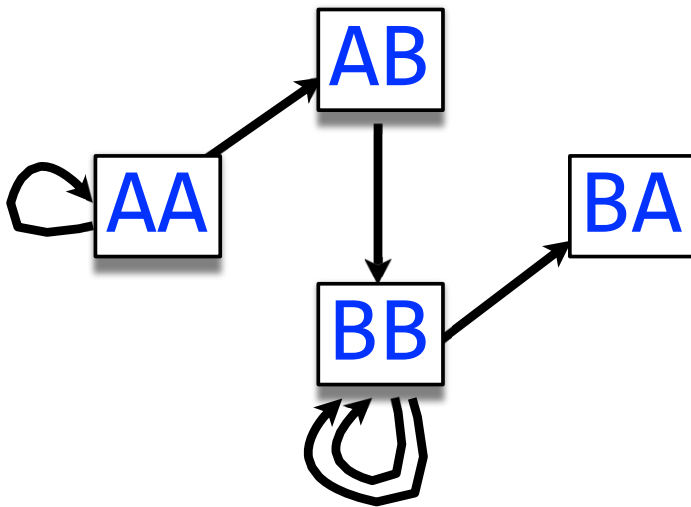
<http://www.langmead-lab.org/teaching-materials/>

# De Bruijn graph

genome: **AAABBBBA**

3-mers: **AAA, AAB, ABB, BBB, BBB, BBA**

L/R 2-mers: **AA, AA    AA, AB    AB, BB    BB, BB    BB, BB    BB, BA**



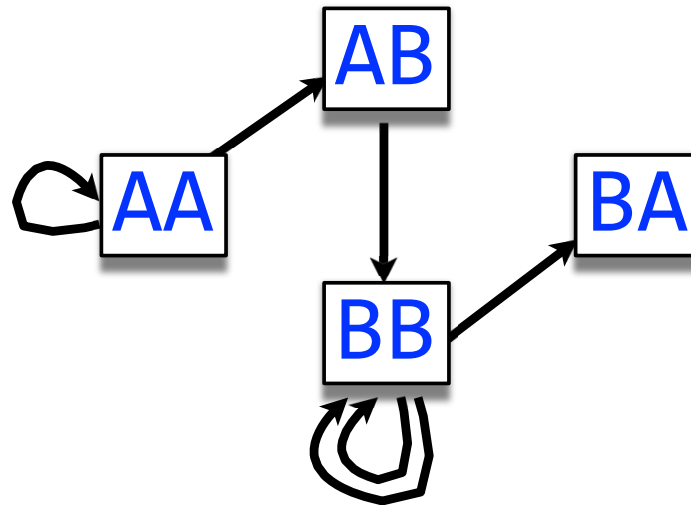
One edge per **every**  $k$ -mer

One node per **distinct**  $k-1$ -mer

Courtesy of [Ben Langmead](#). Used with permission.

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# De Bruijn graph

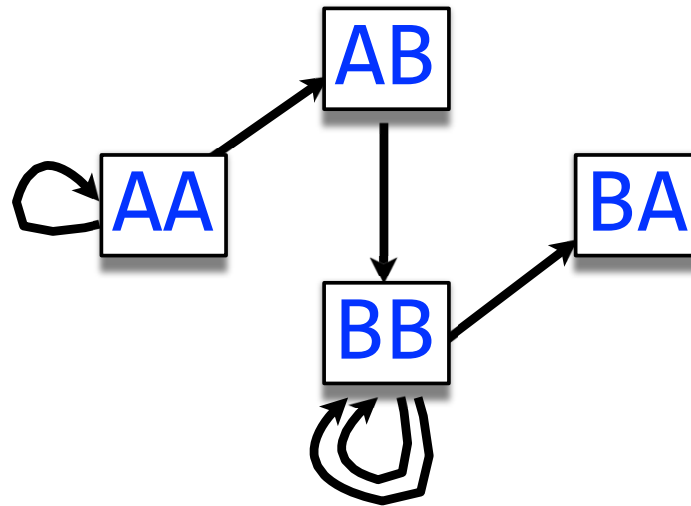


Walk crossing each edge exactly once gives a reconstruction of the genome

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# Assembly = Eulerian walk on De Bruijn graph

