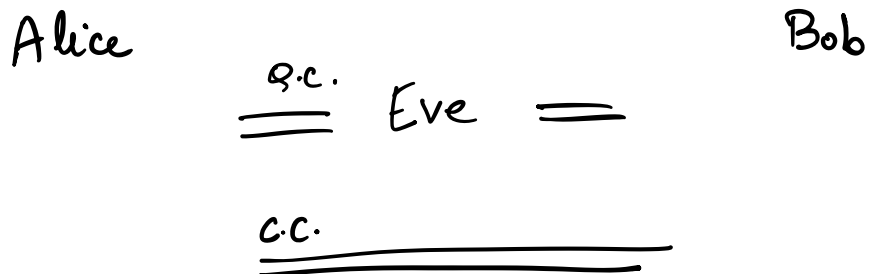


LECTURE -10.

QUANTUM KEY DISTRIBUTION



GOAL :

Alice and Bob agree on a shared key k ,
 such that $(k, \text{Eve's view}) \approx_{\epsilon} (\text{uniform}, \text{Eve's view})$
 \uparrow
statistical indistinguishability

for any two distributions D_1 and D_2 ,
 we say that (D_1, D_2) are ϵ -stat-ind
 $D_1 \approx_{\epsilon} D_2 \iff \Delta(D_1, D_2) \leq \epsilon$.
 Equivalently, $\forall C, \left| \Pr_{d \leftarrow D_1} [C(d)=1] - \Pr_{d \leftarrow D_2} [C(d)=1] \right| \leq \epsilon$

Attempt 0.
Alice

Bob

$$k \leftarrow \{0,1\}^n$$

$$c.c. \xrightarrow{k}$$

But Eve can also observe k .

Attempt 1.

Alice

Bob

$$x \in \{0,1\}^n$$

$$\theta \in \{0,1\}^n$$

$$|\psi\rangle = |\psi_1\rangle \dots |\psi_n\rangle$$

where $|\psi_i\rangle = |x_i\rangle_{\theta_i}$

$$q.c. \xrightarrow{|\psi\rangle}$$

$$c.c. \xleftarrow{\text{I received}}$$

$$c.c. \xrightarrow{\theta = \theta_1 \dots \theta_n}$$

$\forall i$, measure $|\psi_i\rangle$ in basis $\theta_i \rightarrow x_i$

If $\theta_i = 0$ "computational basis" = $\{ |0\rangle |1\rangle \}$

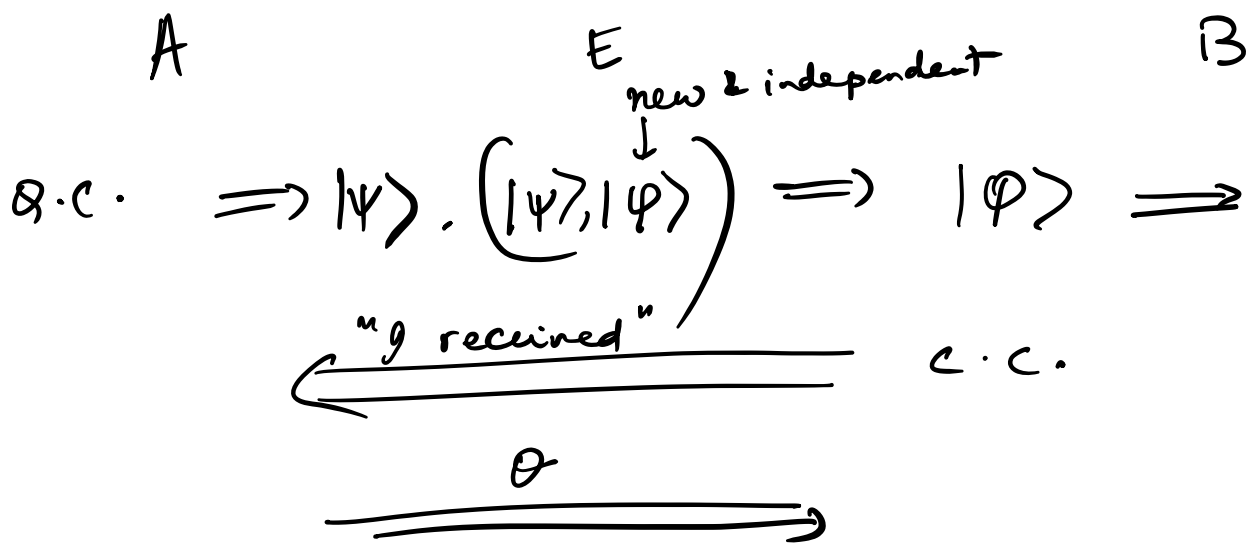
$\theta_i = 1$ "Hadamard basis" = $\{ |+\rangle |-\rangle \}$