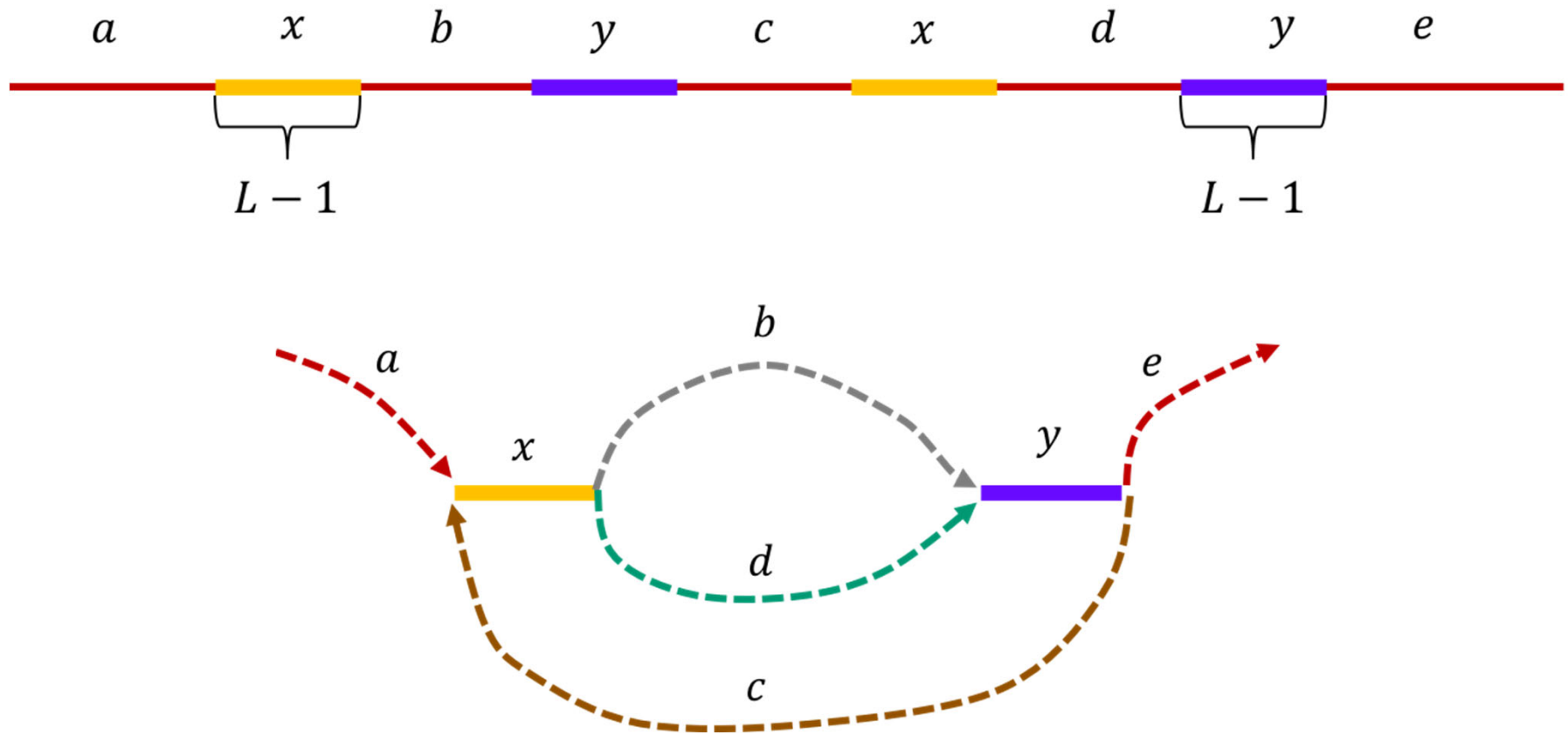


AAABBBBA

Walk crossing each edge exactly once gives a reconstruction of the genome. This is an *Eulerian walk*.