If $x_i = 0$, sample the first element in basis θ_i

If $x_i = 1$, sample the second element in basis θ_i

x_i	$\theta_i = 0$	$\theta_i = 1$
0	$ x_i\rangle_{\theta_i} = 0\rangle$	$ +\rangle$
1	$ 1\rangle$	$ -\rangle$
	$ x_i\rangle_{\theta_i}$	

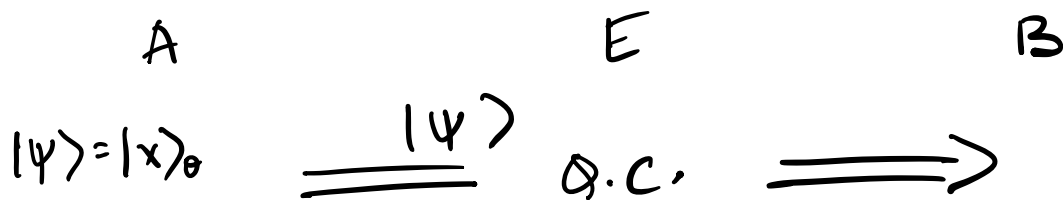
"Conjugate coding"



E measures $|\psi\rangle$ in basis θ .

E and A share a key, E and B share a key.

INSECURE!



← I received

$\theta, S \subseteq [n], \{x_i\}_{i \in S}$

→

$|S| = \frac{n}{2}$

1) $\forall i$ Measure $|\psi_i\rangle$ in θ_i to obtain y_i
 $y = y_1 \dots y_n$

Failed! if not about 2) check if $\forall i \in S, x_i = y_i$
 Passed! 3) Output $\{y_i\}_{i \in [n] \setminus S}$

Output $\{x_i\}_{i \in [n] \setminus S}$

Equivalent experiment

A

EPR pairs

$$\frac{|00\rangle + |11\rangle}{\sqrt{2}} = A_1 B_1$$

⋮

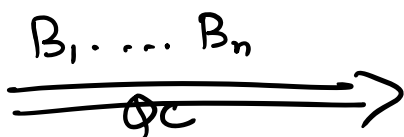
$$= A_n B_n$$

Keeps $A_1 \dots A_n$

$\theta = \theta_1 \dots \theta_n$
measure $A_1 \dots A_n$
in basis $\theta_1 \dots \theta_n$

$$\{x_i\}_{i \in [n] \setminus S}$$

B



$$\frac{1}{2^n} I_n$$

"I received"

$$\{\theta_i\}_{i \in [n]}, S, \{x_i\}_{i \in S}$$

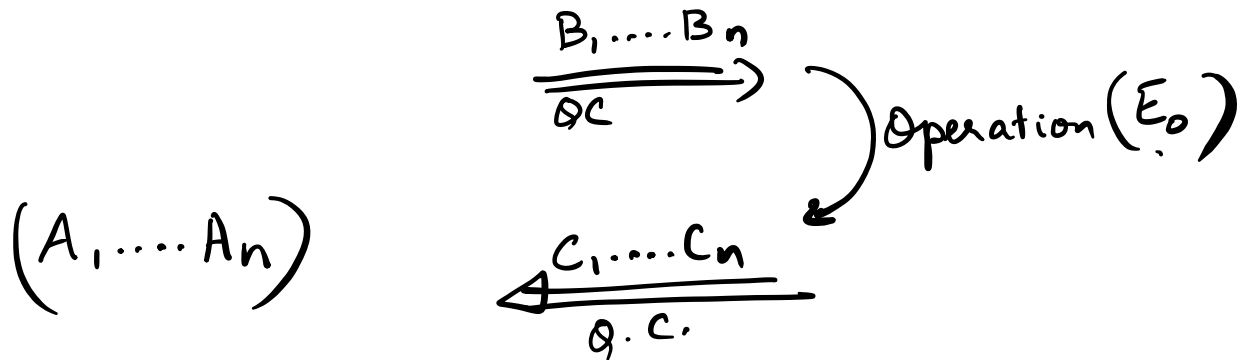
Fails/Passes

measure in θ ,
same check as before.

$$\{y_i\}_{i \in [n] \setminus S}$$

(Alice, Bob)

Eve



C.C. $\mathcal{S}, \{x_i\}_{i \in \mathcal{S}}$

- * Measure A_1, \dots, A_n in basis $\theta_1, \dots, \theta_n \rightarrow X_1, \dots, X_n$
- * Measure C_1, \dots, C_n in basis $\theta_1, \dots, \theta_n$
to obtain Y_1, \dots, Y_n
- * Test if $\{x_i = y_i\}_{i \in \mathcal{S}}$.