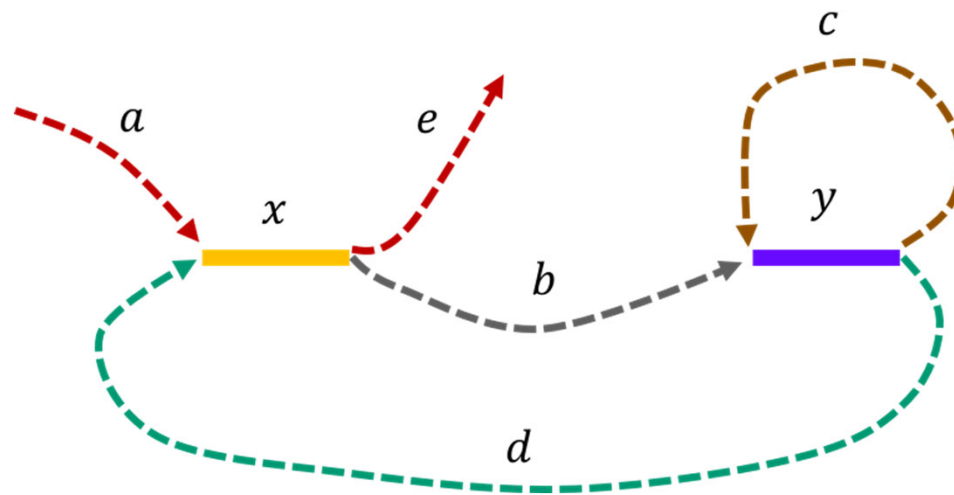
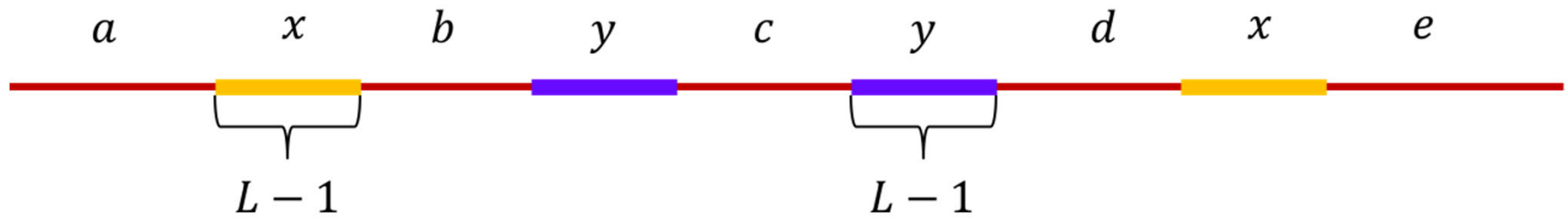
# Why interleaved repeats are dangerous?



The two Eulerian paths that are on the graph:  
 $a-x-b-y-c-x-d-y-e$  and  $a-x-d-y-c-x-b-y-e$

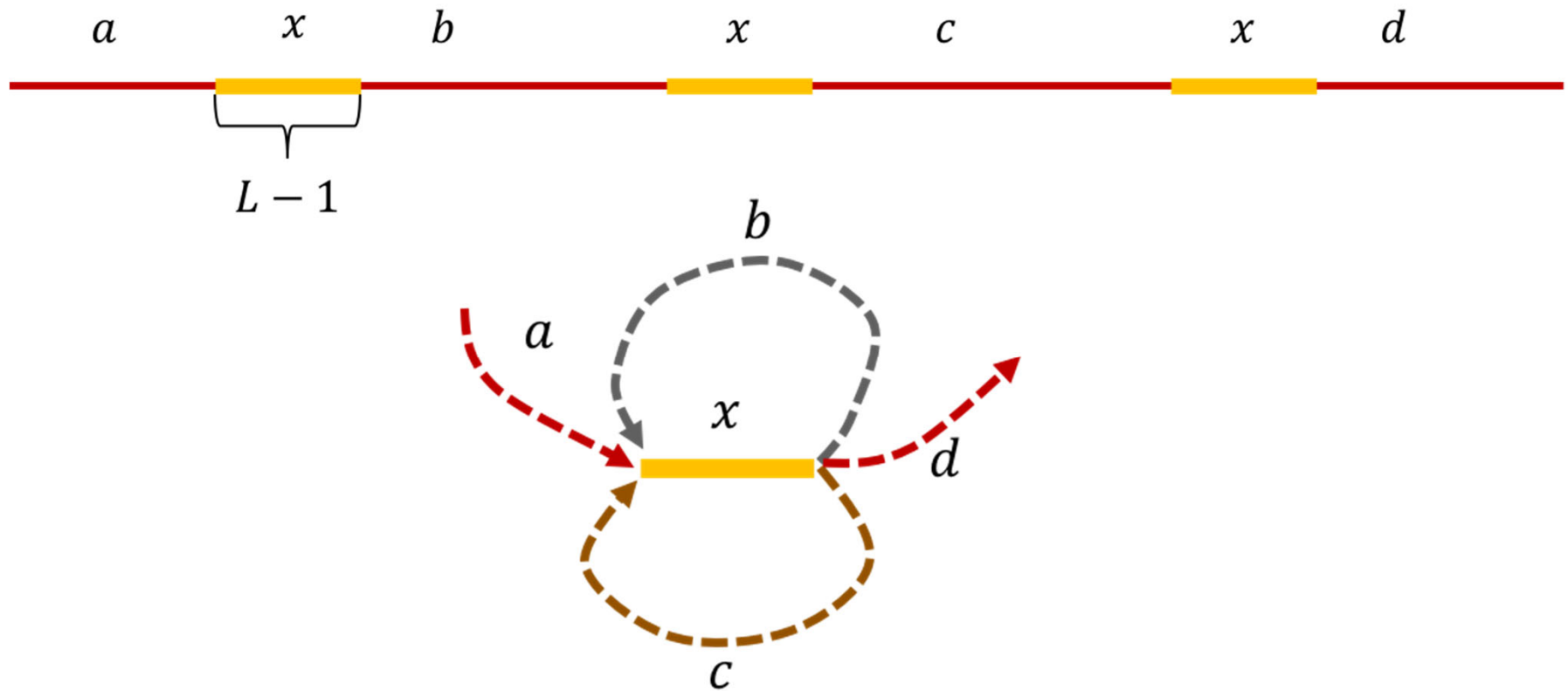
Images from the course [EE 372: Data Science for High-Throughput Sequencing](#).  
taught by David Tse at Stanford

# Why non-interleaved repeats are safe?



The only Eulerian path is:  $a-x-b-y-c-y-d-x-e$

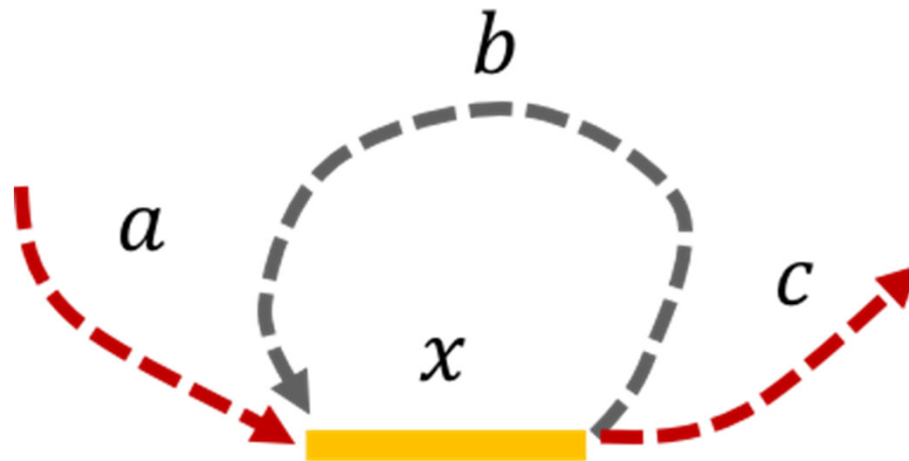
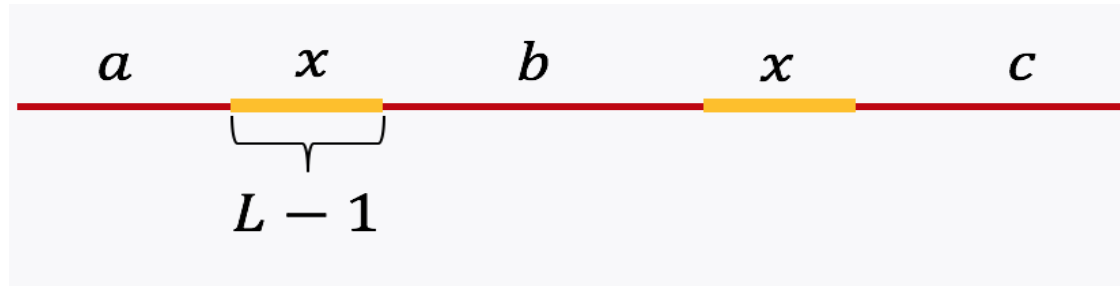
# Why triple repeats are dangerous?



The two Eulerian paths that are on the graph:  
 $a-x-b-x-c-x-d$  and  $a-x-c-x-b-x-d$

Images from the course EE 372: Data Science for High-Throughput Sequencing.  
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# Why double repeats are safe?



The only Eulerian path is:  $a-x-b-x-c$

# Pavel Pevzner's theorem

- **Theorem [Pevzner 1995]:**  
If  $L$ , the read length, is strictly greater than  $\max(\ell_{\text{interleaved}}, \ell_{\text{triple}})$ , then the de Bruijn graph has a unique Eulerian path corresponding to the original genome.