Pass / fail \rightarrow

If pass, output $\{x_i\}_{i \in \mathcal{S}}$ as Alice's "key".
 $\mathcal{S} = [n] \setminus \mathcal{S}$
 $\{y_i\}_{i \in \mathcal{S}}$ as Bob's "key".

$|\Psi\rangle_{ACE_0}$

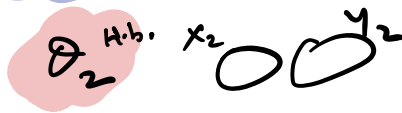
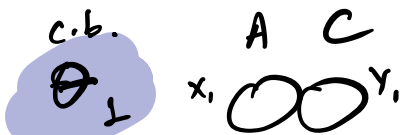
after the check \longrightarrow

$|\Psi\rangle_{\Theta XY E_0}$

Hybrid classical-quantum state

Start with $|\Psi\rangle_{ACE_0}$.

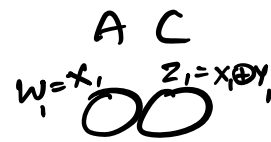
00	\rightarrow	00
01	\rightarrow	01
10	\rightarrow	11
11	\rightarrow	10



...



$\forall i, \text{CNOT}(A_i, C_i)$



...

$$\begin{aligned}
 & H^{\otimes 2} \left(\text{CNOT} \left(H^{\otimes 2} (x_2 y_2) \right) \right) \\
 &= H^{\otimes 2} \left(\text{CNOT} \left(\sum_{\sigma_1, \sigma_2} (-1)^{\langle \sigma_1, \sigma_2, x_2 y_2 \rangle} \right) \left| \sigma_1, \sigma_2 \right\rangle \right) \\
 &= H^{\otimes 2} \left(\sum_{\sigma_1, \sigma_3} (-1)^{\langle \sigma_1, \sigma_3, x_2 y_2 \rangle} \left| \sigma_1, \sigma_3 \right\rangle \right) \\
 &= \sum_{P, Q} \left(\sum_{\sigma_1, \sigma_3} (-1)^{\langle P, Q, \sigma_1, \sigma_3 \rangle} \cdot (-1)^{\langle \sigma_1, \sigma_3 \oplus \sigma_3, x_2 y_2 \rangle} \right) \left| P, Q \right\rangle
 \end{aligned}$$

$$= \sum_{p, q} \left(\sum_{\sigma_1, \sigma_3} (-1)^{\langle pq, \sigma_1, \sigma_3 \rangle} \cdot (-1)^{\langle \sigma_1, \sigma_1 \oplus \sigma_3, x_2 y_2 \rangle} \right) |p, q\rangle$$

$$\begin{aligned} & \sigma_1 x_2 \oplus \sigma_1 y_2 \oplus \sigma_3 y_2 \\ &= \sigma_1 (x_2 \oplus y_2) \oplus \sigma_3 (y_2) \end{aligned}$$

$$\sum_{\sigma_1, \sigma_3} (-1)^{\langle pq, \sigma_1, \sigma_3 \rangle \oplus \langle \sigma_1, \sigma_1 \oplus \sigma_3, x_2 y_2 \rangle}$$


= 0 when the term in exponent is NOT identically 0 $\forall \sigma_1, \sigma_3$.


= 1 $\forall w$.

term in exponent is identically 0 when.

$$\text{identically } 0 \quad \langle pq, \sigma_1, \sigma_3 \rangle = \sigma_1 (x_2 \oplus y_2) \oplus \sigma_3 (y_2)$$

$$\Leftrightarrow p = x_2 \oplus y_2, \quad q = y_2.$$

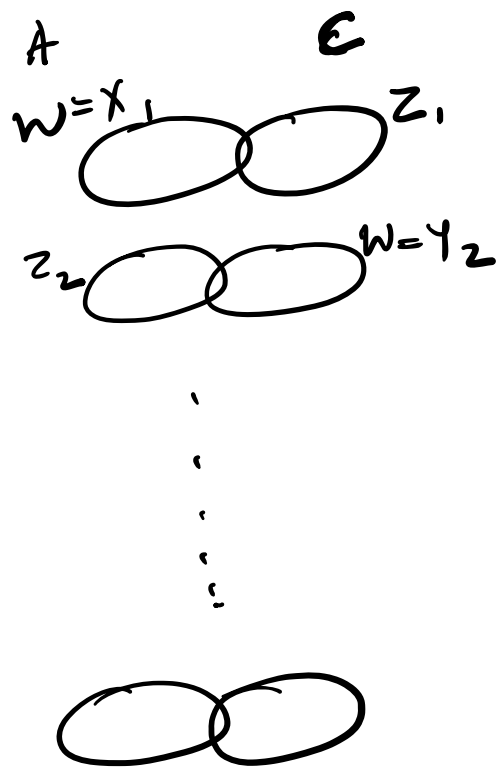
$\theta = 0$ 
c.b.