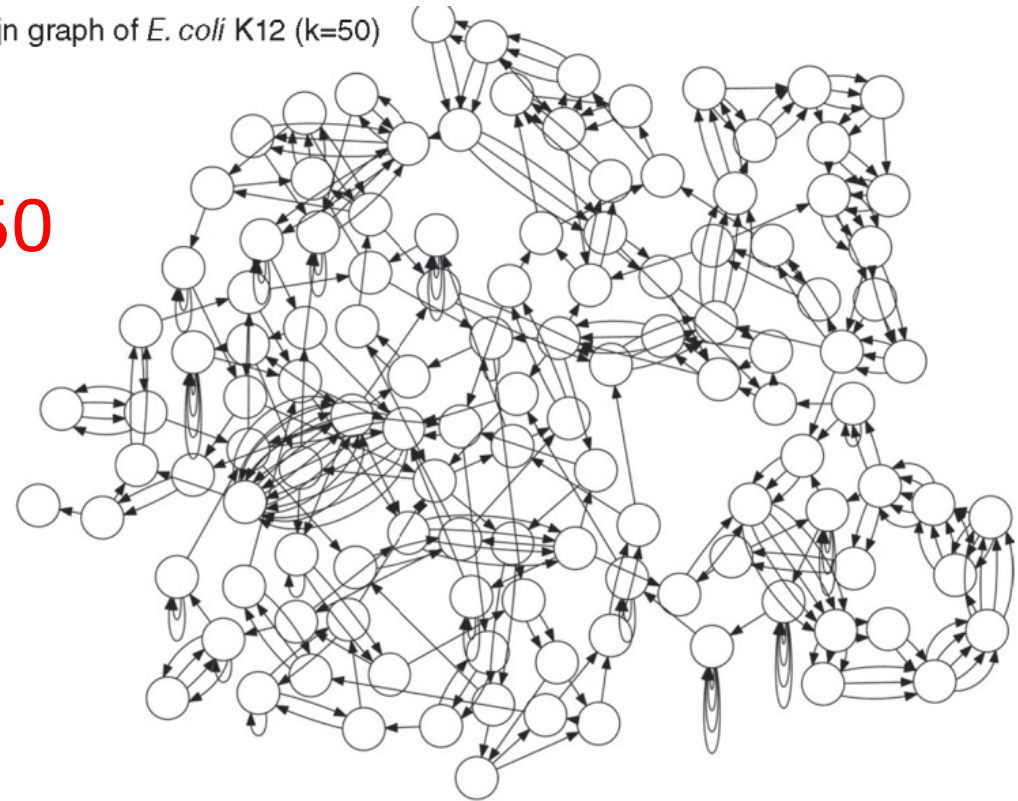


**Pavel Pevzner**  
is the Ronald R. Taylor Chair and  
Distinguished Professor of  
Computer Science and Engineering  
at University of California, San Diego.  
His Alma Mater is  
Moscow Institute of  
Physics and Technology  
in Russia.

# How to assemble a genome with repeats?

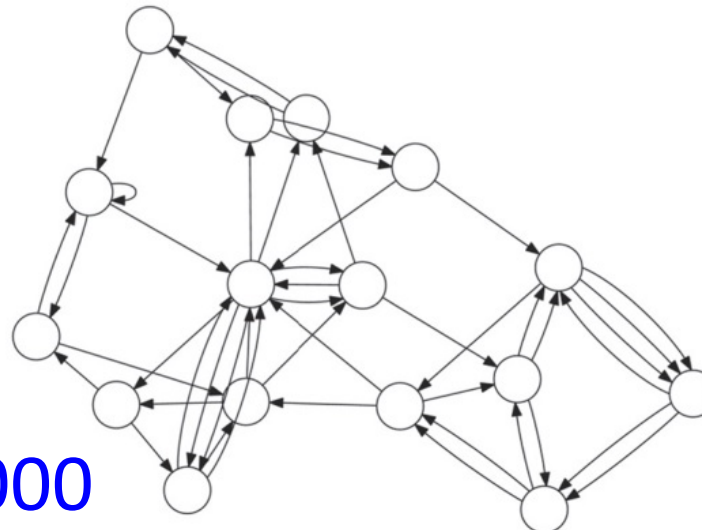
- Answer:  
longer reads
- But:  
cheap sequencing  
=  
short reads

(a) de Bruijn graph of *E. coli* K12 ( $k=50$ )



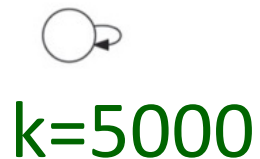
$k=50$

(b) de Bruijn graph ( $k=1,000$ )



$k=1000$

(c) de Bruijn graph ( $k=5,000$ )



$k=5000$

Technology	Read length (bp)
Roche 454	700
Illumina	50–250
SOLiD	50
Ion Torrent	400
Pacific Biosciences	>10,000



Credit: XKCD  
comics

# WHY ARE THERE SLAVES IN THE BIBLE

WHY DO TWINS HAVE DIFFERENT FINGERPRINTS  
WHY ARE AMERICANS AFRAID OF DRAGONS

WHY IS HTTPS CROSSED OUT IN RED  
WHY IS THERE A LINE THROUGH HTTPS  
WHY IS THERE A RED LINE THROUGH HTTPS ON FACEBOOK  
WHY IS HTTPS IMPORTANT