$\theta = 1$ 









then sample $S \subseteq [n]$.

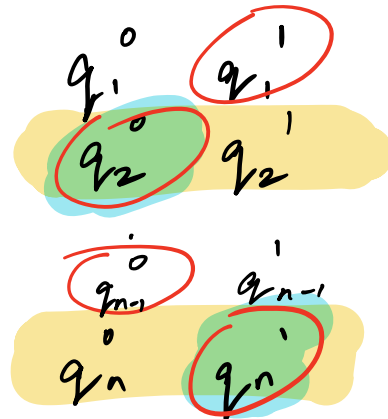
Check if $\sum_{i \in S} z_i = 0$.

If it is 0, then sample $\{w_i\}_{i \in \bar{S}}$.

Classical Sampling and estimation

2n-bit string

Write as $n \times 2$ matrix



$\forall i \in [n]$, sample $j_i \leftarrow \{0, 1\}$.

Let $T = \{(i, j_i)\}_{i \in [n]}$, $\bar{T} = \{(i, 1-j_i)\}_{i \in [n]}$

Sample $S \subseteq T$, then count #1s in q_S
(q size $n/2$) $= W[q_S]$

Use this to estimate $W[q_{\bar{T}}]$. $W[q_{\bar{T}}] \leq 2W[q_S] + \text{smol}$

Step 1: Relate q_T and $q_{\bar{T}}$.

Let $L = \{i: q_i^0 \neq q_i^1\}$ Let $l = |L|$.

$$\begin{aligned} W[q_T] - W[q_{\bar{T}}] &= W[q_{T \setminus L}] - W[q_{\bar{T} \setminus L}] \\ &= 2W[q_{T \setminus L}] - l \end{aligned}$$

$$\therefore W[q_{T \setminus L}] + W[q_{\bar{T} \setminus L}] = l \quad (\text{because 1 is 0, other is 1 in each row})$$

$$\Pr[|W[q_T] - w[q_T]| \geq n\epsilon] \stackrel{(\text{later})}{\leq} 2e^{-n\epsilon^2/2}$$

$$= \Pr[|2W[q_{T|L}] - l| \geq n\epsilon]$$

$$= \Pr[|W[q_{T|L}] - l/2| \geq \frac{n\epsilon}{2}]$$

$q_{T|L}$ is T restricted to positions where $q_i^0 \neq q_i^1$
 So one of them is 0 & the other is 1.

Because T was chosen randomly, every entry in the set $q_{T|L}$ is 0 w.p. $\frac{1}{2}$ & 1 w.p. $\frac{1}{2}$ independently

l such entries, so

$$\mathbb{E}[W[q_{T|L}]] = l/2.$$

Hoeffding's inequality:

For r.v.s X_1, \dots, X_m s.t. $0 \leq X_i \leq 1$,

let $S = X_1 + X_2 + \dots + X_m$

$$\Pr[|S - \mathbb{E}[S]| \geq t] \leq 2e^{-2t^2/m}$$

Substituting $S = W[q_{T|L}]$, $t = \frac{n\epsilon}{2}$, $\Rightarrow 2t^2 = \frac{n^2\epsilon^2}{2}$

$$\Pr[|W[q_{T|L}] - l/2| \geq \frac{n\epsilon}{2}] \leq 2e^{-\frac{n^2\epsilon^2}{2l}} \leq 2e^{-n\epsilon^2/2}$$

Step 2: relate $W[q_T]$ and $W[q_S]$

$$\Pr \left[\left| \frac{W[q_S]}{|S|} - \frac{W[q_T]}{|T|} \right| \geq \delta \right] \leq 2e^{-2\delta^2|S|}$$

Recall,

$$\Pr \left[\left| \frac{W[q_T]}{|T|} - \frac{W[q_{\bar{T}}]}{|T|} \right| \geq \varepsilon \right] \leq 2e^{-n\varepsilon^2/2}$$

$$\Pr \left[\left| \frac{W[q_S]}{|S|} - \frac{W[q_{\bar{T}}]}{|T|} \right| \geq (\varepsilon + \delta) \right]$$

$$\leq 2e^{-\delta^2 n} + 2e^{-\varepsilon^2 n/2}$$

Set ε s.t. $\varepsilon^2/2 = \delta^2$. Then $\leq 4e^{-\delta^2 n}$
 $\varepsilon = \sqrt{2}\delta$

In particular, if $W[q_S] = 0$,

$$w.p. 4e^{-\delta^2 n}, \quad W[q_{\bar{T}}] \leq \delta(1 + \sqrt{2}) \cdot n$$