# QUESTIONS FOUND IN GOOGLE AUTOCOMplete



WHY ARE THERE WEEKS  
WHY DO I FEEL DIZZY

WHY AREN'T ECONOMISTS RICH  
WHY DO AMERICANS CALL IT SOCCER  
WHY ARE MY EARS RINGING  
WHY ARE THERE SO MANY AVENGERS  
WHY ARE THE AVENGERS FIGHTING THE X MEN  
WHY IS WOLVERINE NOT IN THE AVENGERS

WHY ARE THERE SWARMS OF GNATS  
WHY IS THERE PHLEGM  
WHY ARE THERE SO MANY CROWS IN ROCHESTER, MN  
WHY IS PSYCHIC WEAK TO BUG  
WHY DO CHILDREN GET CANCER  
WHY IS POSEIDON ANGRY WITH ODYSSEUS  
WHY IS THERE ICE IN SPACE

# WHY ARE THERE ANTS IN MY LAPTOP

WHY IS EARTH TILTED  
WHY IS SPACE BLACK  
WHY IS OUTER SPACE SO COLD  
WHY ARE THERE PYRAMIDS ON THE MOON  
WHY IS NASA SHUTTING DOWN



WHY IS THERE AN OWL IN MY BACKYARD  
WHY IS THERE AN OWL OUTSIDE MY WINDOW  
WHY IS THERE AN OWL ON THE DOLLAR BILL  
WHY DO OWLS ATTACK PEOPLE  
WHY ARE AK 47s SO EXPENSIVE  
WHY ARE THERE HELICOPTERS CIRCLING MY HOUSE  
WHY ARE THERE GODS  
WHY ARE THERE TWO SPOCKS

WHY ARE DOGS AFRAID OF FIREWORKS  
WHY IS THERE NO KING IN ENGLAND

WHY DO WHALES JUMP  
WHY ARE WITCHES GREEN  
WHY ARE THERE MIRRORS ABOVE BEDS  
WHY DO I SAY UH  
WHY IS SEA SALT BETTER  
WHY ARE THERE TREES IN THE MIDDLE OF FIELDS  
WHY IS THERE NOT A POKEMON MMO  
WHY IS THERE LAUGHING IN TV SHOWS  
WHY ARE THERE DOORS ON THE FREEWAY  
WHY ARE THERE SO MANY SVCHOST.EXE RUNNING  
WHY AREN'T THERE ANY COUNTRIES IN ANTARCTICA  
WHY ARE THERE SCARY SOUNDS IN MINECRAFT  
WHY IS THERE KICKING IN MY STOMACH  
WHY ARE THERE TWO SLASHES AFTER HTTP  
WHY ARE THERE CELEBRITIES  
WHY DO SNAKES EXIST  
WHY DO OYSTERS HAVE PEARLS  
WHY ARE DUCKS CALLED DUCKS  
WHY DO THEY CALL IT THE CLAP  
WHY ARE KYLE AND CARTMAN FRIENDS  
WHY IS THERE AN ARROW ON AANG'S HEAD  
WHY ARE TEXT MESSAGES BLUE  
WHY ARE THERE MUSTACHES ON CLOTHES  
WHY ARE THERE MUSTACHES ON CARS  
WHY ARE THERE MUSTACHES EVERYWHERE  
WHY ARE THERE SO MANY BIRDS IN OHIO  
WHY IS THERE SO MUCH RAIN IN OHIO  
WHY IS OHIO WEATHER SO WEIRD

WHY ARE THERE MALE AND FEMALE BIKES  
WHY ARE THERE TINY SPIDERS IN MY HOUSE  
WHY DO SPIDERS COME INSIDE  
WHY ARE THERE HUGE SPIDERS IN MY HOUSE  
WHY ARE THERE LOTS OF SPIDERS IN MY HOUSE  
WHY ARE THERE SPIDERS IN MY ROOM  
WHY ARE THERE SO MANY SPIDERS IN MY ROOM  
WHY DO SPIDER BITES ITCH  
WHY IS DYING SO SCARY



WHY ARE THERE BRIDESMAIDS  
WHY DO DYING PEOPLE REACH UP  
WHY AREN'T THERE VARICOSE ARTERIES  
WHY ARE OLD KUNGONS DIFFERENT  
WHY IS THERE HELL IF GOD FORGIVES  
WHY IS THERE NO GPS IN LAPTOPS  
WHY DO KNEES CLICK  
WHY AREN'T THERE E GRADES  
WHY IS ISOLATION BAD  
WHY DO BOYS LIKE ME  
WHY DON'T BOYS LIKE ME  
WHY IS THERE ALWAYS A JAVA UPDATE  
WHY ARE THERE RED DOTS ON MY THIGHS  
WHY IS LYING GOOD



WHY IS MT VESUVIUS THERE  
WHY DO THEY SAY T MINUS  
WHY ARE THERE OBELISKS  
WHY ARE WRESTLERS ALWAYS WET  
WHY ARE OCEANS BECOMING MORE ACIDIC  
WHY IS ARWEN DYING  
WHY AREN'T MY QUAIL LAYING EGGS  
WHY AREN'T MY QUAIL EGGS HATCHING  
WHY AREN'T THERE ANY FOREIGN MILITARY BASES IN AMERICA

WHY ARE CIGARETTES LEGAL  
WHY ARE THERE DUCKS IN MY POOL  
WHY IS JESUS WHITE  
WHY IS THERE LIQUID IN MY EAR  
WHY DO Q TIPS FEEL GOOD  
WHY DO GOOD PEOPLE DIE



WHY ARE ULTRASOUNDS IMPORTANT  
WHY ARE ULTRASOUND MACHINES EXPENSIVE  
WHY IS STEALING WRONG

# Geometric Distribution

- A series of **Bernoulli trials** with **probability of success =  $p$** . continued **until the first success**.  $X$  is the number of trials.
- Compare to: Binomial distribution has:
  - Fixed number of trials =  $n$ .  $P(X = x) = C_x^n p^x (1 - p)^{n-x}$
  - Random number of successes =  $x$ .
- Geometric distribution has reversed roles:
  - Random number of trials,  $x$
  - Fixed number of successes, in this case 1.
  - Success always comes in the end: so no combinatorial factor  $C_x^n$
  - $P(X=x) = p(1-p)^{x-1}$  where:  
 $x-1 = 0, 1, 2, \dots$ , the number of failures until the 1<sup>st</sup> success.
- **NOTE OF CAUTION: Matlab, Mathematica**, and many other sources use  $x$  to denote the **number of failures until the first success**. We stick with **Montgomery-Runger notation**



# Geometric Mean & Variance

# Geometric Mean & Variance

$$P(X=x) = p(1-p)^{x-1} = p \cdot q^{x-1}$$

$$\sum_{x=1}^{\infty} P(X=x) = \frac{p}{1-q} = \frac{p}{p} = 1$$

$$q \frac{\partial}{\partial q} \sum_{x=1}^{\infty} P(X=x) = \sum_{x=1}^{\infty} (x-1) P(X=x) = \frac{pq}{(1-q)^2} = \frac{q}{p}$$

$$\langle x \rangle = \sum_{x=1}^{\infty} (x-1) P(X=x) + 1 = \frac{1-p}{p} + 1 = \frac{1}{p}$$

# Geometric Mean & Variance

- If  $X$  is a geometric random variable (according to Montgomery-Bulmer) with parameter  $p$ ,

$$\mu = E(X) = \frac{1}{p} \quad \text{and} \quad \sigma^2 = V(X) = \frac{(1-p)}{p^2} \quad (3-10)$$

- For small  $p$  the **standard deviation**  $\approx$  **mean**
- Very different from Poisson, where it is **variance** = **mean** and **standard deviation** = **mean**<sup>1/2</sup>

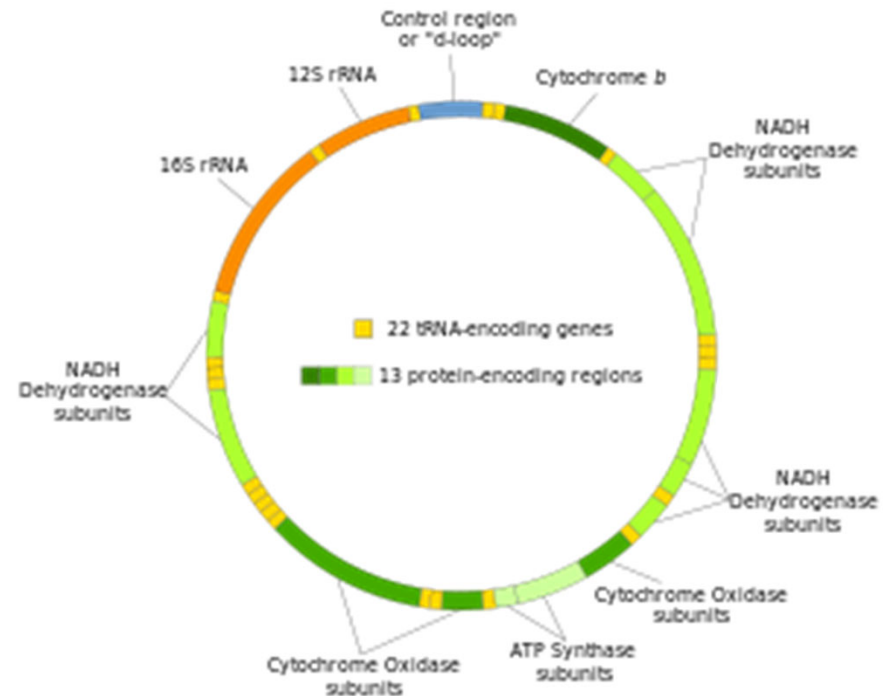
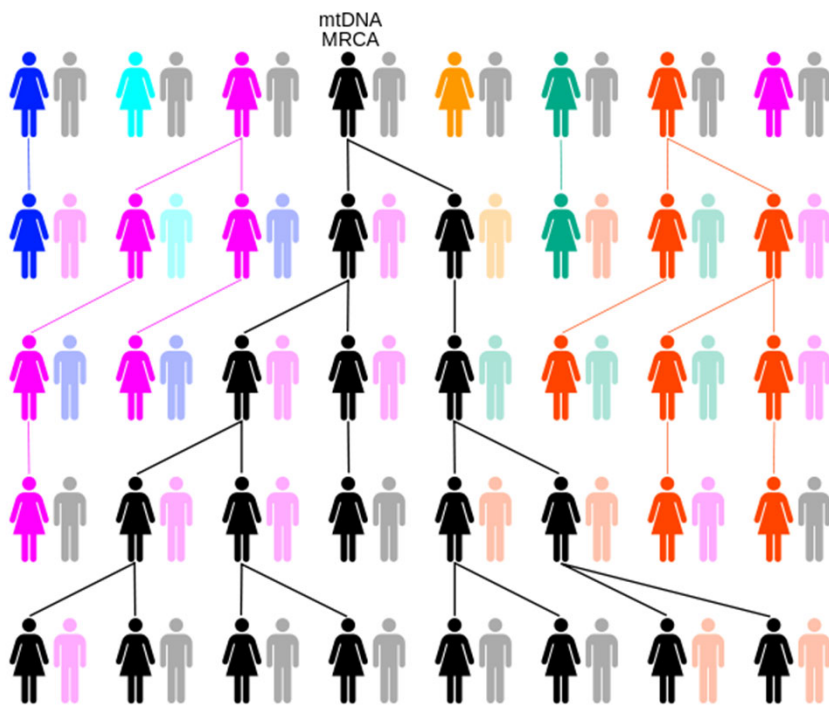
# Matlab exercise

- Find mean, variance, and histogram of 100,000 geometrically-distributed numbers with  $p=0.1$
- Hint: Use help page for random command on how to generate geometrically-distributed random numbers

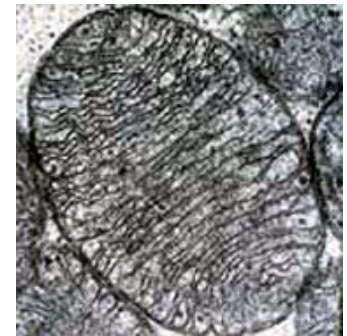
# Matlab: Geometric distributions

- **Stats=100000;**
- **p=0.1;**
- **r2=random('Geometric',p,Stats,1);**
- **r2=r2+1;**
- **disp(mean(r2));**
- **disp(var(r2));**
- **disp(std(r2));**
- **[a,b]=hist(r2, 1:max(r2));**
- **p\_g=a./sum(a);**
- **figure; semilogy(b,p\_g,'ko-');**

# Geometric distribution in biology



- Each of our cells has mitochondria with 16.5kb of mtDNA **inherited only from our mother**
- Human mtDNA has 37 genes encoding 13 proteins, 22+2 tRNA & rRNA
- Mitochondria appeared 1.5-2 billion years ago as a symbiosis between an alpha-proteobacterium (1000s of genes) and an archaeon (of UIUC's Carl R. Woese fame)
- Since that time most mitochondrial genes were transferred into the nucleus
- Plants also have plastids with genomes related to cyanobacteria



# Time to the last common (maternal) ancestor follows geometric distribution

- **Constant population** of  $N$  women
- **Random number** of (female) **offsprings**. Average is 1 (but can be 0 or 2)
- **Randomly** pick **two women**.  
Question: how many **generations  $T$**  since their **last maternal ancestor**?
- $T$  is a random variable What is its PMF:  **$P(T=t)$** ?  
Answer:  $P(T=t)$  follows a **geometric distribution**
- Do these two women have **the same mother**? Yes: **“success”** in finding their last common ancestor ( **$p=1/N$** ).  **$P(T=1)=1/N$** .
- No? **“failure”** ( **$1-p=1-1/N$** ). Go to their mothers and repeat the same question.
- **$P(T=t)=(1-1/N)^{t-1}(1/N) \approx (1/N) \exp(-(t-1)/N)$**
- **$t$**  can be inferred from **the density of differences on mtDNA  $=2\mu t$**

